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STDLIB

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1 User's Guide

The Erlang standard library *STDLIB*.

1.1 The Erlang I/O-protocol

The I/O-protocol in Erlang specifies a way for a client to communicate with an io_server and vice versa. The io_server is a process handling the requests and that performs the requested task on i.e. a device. The client is any Erlang process wishing to read or write data from/to the device.

The common I/O-protocol has been present in OTP since the beginning, but has been fairly undocumented and has also somewhat evolved over the years. In an addendum to Robert Virdings rationale the original I/O-protocol is described. This document describes the current I/O-protocol.

The original I/O-protocol was simple and flexible. Demands for spacial and execution time efficiency has triggered extensions to the protocol over the years, making the protocol larger and somewhat less easy to implement than the original. It can certainly be argumented that the current protocol is to complex, but this text describes how it looks today, not how it should have looked.

The basic ideas from the original protocol still holds. The io_server and client communicate with one single, rather simplistic protocol and no server state is ever present in the client. Any io_server can be used together with any client code and client code need not be aware of the actual device the io_server communicates with.

1.1.1 Protocol basics

As described in Roberts paper, servers and clients communicate using io_request/io_reply tuples as follows:

{io_request, From, ReplyAs, Request}
{io_reply, ReplyAs, Reply}

The client sends an io_request to the io_server and the server eventually sends a corresponding reply.

- From is the pid() of the client, the process which the io_server sends the reply to.
- ReplyAs can be any datum and is simply returned in the corresponding io_reply. The io-module in the Erlang standard library simply uses the pid() of the io_server as the ReplyAs datum, but a more complicated client could have several outstanding io-requests to the same server and would then use i.e. a reference() or something else to differentiate among the incoming io_reply's. The ReplyAs element should be considered opaque by the io_server. Note that the pid() of the server is not explicitly present in the io_reply. The reply can be sent from any process, not necessarily the actual io_server. The ReplyAs element is the only thing that connects one io_request with an io_reply.
- Request and Reply are described below.

When an io_server receives an io_request, it acts upon the actual Request part and eventually sends an io_reply with the corresponding Reply part.

1.1.2 Output requests

To output characters on a device, the following Requests exist:

{put_chars, Encoding, Characters}
{put_chars, Encoding, Module, Function, Args}

- Encoding is either 'latin1' or 'unicode', meaning that the characters are (in case of binaries) encoded as either UTF-8 or iso-latin-1 (pure bytes). A well behaved io_server should also return error if list elements contain integers > 255 when the Encoding is set to latin1. Note that this does not in any way tell how characters should be put on the actual device or how the io_server should handle them. Different io_servers may handle the characters however they want, this simply tells the io_server which format the data is expected to have. In the Module/Function/argument case, the Encoding tells which format the designated function produces. Note that byte-oriented data is simplest sent using latin1 Encoding
- Characters are the data to be put on the device. If encoding is latin1, this is an iolist(). If encoding is unicode, this is an Erlang standard mixed unicode list (one integer in a list per character, characters in binaries represented as UTF-8).
- Module, Function, Args denotes a function which will be called to produce the data (like io_lib:format), Args is a list of arguments to the function. The function should produce data in the given Encoding. The io_server should call the function as apply(Mod, Func, Args) and will put the returned data on the device as if it was sent in a {put_chars, Encoding, Characters} request. If the function returns anything else than a binary or list or throws an exception, an error should be sent back to the client.

The server replies to the client with an io_reply where the Reply element is one of:

ok
{error, Error}

- Error describes the error to the client, which may do whatever it wants with it. The Erlang io-module typically returns it as is.

For backward compatibility the following Requests should also be handled by an io_server (these messages should not be present after R15B of OTP):

{put_chars, Characters}
{put_chars, Module, Function, Args}

These should behave as {put_chars, latin1, Characters} and {put_chars, latin1, Module, Function, Args} respectively.

1.1.3 Input Requests

To read characters from a device, the following Requests exist:

{get_until, Encoding, Prompt, Module, Function, ExtraArgs}

- Encoding denotes how data is to be sent back to the client and what data is sent to the function denoted by Module/Function/Arity. If the function supplied returns data as a list, the data is converted to this encoding. If however the function supplied returns data in some other format, no conversion can be done and it's up to the client supplied function to return data in a proper way. If Encoding is latin1, lists of integers 0..255 or binaries containing plain bytes are sent back to the client when possible, if Encoding is unicode, lists with integers in the whole unicode range or binaries encoded in UTF-8 are sent to the client. The user supplied function will always see lists of integers, never binaries, but the list may contain numbers > 255 if the Encoding is 'unicode'.
- Prompt is a list of characters (not mixed, no binaries) or an atom() to be output as a prompt for input on the device. The Prompt is often ignored by the io_server and a Prompt set to '' should always be ignored (and result in nothing being written to the device).
- Module, Function, ExtraArgs denotes a function and arguments to determine when enough data is written. The function should take two additional arguments, the last state, and a list of characters. The function should return one of:

  {done, Result, RestChars}
  {more, Continuation}

  The Result can be any Erlang term, but if it is a list(), the io_server may convert it to a binary() of appropriate format before returning it to the client, if the server is set in binary mode (see below).

The function will be called with the data the io_server finds on its device, returning {done, Result, RestChars} when enough data is read (in which case Result is sent to the client and RestChars are kept in the io_server as a buffer for subsequent input) or {more, Continuation}, indicating that more characters are needed to complete the request. The Continuation will be sent as the state in subsequent calls to the function when more characters are available. When no more characters are available, the function shall return {done,eof,Rest}. The initial state is the empty list and the data when an end of file is reached on the device is the atom 'eof'. An emulation of the get_line request could be (inefficiently) implemented using the following functions:

```
-module(demo).
-export([until_newline/3, get_line/1]).

until_newline(_ThisFar,eof,_MyStopCharacter) ->
    {done,eof,[]};
until_newline(ThisFar,CharList,MyStopCharacter) ->
    case lists:splitwith(fun(X) -> X =/= MyStopCharacter end,  CharList) of
 {L,[]} ->
            {more,ThisFar++L};
 {L2,[MyStopCharacter|Rest]} ->
     {done,ThisFar++L2++[MyStopCharacter],Rest}
    end.

get_line(IoServer) ->
    IoServer ! {io_request, self(), IoServer, {get_until, unicode, '',
                                               ?MODULE, until_newline, [$\n]}},
    receive
        {io_reply, IoServer, Data} ->
     Data
    end.
```

Note especially that the last element in the Request tuple ([$\n]) is appended to the argument list when the function is called. The function should be called like apply(Module, Function, [State, Data | ExtraArgs]) by the io_server

A defined number of characters is requested using this Request:

{get_chars, Encoding, Prompt, N}

- Encoding and Prompt as for get_until.
- N is the number of characters to be read from the device.

A single line (like in the example above) is requested with this Request:

{get_line, Encoding, Prompt}

- Encoding and prompt as above.

Obviously, get_chars and get_line could be implemented with the get_until request (and indeed was originally), but demands for efficiency has made these additions necessary.

The server replies to the client with an io_reply where the Reply element is one of:

Data
eof
{error, Error}

- Data is the characters read, in either list or binary form (depending on the io_server mode, see below).
- Error describes the error to the client, which may do whatever it wants with it. The Erlang io-module typically returns it as is.
- eof is returned when input end is reached and no more data is available to the client process.

For backward compatibility the following Requests should also be handled by an io_server (these messages should not be present after R15B of OTP):

{get_until, Prompt, Module, Function, ExtraArgs}
{get_chars, Prompt, N}
{get_line, Prompt}

These should behave as {get_until, latin1, Prompt, Module, Function, ExtraArgs}, {get_chars, latin1, Prompt, N} and {get_line, latin1, Prompt} respectively.

1.1.4 I/O-server modes

Demands for efficiency when reading data from an io_server has not only lead to the addition of the get_line and get_chars requests, but has also added the concept of io_server options. No options are mandatory to implement, but all io_servers in the Erlang standard libraries honor the 'binary' option, which allows the Data in the io_reply to be binary instead of in list form *when possible*. If the data is sent as a binary, Unicode data will be sent in the standard Erlang unicode format, i.e. UTF-8 (note that the function in get_until still gets list data regardless of the io_server mode).

Note that i.e. the `get_until` request allows for a function with the data specified as always being a list. Also the return value data from such a function can be of any type (as is indeed the case when an io:fread request is sent to an io_server). The client has to be prepared for data received as answers to those requests to be in a variety of forms, but the server should convert the results to binaries whenever possible (i.e. when the function supplied to get_until actually returns a list). The example shown later in this text does just that.

An I/O-server in binary mode will affect the data sent to the client, so that it has to be able to handle binary data. For convenience, it is possible to set and retrieve the modes of an io_server using the following I/O-requests:

{setopts, Opts}

- Opts is a list of options in the format recognized by proplists (and of course by the io_server itself).

As an example, the io_server for the interactive shell (in group.erl) understands the following options:

{binary, bool()} (or 'binary'/'list')
{echo, bool()}
{expand_fun, fun()}
{encoding, 'unicode'/'latin1'} (or 'unicode'/'latin1')

- of which the 'binary' and 'encoding' options are common for all io_servers in OTP, while 'echo' and 'expand' is valid only for this io_server. It's worth noting that the 'unicode' option notifies how characters are actually put on the physical device, i.e. if the terminal per se is unicode aware, it does not affect how characters are sent in the I/O-protocol, where each request contains encoding information for the provided or returned data.

The server should send one of the following as Reply:

ok
{error, Error}

An error (preferably enotsup) is to be expected if the option is not supported by the io_server (like if an 'echo' option is sent in a setopt Request to a plain file).

To retrieve options, this message is used:

getopts

The 'getopts' message requests a complete list of all options supported by the io_server as well as their current values.

The server replies:

OptList
{error,Error}

- OptList is a list of tuples {Option, Value} where Option is always an atom.

1.1.5 Multiple I/O requests

The Request element can in itself contain several Requests by using the following format:

{requests, Requests}

- Requests is a list of valid Request tuples for the protocol, they shall be executed in the order in which they appear in the list and the execution should continue until one of the requests result in an error or the list is consumed. The result of the last request is sent back to the client.

The server can for a list of requests send any of the valid results in the reply:

ok
{ok, Data}
{ok, Options}
{error, Error}

- depending on the actual requests in the list.

1.1.6 Optional I/O-requests

The following I/O request is optional to implement and a client should be prepared for an error return:

{get_geometry, Geometry}

- Geometry is either the atom 'rows' or the atom 'columns'.

The server should send the Reply as:

{ok, N}
{error, Error}

- N is the number of character rows or columns the device has, if applicable to the device the io_server handles, otherwise {error, enotsup} is a good answer.

1.1.7 Unimplemented request types:

If an io_server encounters a request it does not recognize (i.e. the io_request tuple is in the expected format, but the actual Request is unknown), the server should send a valid reply with the error tuple:

{error, request}

This makes it possible to extend the protocol with optional messages and for the clients to be somewhat backwards compatible.

1.1.8 An annotated and working example io_server:

An io_server is any process capable of handling the protocol. There is no generic io_server behavior, but could well be. The framework is simple enough, a process handling incoming requests, usually both io_requests and other device-specific requests (for i.e. positioning , closing etc.).

Our example io_server stores characters in an ets table, making up a fairly crude ram-file (it is probably not useful, but working).

The module begins with the usual directives, a function to start the server and a main loop handling the requests:

```
-module(ets_io_server).

-export([start_link/0, init/0, loop/1, until_newline/3, until_enough/3]).

-define(CHARS_PER_REC, 10).
```

```
-record(state, {
   table,
   position, % absolute
   mode % binary | list
  }).

start_link() ->
    spawn_link(?MODULE,init,[]).

init() ->
    Table = ets:new(noname,[ordered_set]),
    ?MODULE:loop(#state{table = Table, position = 0, mode=list}).

loop(State) ->
    receive
 {io_request, From, ReplyAs, Request} ->
     case request(Request,State) of
  {Tag, Reply, NewState} when Tag =:= ok; Tag =:= error ->
      reply(From, ReplyAs, Reply),
      ?MODULE:loop(NewState);
  {stop, Reply, _NewState} ->
      reply(From, ReplyAs, Reply),
      exit(Reply)
     end;
 %% Private message
 {From, rewind} ->
     From ! {self(), ok},
     ?MODULE:loop(State#state{position = 0});
 _Unknown ->
     ?MODULE:loop(State)
    end.
```

The main loop receives messages from the client (which might be using the io-module to send requests). For each request the function request/2 is called and a reply is eventually sent using the reply/3 function.

The "private" message {From, rewind} results in the current position in the pseudo-file to be reset to 0 (the beginning of the "file"). This is a typical example of device-specific messages not being part of the I/O-protocol. It is usually a bad idea to embed such private messages in io_request tuples, as that might be confusing to the reader.

Let's look at the reply function first...

```
reply(From, ReplyAs, Reply) ->
    From ! {io_reply, ReplyAs, Reply}.
```

Simple enough, it sends the io_reply tuple back to the client, providing the ReplyAs element received in the request along with the result of the request, as described above.

Now look at the different requests we need to handle. First the requests for writing characters:

```
request({put_chars, Encoding, Chars}, State) ->
    put_chars(unicode:characters_to_list(Chars,Encoding),State);
request({put_chars, Encoding, Module, Function, Args}, State) ->
    try
 request({put_chars, Encoding, apply(Module, Function, Args)}, State)
    catch
 _:_ ->
     {error, {error,Function}, State}
```

```
    end;
```

The Encoding tells us how the characters in the message are represented. We want to store the characters as lists in the ets-table, so we convert them to lists using the unicode:characters_to_list/2 function. The conversion function conveniently accepts the encoding types unicode or latin1, so we can use the Encoding parameter directly.

When Module, Function and Arguments are provided, we simply apply it and do the same thing with the result as if the data was provided directly.

Let's handle the requests for retrieving data too:

```
request({get_until, Encoding, _Prompt, M, F, As}, State) ->
    get_until(Encoding, M, F, As, State);
request({get_chars, Encoding, _Prompt, N}, State) ->
    %% To simplify the code, get_chars is implemented using get_until
    get_until(Encoding, ?MODULE, until_enough, [N], State);
request({get_line, Encoding, _Prompt}, State) ->
    %% To simplify the code, get_line is implemented using get_until
    get_until(Encoding, ?MODULE, until_newline, [$\n], State);
```

Here we have cheated a little by more or less only implementing get_until and using internal helpers to implement get__chars and get_line. In production code, this might be to inefficient, but that of course depends on the frequency of the different requests. Before we start actually implementing the functions put_chars/2 and get_until/5, lets look into the few remaining requests:

```
request({get_geometry,_}, State) ->
    {error, {error,enotsup}, State};
request({setopts, Opts}, State) ->
    setopts(Opts, State);
request(getopts, State) ->
    getopts(State);
request({requests, Reqs}, State) ->
     multi_request(Reqs, {ok, ok, State});
```

The get_geometry request has no meaning for this io_server, so the reply will be {error, enotsup}. The only option we handle is the binary/list option, which is done in separate functions.

The multi-request tag (requests) is handled in a separate loop function applying the requests in the list one after another, returning the last result.

What's left is to handle backward compatibility and the file-module (which uses the old requests until backward compatibility with pre-R13 nodes is no longer needed). Note that the io_server will not work with a simple file:write if these are not added:

```
request({put_chars,Chars}, State) ->
    request({put_chars,latin1,Chars}, State);
request({put_chars,M,F,As}, State) ->
    request({put_chars,latin1,M,F,As}, State);
request({get_chars,Prompt,N}, State) ->
    request({get_chars,latin1,Prompt,N}, State);
request({get_line,Prompt}, State) ->
    request({get_line,latin1,Prompt}, State);
request({get_until, Prompt,M,F,As}, State) ->
    request({get_until,latin1,Prompt,M,F,As}, State);
```

Ok, what's left now is to return {error, request} if the request is not recognized:

```
request(_Other, State) ->
    {error, {error, request}, State}.
```

Let's move further and actually handle the different requests, first the fairly generic multi-request type:

```
multi_request([R|Rs], {ok, _Res, State}) ->
    multi_request(Rs, request(R, State));
multi_request([_|_], Error) ->
    Error;
multi_request([], Result) ->
    Result.
```

We loop through the requests one at the time, stopping when we either encounter an error or the list is exhausted. The last return value is sent back to the client (it's first returned to the main loop and then sent back by the function io_reply).

The getopt and setopt requests is also simple to handle, we just change or read our state record:

```
setopts(Opts0,State) ->
    Opts = proplists:unfold(
      proplists:substitute_negations(
        [{list,binary}],
        Opts0)),
    case check_valid_opts(Opts) of
 true ->
         case proplists:get_value(binary, Opts) of
      true ->
   {ok,ok,State#state{mode=binary}};
      false ->
   {ok,ok,State#state{mode=binary}};
      _ ->
   {ok,ok,State}
  end;
 false ->
     {error,{error,enotsup},State}
    end.
check_valid_opts([]) ->
    true;
check_valid_opts([{binary,Bool}|T]) when is_boolean(Bool) ->
    check_valid_opts(T);
check_valid_opts(_) ->
    false.

getopts(#state{mode=M} = S) ->
    {ok,[{binary, case M of
        binary ->
     true;
        _ ->
     false
    end}],S}.
```

As a convention, all io_servers handle both {setopts, [binary]}, {setopts, [list]} and {setopts,[{binary, bool()}]}, hence the trick with proplists:substitute_negations/2 and proplists:unfold/1. If invalid options are sent to us, we send {error,enotsup} back to the client.

The getopts request should return a list of {Option, Value} tuples, which has the twofold function of providing both the current values and the available options of this io_server. We have only one option, and hence return that.

So far our io_server has been fairly generic (except for the rewind request handled in the main loop and the creation of an ets table). Most io_servers contain code similar to what's above.

To make the example runnable, we now start implementing the actual reading and writing of the data to/from the ets-table. First the put_chars function:

```
put_chars(Chars, #state{table = T, position = P} = State) ->
    R = P div ?CHARS_PER_REC,
    C = P rem ?CHARS_PER_REC,
    [ apply_update(T,U) || U <- split_data(Chars, R, C) ],
    {ok, ok, State#state{position = (P + length(Chars))}}.
```

We already have the data as (Unicode) lists and therefore just split the list in runs of a predefined size and put each run in the table at the current position (and forward). The functions split_data/3 and apply_update/2 are implemented below.

Now we want to read data from the table. The get_until function reads data and applies the function until it says it's done. The result is sent back to the client:

```
get_until(Encoding, Mod, Func, As,
   #state{position = P, mode = M, table = T} = State) ->
    case get_loop(Mod,Func,As,T,P,[]) of
 {done,Data,_,NewP} when is_binary(Data); is_list(Data) ->
     if
  M =:= binary ->
      {ok,
       unicode:characters_to_binary(Data,unicode,Encoding),
       State#state{position = NewP}};
  true ->
      case check(Encoding,
    unicode:characters_to_list(Data, unicode)) of
   {error, _} = E ->
       {error, E, State};
   List ->
       {ok, List,
        State#state{position = NewP}}
      end
     end;
 {done,Data,_,NewP} ->
     {ok, Data, State#state{position = NewP}};
 Error ->
     {error, Error, State}
    end.

get_loop(M,F,A,T,P,C) ->
    {NewP,L} = get(P,T),
    case catch apply(M,F,[C,L|A]) of
 {done, List, Rest} ->
     {done, List, [], NewP - length(Rest)};
 {more, NewC} ->
     get_loop(M,F,A,T,NewP,NewC);
 _ ->
     {error,F}
    end.
```

Here we also handle the mode (binary or list) that can be set by the setopts request. By default, all OTP io_servers send data back to the client as lists, but switching mode to binary might increase efficiency if the server handles it in an

appropriate way. The implementation of get_until is hard to get efficient as the supplied function is defined to take lists as arguments, but get_chars and get_line can be optimized for binary mode. This example does not optimize anything however. It is important though that the returned data is of the right type depending on the options set, so we convert the lists to binaries in the correct encoding *if possible* before returning. The function supplied in the get_until request may, as its final result return anything, so only functions actually returning lists can get them converted to binaries. If the request contained the encoding tag unicode, the lists can contain all unicode codepoints and the binaries should be in UTF-8, if the encoding tag was latin1, the client should only get characters in the range 0..255. The function check/2 takes care of not returning arbitrary unicode codepoints in lists if the encoding was given as latin1. If the function did not return a list, the check cannot be performed and the result will be that of the supplied function untouched.

Now we are more or less done. We implement the utility functions below to actually manipulate the table:

```
check(unicode, List) ->
    List;
check(latin1, List) ->
    try
 [ throw(not_unicode) || X <- List,
    X > 255 ],
 List
    catch
 throw:_ ->
     {error,{cannot_convert, unicode, latin1}}
    end.
```

The function check takes care of providing an error tuple if unicode codepoints above 255 is to be returned if the client requested latin1.

The two functions until_newline/3 and until_enough/3 are helpers used together with the get_until function to implement get_chars and get_line (inefficiently):

```
until_newline([],eof,_MyStopCharacter) ->
    {done,eof,[]};
until_newline(ThisFar,eof,_MyStopCharacter) ->
    {done,ThisFar,[]};
until_newline(ThisFar,CharList,MyStopCharacter) ->
    case lists:splitwith(fun(X) -> X =/= MyStopCharacter end,  CharList) of
 {L,[]} ->
           {more,ThisFar++L};
 {L2,[MyStopCharacter|Rest]} ->
     {done,ThisFar++L2++[MyStopCharacter],Rest}
    end.

until_enough([],eof,_N) ->
    {done,eof,[]};
until_enough(ThisFar,eof,_N) ->
    {done,ThisFar,[]};
until_enough(ThisFar,CharList,N)
  when length(ThisFar) + length(CharList) >= N ->
    {Res,Rest} = my_split(N,ThisFar ++ CharList, []),
    {done,Res,Rest};
until_enough(ThisFar,CharList,_N) ->
    {more,ThisFar++CharList}.
```

As can be seen, the functions above are just the type of functions that should be provided in get_until requests.

Now we only need to read and write the table in an appropriate way to complete the server:

```
get(P,Tab) ->
    R = P div ?CHARS_PER_REC,
    C = P rem ?CHARS_PER_REC,
    case ets:lookup(Tab,R) of
 [] ->
     {P,eof};
 [{R,List}] ->
     case my_split(C,List,[]) of
  {_,[]} ->
      {P+length(List),eof};
  {_,Data} ->
      {P+length(Data),Data}
     end
    end.

my_split(0,Left,Acc) ->
    {lists:reverse(Acc),Left};
my_split(_,[],Acc) ->
    {lists:reverse(Acc),[]};
my_split(N,[H|T],Acc) ->
    my_split(N-1,T,[H|Acc]).

split_data([],_,_) ->
    [];
split_data(Chars, Row, Col) ->
    {This,Left} = my_split(?CHARS_PER_REC - Col, Chars, []),
    [ {Row, Col, This} | split_data(Left, Row + 1, 0) ].

apply_update(Table, {Row, Col, List}) ->
    case ets:lookup(Table,Row) of
 [] ->
     ets:insert(Table,{Row, lists:duplicate(Col,0) ++ List});
 [{Row, OldData}] ->
     {Part1,_} = my_split(Col,OldData,[]),
     {_,Part2} = my_split(Col+length(List),OldData,[]),
     ets:insert(Table,{Row, Part1 ++ List ++ Part2})
    end.
```

The table is read or written in chunks of ?CHARS_PER_REC, overwriting when necessary. The implementation is obviously not efficient, it is just working.

This concludes the example. It is fully runnable and you can read or write to the io_server by using i.e. the io_module or even the file module. It's as simple as that to implement a fully fledged io_server in Erlang.

1.2 Using Unicode in Erlang

Implementing support for Unicode character sets is an ongoing process. The Erlang Enhancement Proposal (EEP) 10 outlines the basics of Unicode support and also specifies a default encoding in binaries that all Unicode-aware modules should handle in the future.

The functionality described in EEP10 is implemented in Erlang/OTP as of R13A, but that's by no means the end of it. More functionality will be needed in the future and more OTP-libraries might need updating to cope with Unicode data. One example of future development is obvious when reading this manual, our documentation format is limited to the ISO-latin-1 character range, why no Unicode characters beyond that range will occur in this document.

This guide outlines the current Unicode support and gives a couple of recipes for working with Unicode data.

Erlang对Unicode并不是完全支持

unicode是

1.2.1 What Unicode is

Unicode is a standard defining codepoints (numbers) for all known, living or dead, scripts. In principle, every known symbol used in any language has a Unicode codepoint.

Unicode codepoints are defined and published by the *Unicode Consortium*, which is a non profit organization.

Support for Unicode is increasing throughout the world of computing, as the benefits of one common character set are overwhelming when programs are used in a global environment.

Along with the base of the standard, the codepoints for all the scripts, there are a couple of encoding standards available. Different operating systems and tools support different encodings. For example Linux and MacOS X has chosen the UTF-8 encoding, which is backwards compatible with 7-bit ASCII and therefore affects programs written in plain English the least. Windows® on the other hand supports a limited version of UTF-16, namely all the code planes where the characters can be stored in one single 16-bit entity, which includes most living languages.

The most widely spread encodings are:

UTF-8
: Each character is stored in one to four bytes depending on codepoint. The encoding is backwards compatible with 7-bit ASCII as all 7-bit characters are stored in one single byte as is. The characters beyond codepoint 127 are stored in more bytes, letting the most significant bit in the first character indicate a multi-byte character. For details on the encoding, the RFC is publicly available.

UTF-16
: This encoding has many similarities to UTF-8, but the basic unit is a 16-bit number. This means that all characters occupy at least two bytes, some high numbers even four bytes. Some programs and operating systems claiming to use UTF-16 only allows for characters that can be stored in one 16-bit entity, which is usually sufficient to handle living languages. As the basic unit is more than one byte, byte-order issues occur, why UTF-16 exists in both a big-endian and little-endian variant.

UTF-32
: The most straight forward representation, each character is stored in one single 32-bit number. There is no need for escapes or any variable amount of entities for one character, all Unicode codepoints can be stored in one single 32-bit entity. As with UTF-16, there are byte-order issues, UTF-32 can be both big- and little-endian.

UCS-4
: Basically the same as UTF-32, but without some Unicode semantics, defined by IEEE and has little use as a separate encoding standard. For all normal (and possibly abnormal) usages, UTF-32 and UCS-4 are interchangeable.

Certain ranges of characters are left unused and certain ranges are even deemed invalid. The most notable invalid range is 16#D800 - 16#DFFF, as the UTF-16 encoding does not allow for encoding of these numbers. It can be speculated that the UTF-16 encoding standard was, from the beginning, expected to be able to hold all Unicode characters in one 16-bit entity, but then had to be extended, leaving a whole in the Unicode range to cope with backward compatibility.

Additionally, the codepoint 16#FEFF is used for byte order marks (BOM's) and use of that character is not encouraged in other contexts than that. It actually is valid though, as the character "ZWNBS" (Zero Width Non Breaking Space). BOM's are used to identify encodings and byte order for programs where such parameters are not known in advance. Byte order marks are more seldom used than one could expect, put their use is becoming more widely spread as they provide the means for programs to make educated guesses about the Unicode format of a certain file.

1.2.2 Standard Unicode representation in Erlang

In Erlang, strings are actually lists of integers. A string is defined to be encoded in the ISO-latin-1 (ISO8859-1) character set, which is, codepoint by codepoint, a sub-range of the Unicode character set.

The standard list encoding for strings is therefore easily extendible to cope with the whole Unicode range: A Unicode string in Erlang is simply a list containing integers, each integer being a valid Unicode codepoint and representing one character in the Unicode character set.

Erlang的string实际上就是整数元素组成的list,注意string的encode是使用ISO-latin-1字符集,这个字符集是Unicode的子集.
标准的list做encoding很容易扩展到整个unicode编码:编码是整数和字符的对应关系,只要list中的整函数是有效的Unicode codepoint
就可以找到对应的字符; Erlang的string实际上就是整数元素组成的list,注意string的encode是使用ISO-latin-1字符集,
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字符的对应关系,只要list中的整函数是有效的Unicode codepoint就可以找到对应的字符;

二进制数据处理起来就麻烦一些了,二进制数据是紧凑排列的:一个字节代表一个字符,而不是两个字(word)一个字符,
这里如果存在疑问可以查看erlang官方文档中关于内存消耗的列表:
http://www.erlang.org/doc/efficiency_guide/advanced.html .
平时我们使用的erlang:list_to_binary,常规的Erlang string(ISO-latin编码的string)
可以逐字节逐字符顺利转成binary.但是超出ISO-latin编码范围就会出错了,看下面的例子:

```
Eshell V5.9  (abort with ^G)
1> L=[10,12,23,45].
[10,12,23,45]
2> list_to_binary(L).
<<10,12,23,45>>
3> L2= "中国".
[20013,22269]
4>4> list_to_binary(L2).
** exception error: bad argument
     in function  list_to_binary/1
        called as list_to_binary([20013,22269])
5>
```


Regular Erlang strings in ISO-latin-1 are a subset of there Unicode strings.

Binaries on the other hand are more troublesome. For performance reasons, programs often store textual data in binaries instead of lists, mainly because they are more compact (one byte per character instead of two words per character, as is the case with lists). Using erlang:list_to_binary/1, a regular Erlang string can be converted into a binary, effectively using the ISO-latin-1 encoding in the binary - one byte per character. This is very convenient for those regular Erlang strings, but cannot be done for Unicode lists.

As the UTF-8 encoding is widely spread and provides the most compact storage, it is selected as the standard encoding of Unicode characters in binaries for Erlang.

The standard binary encoding is used whenever a library function in Erlang should cope with Unicode data in binaries, but is of course not enforced when communicating externally. Functions and bit-syntax exist to encode and decode both UTF-8, UTF-16 and UTF-32 in binaries. Library functions dealing with binaries and Unicode in general, however, only deal with the default encoding.

Character data may be combined from several sources, sometimes available in a mix of strings and binaries. Erlang has for long had the concept of iodata or iolists, where binaries and lists can be combined to represent a sequence of bytes. In the same way, the Unicode aware modules often allow for combinations of binaries and lists where the binaries have characters encoded in UTF-8 and the lists contain such binaries or numbers representing Unicode codepoints:

```
unicode_binary() = binary() with characters encoded in UTF-8 coding standard
unicode_char() = integer() representing valid unicode codepoint

chardata() = charlist() | unicode_binary()

charlist() = [unicode_char() | unicode_binary() | charlist()]
  a unicode_binary is allowed as the tail of the list
```

The module `unicode` in stdlib even supports similar mixes with binaries containing other encodings than UTF-8, but that is a special case to allow for conversions to and from external data:

```
external_unicode_binary() = binary() with characters coded in a user specified Unicode
  encoding other than UTF-8 (UTF-16 or UTF-32)

external_chardata() = external_charlist() | external_unicode_binary()

external_charlist() = [unicode_char() | external_unicode_binary() | external_charlist()]
  an external_unicode_binary is allowed as the tail of the list
```

1.2.3 Basic language support for Unicode

First of all, Erlang is still defined to be written in the ISO-latin-1 character set. Functions have to be named in that character set, atoms are restricted to ISO-latin-1 and regular strings are still lists of characters 0..255 in the ISO-latin-1 encoding. This has not (yet) changed, but the language has been slightly extended to cope with Unicode characters and encodings.

Bit-syntax

The bit-syntax contains types for coping with binary data in the three main encodings. The types are named `utf8`, `utf16` and `utf32` respectively. The `utf16` and `utf32` types can be in a big- or little-endian variant:

```
<<Ch/utf8,_/binary>> = Bin1,
<<Ch/utf16-little,_/binary>> = Bin2,
```

```
Bin3 = <<$H/utf32-little, $e/utf32-little, $l/utf32-little, $l/utf32-little,
              $o/utf32-little>>,
```

For convenience, literal strings can be encoded with a Unicode encoding in binaries using the following (or similar) syntax:

```
Bin4 = <<"Hello"/utf16>>,
```

String- and character-literals

> ### Warning:
>
> The literal syntax described here may be subject to change in R13B, it has not yet passed the usual process for language changes approval.

It is convenient to be able to write a list of Unicode characters in the string syntax. However, the language specifies strings as being in the ISO-latin-1 character set which the compiler tool chain as well as many other tools expect.

Also the source code is (for now) still expected to be written using the ISO-latin-1 character set, why Unicode characters beyond that range cannot be entered in string literals.

Shell

To make it easier to enter Unicode characters in the shell, it allows strings with Unicode characters on input, immediately converting them to regular lists of integers. They will, by the evaluator etc be viewed as if they were input using the regular list syntax, which is - in the end - how the language actually treats them. They will in the same way not be output as strings by i.e `io:write/2` or `io:format/3` unless the format string supplied to `io:format` uses the Unicode translation modifier (which we will talk about later).

For source code, there is an extension to the \OOO (backslash followed by three octal numbers) and \xHH (backslash followed by 'x', followed by two hexadecimal characters) syntax, namely \x{H ...} (a backslash followed by an 'x', followed by left curly bracket, any number of hexadecimal digits and a terminating right curly bracket). This allows for entering characters of any codepoint literally in a string. The string is immediately converted into a list by the scanner however, which is obvious when calling it directly:

```
1> erl_scan:string("\"X\".").
{ok,[{string,1,"X"},{dot,1}],1}
2> erl_scan:string("\"\x{400}\".").
{ok,[{'[',1},{integer,1,1024},{']',1},{dot,1}],1}
```

Character literals, or rather integers representing Unicode codepoints can be expressed in a similar way using $ \x{H ...}:

```
4> $\x{400}.
1024
```

This also is a translation by the scanner:

```
5> erl_scan:string("$Y.").
{ok,[{char,1,89},{dot,1}],1}
```

```
6> erl_scan:string("$\x{400}.").
{ok,[{integer,1,1024},{dot,1}],1}
```

In the shell, if using a Unicode input device, '$' can be followed directly by a Unicode character producing an integer. In the following example, let's imagine the character 'c' is actually a Cyrillic 's' (looking fairly similar):

```
7> $c.
1089
```

The literal syntax allowing Unicode characters is to be viewed as "syntactic sugar", but is, as such, fairly useful.

1.2.4 The interactive shell

The interactive Erlang shell, when started towards a terminal or started using the `werl` command on windows, can support Unicode input and output.

On Windows®, proper operation requires that a suitable font is installed and selected for the Erlang application to use. If no suitable font is available on your system, try installing the DejaVu fonts (dejavu-fonts.org), which are freely available and then select that font in the Erlang shell application.

On Unix®-like operating systems, the terminal should be able to handle UTF-8 on input and output (modern versions of XTerm, KDE konsole and the Gnome terminal do for example) and your locale settings have to be proper. As an example, my LANG environment variable is set as this:

```
$ echo $LANG
en_US.UTF-8
```

Actually, most systems handle the LC_CTYPE variable before LANG, so if that is set, it has to be set to UTF-8:

```
$ echo $LC_CTYPE
en_US.UTF-8
```

The LANG or LC_CTYPE setting should be consistent with what the terminal is capable of, there is no portable way for Erlang to ask the actual terminal about its UTF-8 capacity, we have to rely on the language and character type settings.

To investigate what Erlang thinks about the terminal, the `io:getopts()` call can be used when the shell is started:

```
$ LC_CTYPE=en_US.ISO-8859-1 erl
Erlang R13A (erts-5.7) [source] [64-bit] [smp:4:4] [rq:4] [async-threads:0] [kernel-poll:false]

Eshell V5.7  (abort with ^G)
1> lists:keyfind(encoding,1,io:getopts()).
{encoding,latin1}
2> q().
ok
$ LC_CTYPE=en_US.UTF-8 erl
Erlang R13A (erts-5.7) [source] [64-bit] [smp:4:4] [rq:4] [async-threads:0] [kernel-poll:false]

Eshell V5.7  (abort with ^G)
1> lists:keyfind(encoding,1,io:getopts()).
{encoding,unicode}
2>
```

When (finally?) everything is in order with the locale settings, fonts and the terminal emulator, you probably also have discovered a way to input characters in the script you desire. For testing, the simplest way is to add some keyboard mappings for other languages, usually done with some applet in your desktop environment. In my KDE environment, I start the KDE Control Center (Personal Settings), select "Regional and Accessibility" and then "Keyboard Layout". On Windows XP®, I start Control Panel->Regional and Language Options, select the Language tab and click the Details... button in the square named "Text services and input Languages". Your environment probably provides similar means of changing the keyboard layout. Make sure you have a way to easily switch back and forth between keyboards if you are not used to this, entering commands using a Cyrillic character set is, as an example, not easily done in the Erlang shell.

Now you are set up for some Unicode input and output. The simplest thing to do is of course to enter a string in the shell:

```
$ erl
Erlang R13A (erts-5.7) [source] [64-bit] [smp:4:4] [rq:4] [async-threads:0] [ker
nel-poll:false]

Eshell V5.7  (abort with ^G)
1> lists:keyfind(encoding,1,io:getopts()).
{encoding,unicode}
2> "уницоде".
[1091,1085,1080,1094,1086,1076,1077]
3> io:format("~ts~n",[v(2)]).
уницоде
ok
4>
```

Figure 2.1: Cyrillic characters in an Erlang shell

While strings can be input as Unicode characters, the language elements are still limited to the ISO-latin-1 character set. Only character constants and strings are allowed to be beyond that range:

```
Eshell V5.7  (abort with ^G)
1> $ξ.
958
2> θνιψοβε.
* 1: illegal character
2>
```

Figure 2.2: Unicode characters in allowed and disallowed context

1.2.5 Unicode file names

Most modern operating systems support Unicode file names in some way or another. There are several different ways to do this and Erlang by default treats the different approaches differently:

Mandatory Unicode file naming

> Windows and, for most common uses, MacOSX enforces Unicode support for file names. All files created in the filesystem has names that can consistently be interpreted. In MacOSX, all file names are retrieved in UTF-8 encoding, while Windows have selected an approach where each system call handling file names has a special Unicode aware variant, giving much the same effect. There are no file names on these systems that are not Unicode file names, why the default behavior of the Erlang VM is to work in "Unicode file name translation mode", meaning that a file name can be given as a Unicode list and that will be automatically translated to the proper name encoding for the underlying operating and file system.
>
> Doing i.e. a `file:list_dir/1` on one of these systems may return Unicode lists with codepoints beyond 255, depending on the content of the actual filesystem.
>
> As the feature is fairly new, you may still stumble upon non core applications that cannot handle being provided with file names containing characters with codepoints larger than 255, but the core Erlang system should have no problems with Unicode file names.

Transparent file naming

Most Unix operating systems have adopted a simpler approach, namely that Unicode file naming is not enforced, but by convention. Those systems usually use UTF-8 encoding for Unicode file names, but do not enforce it. On such a system, a file name containing characters having codepoints between 128 and 255 may be named either as plain ISO-latin-1 or using UTF-8 encoding. As no consistency is enforced, the Erlang VM can do no consistent translation of all file names. If the VM would automatically select encoding based on heuristics, one could get unexpected behavior on these systems, therefore file names not being encoded in UTF-8 are returned as "raw file names" if Unicode file naming support is turned on.

A raw file name is not a list, but a binary. Many non core applications still does not handle file names given as binaries, why such raw names are avoided by default. This means that systems having implemented Unicode file naming through transparent file systems and an UTF-8 convention, does not by default have Unicode file naming turned on. Explicitly turning Unicode file name handling on for these types of systems is considered experimental.

The Unicode file naming support was introduced with OTP release R14B01. A VM operating in Unicode file mode can work with files having names in any language or character set (as long as it's supported by the underlying OS and file system). The Unicode character list is used to denote file or directory names and if the file system content is listed, you will also be able to get Unicode lists as return value. The support lies in the kernel and stdlib modules, why most applications (that does not explicitly require the file names to be in the ISO-latin-1 range) will benefit from the Unicode support without change.

On Operating systems with mandatory Unicode file names, this means that you more easily conform to the file names of other (non Erlang) applications, and you can also process file names that, at least on Windows, where completely inaccessible (due to having names that could not be represented in ISO-latin-1). Also you will avoid creating incomprehensible file names on MacOSX as the vfs layer of the OS will accept all your file names as UTF-8 and will not rewrite them.

For most systems, turning on Unicode file name translation is no problem even if it uses transparent file naming. Very few systems have mixed file name encodings. A consistent UTF-8 named system will work perfectly in Unicode file name mode. It is still however considered experimental in R14B01. Unicode file name translation is turned on with the `+fnu` switch to the `erl` program. If the VM is started in Unicode file name translation mode, `file:native_name_encoding/0` will return the atom `utf8`.

In Unicode file name mode, file names given to the BIF `open_port/2` with the option `{spawn_executable,...}` are also interpreted as Unicode. So is the parameter list given in the `args` option available when using `spawn_executable`. The UTF-8 translation of arguments can be avoided using binaries, see the discussion about raw file names below.

It is worth noting that the file `encoding` options given when opening a file has nothing to do with the file *name* encoding convention. You can very well open files containing UTF-8 but having file names in ISO-latin-1 or vice versa.

> **Note:**
> Erlang drivers and NIF shared objects still can not be named with names containing codepoints beyond 127. This is a known limitation to be removed in a future release. Erlang modules however can, but it is definitely not a good idea and is still considered experimental.

Notes about raw file names and automatic file name conversion

Raw file names is introduced together with Unicode file name support in erts-5.8.2 (OTP R14B01). The reason the "raw file names" is introduced in the system is to be able to consistently represent file names given in different encodings on the same system. Having the VM automatically translate a file name that is not in UTF-8 when to a list of Unicode characters might seem practical, but this would open up for both duplicate file names and other inconsistent behavior. Consider a directory containing a file named "björn" in ISO-latin-1, while the Erlang VM is operating in Unicode file name mode (and therefore expecting UTF-8 file naming). The ISO-latin-1 name is not valid UTF-8 and one could

be tempted to think that automatic conversion in for example `file:list_dir/1` is a good idea. But what would happen if we later tried to open the file and has the name as a Unicode list (magically converted from the ISO-latin-1 file name)? The VM will convert the file name given to UTF-8, as this is the encoding expected. Effectively this means trying to open the file named <<"björn"/utf8>>. This file does not exist, and even if it existed it would not be the same file as the one that was listed. We could even create two files named "björn", one named in the UTF-8 encoding and one not. If `file:list_dir/1` would automatically convert the ISO-latin-1 file name to a list, we would get two identical file names as the result. To avoid this, we need to differentiate between file names being properly encoded according to the Unicode file naming convention (i.e. UTF-8) and file names being invalid under the encoding. This is done by representing invalid encoding as "raw" file names, i.e. as binaries.

The core system of Erlang (kernel and stdlib) accept raw file names except for loadable drivers and executables invoked using `open_port({spawn, ...} ...)`. `open_port({spawn_executable, ...} ...)` however does accept them. As mentioned earlier, the arguments given in the option list to `open_port({spawn_executable, ...} ...)` undergo the same conversion as the file names, meaning that the executable will be provided with arguments in UTF-8 as well. This translation is avoided consistently with how the file names are treated, by giving the argument as a binary.

To force Unicode file name translation mode on systems where this is not the default is considered experimental in OTP R14B01 due to the raw file names possibly being a new experience to the programmer and that the non core applications of OTP are not tested for compliance with raw file names yet. Unicode file name translation is expected to be default in future releases.

If working with raw file names, one can still conform to the encoding convention of the Erlang VM by using the `file:native_name_encoding/0` function, which returns either the atom `latin1` or the atom `utf8` depending on the file name translation mode. On Linux, an VM started without explicitly stating the file name translation mode will default to `latin1` as the native file name encoding, why file names on the disk encoded as UTF-8 will be returned as a list of the names interpreted as ISO-latin-1. The "UTF-8 list" is not a practical type for displaying or operating on in Erlang, but it is backward compatible and usable in all functions requiring a file name. On Windows and MacOSX, the default behavior is that of file name translation, why the `file:native_name_encoding/0` by default returns `utf8` on those systems (the fact that Windows actually does not use UTF-8 on the file system level can safely be ignored by the Erlang programmer). The default behavior can be changed using the `+fnu` or `+fnl` options to the VM, see the `erl` command manual page.

Even if you are operating without Unicode file naming translation automatically done by the VM, you can access and create files with names in UTF-8 encoding by using raw file names encoded as UTF-8. Enforcing the UTF-8 encoding regardless of the mode the Erlang VM is started in might, in some circumstances be a good idea, as the convention of using UTF-8 file names is spreading.

Notes about MacOSX

MacOSXs vfs layer enforces UTF-8 file names in a quite aggressive way. Older versions did this by simply refusing to create non UTF-8 conforming file names, while newer versions replace offending bytes with the sequence "%HH", where HH is the original character in hexadecimal notation. As Unicode translation is enabled by default on MacOSX, the only way to come up against this is to either start the VM with the `+fnl` flag or to use a raw file name in `latin1` encoding. In that case, the file can not be opened with the same name as the one used to create this. The problem is by design in newer versions of MacOSX.

MacOSX also reorganizes the names of files so that the representation of accents etc is denormalized, i.e. the character ö is represented as the codepoints [111,776], where 111 is the character `o` and 776 is a special accent character. This type of denormalized Unicode is otherwise very seldom used and Erlang normalizes those file names on retrieval, so that denormalized file names is not passed up to the Erlang application. In Erlang the file name "björn" is retrieved as [98,106,246,114,110], not as [98,106,117,776,114,110], even though the file system might think differently.

1.2.6 Unicode-aware modules

Most of the modules in Erlang/OTP are of course Unicode-unaware in the sense that they have no notion of Unicode and really shouldn't have. Typically they handle non-textual or byte-oriented data (like `gen_tcp` etc).

Modules that actually handle textual data (like `io_lib`, `string` etc) are sometimes subject to conversion or extension to be able to handle Unicode characters.

Fortunately, most textual data has been stored in lists and range checking has been sparse, why modules like `string` works well for Unicode lists with little need for conversion or extension.

Some modules are however changed to be explicitly Unicode-aware. These modules include:

`unicode`

The module *unicode* is obviously Unicode-aware. It contains functions for conversion between different Unicode formats as well as some utilities for identifying byte order marks. Few programs handling Unicode data will survive without this module.

`io`

The *io* module has been extended along with the actual I/O-protocol to handle Unicode data. This means that several functions require binaries to be in UTF-8 and there are modifiers to formatting control sequences to allow for outputting of Unicode strings.

`file`, `group` and `user`

I/O-servers throughout the system are able both to handle Unicode data and has options for converting data upon actual output or input to/from the device. As shown earlier, the *shell* has support for Unicode terminals and the *file* module allows for translation to and from various Unicode formats on disk.

The actual reading and writing of files with Unicode data is however not best done with the `file` module as its interface is byte oriented. A file opened with a Unicode encoding (like UTF-8), is then best read or written using the *io* module.

`re`

The *re* module allows for matching Unicode strings as a special option. As the library is actually centered on matching in binaries, the Unicode support is UTF-8-centered.

`wx`

The *wx* graphical library has extensive support for Unicode text

The module *string* works perfect for Unicode strings as well as for ISO-latin-1 strings with the exception of the language-dependent *to_upper* and *to_lower* functions, which are only correct for the ISO-latin-1 character set. Actually they can never function correctly for Unicode characters in their current form, there are language and locale issues as well as multi-character mappings to consider when conversion text between cases. Converting case in an international environment is a big subject not yet addressed in OTP.

1.2.7 Unicode recipes

When starting with Unicode, one often stumbles over some common issues. I try to outline some methods of dealing with Unicode data in this section.

Byte order marks

A common method of identifying encoding in text-files is to put a byte order mark (BOM) first in the file. The BOM is the codepoint 16#FEFF encoded in the same way as the rest of the file. If such a file is to be read, the first few bytes (depending on encoding) is not part of the actual text. This code outlines how to open a file which is believed to have a BOM and set the files encoding and position for further sequential reading (preferably using the *io* module). Note that error handling is omitted from the code:

```
open_bom_file_for_reading(File) ->
    {ok,F} = file:open(File,[read,binary]),
    {ok,Bin} = file:read(F,4),
    {Type,Bytes} = unicode:bom_to_encoding(Bin),
    file:position(F,Bytes),
    io:setopts(F,[{encoding,Type}]),
    {ok,F}.
```

The `unicode:bom_to_encoding/1` function identifies the encoding from a binary of at least four bytes. It returns, along with an term suitable for setting the encoding of the file, the actual length of the BOM, so that the file position can be set accordingly. Note that `file:position` always works on byte-offsets, so that the actual byte-length of the BOM is needed.

To open a file for writing and putting the BOM first is even simpler:

```
open_bom_file_for_writing(File,Encoding) ->
    {ok,F} = file:open(File,[write,binary]),
    ok = file:write(File,unicode:encoding_to_bom(Encoding)),
    io:setopts(F,[{encoding,Encoding}]),
    {ok,F}.
```

In both cases the file is then best processed using the `io` module, as the functions in `io` can handle codepoints beyond the ISO-latin-1 range.

Formatted input and output

When reading and writing to Unicode-aware entities, like the User or a file opened for Unicode translation, you will probably want to format text strings using the functions in *io* or *io_lib*. For backward compatibility reasons, these functions don't accept just any list as a string, but require e special "translation modifier" when working with Unicode texts. The modifier is "t". When applied to the "s" control character in a formatting string, it accepts all Unicode codepoints and expect binaries to be in UTF-8:

```
1> io:format("~ts~n",[<<"åäö"/utf8>>]).
åäö
ok
2> io:format("~s~n",[<<"åäö"/utf8>>]).
Ã¥Ã¤Ã¶
ok
```

Obviously the second `io:format` gives undesired output because the UTF-8 binary is not in latin1. Because ISO-latin-1 is still the defined character set of Erlang, the non prefixed "s" control character expects ISO-latin-1 in binaries as well as lists.

As long as the data is always lists, the "t" modifier can be used for any string, but when binary data is involved, care must be taken to make the tight choice of formatting characters.

The function `format` in `io_lib` behaves similarly. This function is defined to return a deep list of characters and the output could easily be converted to binary data for outputting on a device of any kind by a simple `erlang:list_to_binary`. When the translation modifier is used, the list can however contain characters that cannot be stored in one byte. The call to `erlang:list_to_binary` will in that case fail. However, if the io_server you want to communicate with is Unicode-aware, the list returned can still be used directly:

```
Eshell V5.7  (abort with ^G)
1> io_lib:format("~ts~n",["θνιψοδε"]).
[[952,957,953,968,959,948,949],"\n"]
2> io:put_chars(io_lib:format("~ts~n",["θνιψοδε"])).
θνιψοδε
ok
```

Figure 2.3: io_lib:format with Unicode translation

The Unicode string is returned as a Unicode list, why the return value of `io_lib:format` no longer qualifies as a regular Erlang string (the function *io_lib:deep_char_list* will, as an example, return `false`). The Unicode list is however valid input to the *io:put_chars* function, so data can be output on any Unicode capable device anyway. If the device is a terminal, characters will be output in the \x{H ...} format if encoding is `latin1` otherwise in UTF-8 (for the non-interactive terminal - "oldshell" or "noshell") or whatever is suitable to show the character properly (for an interactive terminal - the regular shell). The bottom line is that you can always send Unicode data to the `standard_io` device. Files will however only accept Unicode codepoints beyond ISO-latin-1 if `encoding` is set to something else than `latin1`.

Heuristic identification of UTF-8

While it's strongly encouraged that the actual encoding of characters in binary data is known prior to processing, that is not always possible. On a typical Linux® system, there is a mix of UTF-8 and ISO-latin-1 text files and there are seldom any BOM's in the files to identify them.

UTF-8 is designed in such a way that ISO-latin-1 characters with numbers beyond the 7-bit ASCII range are seldom considered valid when decoded as UTF-8. Therefore one can usually use heuristics to determine if a file is in UTF-8 or if it is encoded in ISO-latin-1 (one byte per character) encoding. The `unicode` module can be used to determine if data can be interpreted as UTF-8:

```
heuristic_encoding_bin(Bin) when is_binary(Bin) ->
    case unicode:characters_to_binary(Bin,utf8,utf8) of
 Bin ->
     utf8;
 _ ->
     latin1
    end.
```

If one does not have a complete binary of the file content, one could instead chunk through the file and check part by part. The return-tuple `{incomplete,Decoded,Rest}` from `unicode:characters_to_binary/{1,2,3}` comes in handy. The incomplete rest from one chunk of data read from the file is prepended to the next chunk and we therefore circumvent the problem of character boundaries when reading chunks of bytes in UTF-8 encoding:

```
heuristic_encoding_file(FileName) ->
    {ok,F} = file:open(FileName,[read,binary]),
    loop_through_file(F,<<>>,file:read(F,1024)).

loop_through_file(_,<<>>,eof) ->
    utf8;
loop_through_file(_,_,eof) ->
    latin1;
loop_through_file(F,Acc,{ok,Bin}) when is_binary(Bin) ->
    case unicode:characters_to_binary([Acc,Bin]) of
 {error,_,_} ->
     latin1;
 {incomplete,_,Rest} ->
     loop_through_file(F,Rest,file:read(F,1024));
 Res when is_binary(Res) ->
```

```
     loop_through_file(F,<<>>,file:read(F,1024))
    end.
```

Another option is to try to read the whole file in utf8 encoding and see if it fails. Here we need to read the file using `io:get_chars/3`, as we have to succeed in reading characters with a codepoint over 255:

```
heuristic_encoding_file2(FileName) ->
    {ok,F} = file:open(FileName,[read,binary,{encoding,utf8}]),
    loop_through_file2(F,io:get_chars(F,'',1024)).

loop_through_file2(_,eof) ->
    utf8;
loop_through_file2(_,{error,_Err}) ->
    latin1;
loop_through_file2(F,Bin) when is_binary(Bin) ->
    loop_through_file2(F,io:get_chars(F,'',1024)).
```

2 Reference Manual

The Standard Erlang Libraries application, *STDLIB*, contains modules for manipulating lists, strings and files etc.

性能评测

数据结构	写1w次时间(us)	读1w次时间(us)
array	5398	2808
dict	19081	3393
sets	15267	3009
gb_trees	42864	5061
gb_sets	46527	4902
orddict	2812805	1248887
ordsets	1699518	964396

Erlang 数据结构模块

- array - 动态数组
- dict - 动态散列表实现的字典
- sets - 动态散列表实现的集合
- queue - 双向队列
- gb_trees - 平衡二叉查找树，可作为有序字典使用
- gb_sets - 平衡二叉查找树实现的有序集合
- orddict - 列表实现的有序字典
- ordsets - 列表实现的有序集合

queue头部入队1w次时间(us)	queue头部查看1w次时间(us)	queue头部出队1w次时间(us)
940	906	1395
queue尾部入队1w次时间(us)	queue尾部查看1w次时间(us)	queue尾部出队1w次时间(us)
1036	809	1169

总结

- 要求遍历的顺序性时，尽量避免使用效率低下的 orddict 和 ordsets，用 gb_trees 和 gb_sets 代替它们；
- 不要求遍历的顺序性时，尽量使用 dict 和 sets，因它们比 gb_trees 和 gb_sets 的效率高；
- 平均来看 queue 的入队出队操作基本上同列表头部操作开销一样，效率很高，只需要双向队列访问模式的地方推荐使用；

在数据量较小的时候可以使用proplists和orddict,orddict的在复杂性和效率之间的最佳值是75.
There are basically two key-value structures/modules to deal with larger amounts of data: dicts and gb_trees.

STDLIB

Application

The STDLIB is mandatory in the sense that the minimal system based on Erlang/OTP consists of Kernel and STDLIB. The STDLIB application contains no services.

Configuration

The following configuration parameters are defined for the STDLIB application. See `app(4)` for more information about configuration parameters.

`shell_esc = icl | abort`

This parameter can be used to alter the behaviour of the Erlang shell when ^G is pressed.

`restricted_shell = module()`

This parameter can be used to run the Erlang shell in restricted mode.

`shell_catch_exception = bool()`

This parameter can be used to set the exception handling of the Erlang shell's evaluator process.

`shell_history_length = integer() >= 0`

This parameter can be used to determine how many commands are saved by the Erlang shell.

`shell_prompt_func = {Mod, Func} | default`

where

- `Mod = atom()`
- `Func = atom()`

This parameter can be used to set a customized Erlang shell prompt function.

`shell_saved_results = integer() >= 0`

This parameter can be used to determine how many results are saved by the Erlang shell.

See Also

app(4), *application(3)*, *shell(3)*,

array

可扩展数组;可固定大小,可按需自动增长.如果没有明确赋值会使用默认值undefined.
要区分是否赋值过可以使用其它的值,从不自动收缩,如果索引i的位置被赋值了,
中间的[0,i]都会处于可访问的状态,除非显示调用resize/2

Erlang module

Functional, extendible arrays. Arrays can have fixed size, or can grow automatically as needed. A default value is used for entries that have not been explicitly set.

Arrays uses *zero* based indexing. This is a deliberate design choice and differs from other erlang datastructures, e.g. tuples. 注意:索引计数从0开始,这是一个"深思熟虑"的设计决策

Unless specified by the user when the array is created, the default value is the atom `undefined`. There is no difference between an unset entry and an entry which has been explicitly set to the same value as the default one (cf. *reset/2*). If you need to differentiate between unset and set entries, you must make sure that the default value cannot be confused with the values of set entries.

The array never shrinks automatically; if an index `I` has been used successfully to set an entry, all indices in the range [0,`I`] will stay accessible unless the array size is explicitly changed by calling *resize/2*.

Examples: 不自动收缩

```
%% Create a fixed-size array with entries 0-9 set to 'undefined'
A0 = array:new(10).
10 = array:size(A0).

%% Create an extendible array and set entry 17 to 'true',
%% causing the array to grow automatically
A1 = array:set(17, true, array:new()).
18 = array:size(A1).

%% Read back a stored value
true = array:get(17, A1).

%% Accessing an unset entry returns the default value
undefined = array:get(3, A1).

%% Accessing an entry beyond the last set entry also returns the
%% default value, if the array does not have fixed size
undefined = array:get(18, A1).

%% "sparse" functions ignore default-valued entries
A2 = array:set(4, false, A1).
[{4, false}, {17, true}] = array:sparse_to_orddict(A2).

%% An extendible array can be made fixed-size later
A3 = array:fix(A2).

%% A fixed-size array does not grow automatically and does not
%% allow accesses beyond the last set entry
{'EXIT',{badarg,_}} = (catch array:set(18, true, A3)).
{'EXIT',{badarg,_}} = (catch array:get(18, A3)).
```

DATA TYPES

`array()`

A functional, extendible array. The representation is not documented and is subject to change without notice. Note that arrays cannot be directly compared for equality.

数组不可以直接比较大小

Exports

```
default(#array{default = D}) -> D;
default(_) -> erlang:error(badarg).
```

`default(Array::array()) -> term()`

Get the value used for uninitialized entries.

See also: new/2.

```
fix(#array{}=A) ->
  A#array{max = 0}.
```

`fix(Array::array()) -> array()`

Fix the size of the array. This prevents it from growing automatically upon insertion; see also *set/3*.

See also: relax/1.

`foldl(Function, InitialAcc::term(), Array::array()) -> term()`

Types:

Function = (Index::integer(), Value::term(), Acc::term()) -> term()

Fold the elements of the array using the given function and initial accumulator value. The elements are visited in order from the lowest index to the highest. If `Function` is not a function, the call fails with reason `badarg`.

See also: foldr/3, map/2, sparse_foldl/3.

`foldr(Function, InitialAcc::term(), Array::array()) -> term()`

Types:

Function = (Index::integer(), Value::term(), Acc::term()) -> term()

Fold the elements of the array right-to-left using the given function and initial accumulator value. The elements are visited in order from the highest index to the lowest. If `Function` is not a function, the call fails with reason `badarg`.

See also: foldl/3, map/2.

`from_list(List::list()) -> array()`

Equivalent to *from_list(List, undefined)*.

`from_list(List::list(), Default::term()) -> array()`

Convert a list to an extendible array. `Default` is used as the value for uninitialized entries of the array. If `List` is not a proper list, the call fails with reason `badarg`.

See also: new/2, to_list/1.

`from_orddict(Orddict::list()) -> array()`

Equivalent to *from_orddict(Orddict, undefined)*.

`from_orddict(List::list(), Default::term()) -> array()`

Convert an ordered list of pairs `{Index, Value}` to a corresponding extendible array. `Default` is used as the value for uninitialized entries of the array. If `List` is not a proper, ordered list of pairs whose first elements are nonnegative integers, the call fails with reason `badarg`.

See also: new/2, to_orddict/1.

get(I::integer(), Array::array()) -> term()

Get the value of entry `I`. If `I` is not a nonnegative integer, or if the array has fixed size and `I` is larger than the maximum index, the call fails with reason `badarg`.

If the array does not have fixed size, this function will return the default value for any index `I` greater than `size(Array)-1`.

See also: set/3.

is_array(X::term()) -> bool()

Returns `true` if `X` appears to be an array, otherwise `false`. Note that the check is only shallow; there is no guarantee that `X` is a well-formed array representation even if this function returns `true`.

is_fix(Array::array()) -> bool()

Check if the array has fixed size. Returns `true` if the array is fixed, otherwise `false`.

See also: fix/1.

map(Function, Array::array()) -> array()

Types:

Function = (Index::integer(), Value::term()) -> term()

Map the given function onto each element of the array. The elements are visited in order from the lowest index to the highest. If `Function` is not a function, the call fails with reason `badarg`.

See also: foldl/3, foldr/3, sparse_map/2.

new() -> array()

Create a new, extendible array with initial size zero.

See also: new/1, new/2.

new(Options::term()) -> array()

Create a new array according to the given options. By default, the array is extendible and has initial size zero. Array indices start at 0.

`Options` is a single term or a list of terms, selected from the following:

`N::integer()` or `{size, N::integer()}`

Specifies the initial size of the array; this also implies `{fixed, true}`. If `N` is not a nonnegative integer, the call fails with reason `badarg`.

`fixed` or `{fixed, true}`

Creates a fixed-size array; see also *fix/1*.

`{fixed, false}`

Creates an extendible (non fixed-size) array.

`{default, Value}`

Sets the default value for the array to `Value`.

Options are processed in the order they occur in the list, i.e., later options have higher precedence.

The default value is used as the value of uninitialized entries, and cannot be changed once the array has been created.

Examples:

```
array:new(100)
```

creates a fixed-size array of size 100.

```
array:new({default,0})
```

creates an empty, extendible array whose default value is 0.

```
array:new([{size,10},{fixed,false},{default,-1}])
```

creates an extendible array with initial size 10 whose default value is -1.

See also: fix/1, *from_list/2*, *get/2*, *new/0*, *new/2*, *set/3*.

new(Size::integer(), Options::term()) -> array()

Create a new array according to the given size and options. If `Size` is not a nonnegative integer, the call fails with reason `badarg`. By default, the array has fixed size. Note that any size specifications in `Options` will override the `Size` parameter.

If `Options` is a list, this is simply equivalent to `new([{size, Size} | Options]`, otherwise it is equivalent to `new([{size, Size} | [Options]]`. However, using this function directly is more efficient.

Example:

```
array:new(100, {default,0})
```

creates a fixed-size array of size 100, whose default value is 0.

See also: new/1.

relax(Array::array()) -> array()

Make the array resizable. (Reverses the effects of *fix/1*.)

See also: fix/1.

reset(I::integer(), Array::array()) -> array()

Reset entry `I` to the default value for the array. If the value of entry `I` is the default value the array will be returned unchanged. Reset will never change size of the array. Shrinking can be done explicitly by calling *resize/2*.

If `I` is not a nonnegative integer, or if the array has fixed size and `I` is larger than the maximum index, the call fails with reason `badarg`; cf. *set/3*

See also: new/2, *set/3*.

resize(Array::array()) -> array()

Change the size of the array to that reported by *sparse_size/1*. If the given array has fixed size, the resulting array will also have fixed size.

See also: resize/2, *sparse_size/1*.

resize(Size::integer(), Array::array()) -> array()

Change the size of the array. If `Size` is not a nonnegative integer, the call fails with reason `badarg`. If the given array has fixed size, the resulting array will also have fixed size.

set(I::integer(), Value::term(), Array::array()) -> array()

Set entry `I` of the array to `Value`. If `I` is not a nonnegative integer, or if the array has fixed size and `I` is larger than the maximum index, the call fails with reason `badarg`.

If the array does not have fixed size, and `I` is greater than `size(Array)-1`, the array will grow to size `I+1`.

See also: get/2, reset/2.

size(Array::array()) -> integer()

Get the number of entries in the array. Entries are numbered from 0 to `size(Array)-1`; hence, this is also the index of the first entry that is guaranteed to not have been previously set.

See also: set/3, sparse_size/1.

sparse_foldl(Function, InitialAcc::term(), Array::array()) -> term()

Types:

Function = (Index::integer(), Value::term(), Acc::term()) -> term()

Fold the elements of the array using the given function and initial accumulator value, skipping default-valued entries. The elements are visited in order from the lowest index to the highest. If `Function` is not a function, the call fails with reason `badarg`.

See also: foldl/3, sparse_foldr/3.

sparse_foldr(Function, InitialAcc::term(), Array::array()) -> term()

Types:

Function = (Index::integer(), Value::term(), Acc::term()) -> term()

Fold the elements of the array right-to-left using the given function and initial accumulator value, skipping default-valued entries. The elements are visited in order from the highest index to the lowest. If `Function` is not a function, the call fails with reason `badarg`.

See also: foldr/3, sparse_foldl/3.

sparse_map(Function, Array::array()) -> array()

Types:

Function = (Index::integer(), Value::term()) -> term()

Map the given function onto each element of the array, skipping default-valued entries. The elements are visited in order from the lowest index to the highest. If `Function` is not a function, the call fails with reason `badarg`.

See also: map/2.

sparse_size(A::array()) -> integer()

Get the number of entries in the array up until the last non-default valued entry. In other words, returns `I+1` if `I` is the last non-default valued entry in the array, or zero if no such entry exists.

See also: resize/1, size/1.

sparse_to_list(Array::array()) -> list()

Converts the array to a list, skipping default-valued entries.

See also: to_list/1.

sparse_to_orddict(Array::array()) -> [{Index::integer(), Value::term()}]

Convert the array to an ordered list of pairs `{Index, Value}`, skipping default-valued entries.

See also: to_orddict/1.

to_list(Array::array()) -> list()

Converts the array to a list.

See also: from_list/2, sparse_to_list/1.

to_orddict(Array::array()) -> [{Index::integer(), Value::term()}]

Convert the array to an ordered list of pairs `{Index, Value}`.

See also: from_orddict/2, sparse_to_orddict/1.

array

[1] 考虑到函数式语言数据不可变的特性，array 使用基数为 10 的 tuple tree 形式实现，
将元素更改发生时需要复制的数据限制在很小的范围内，以便提高效率。

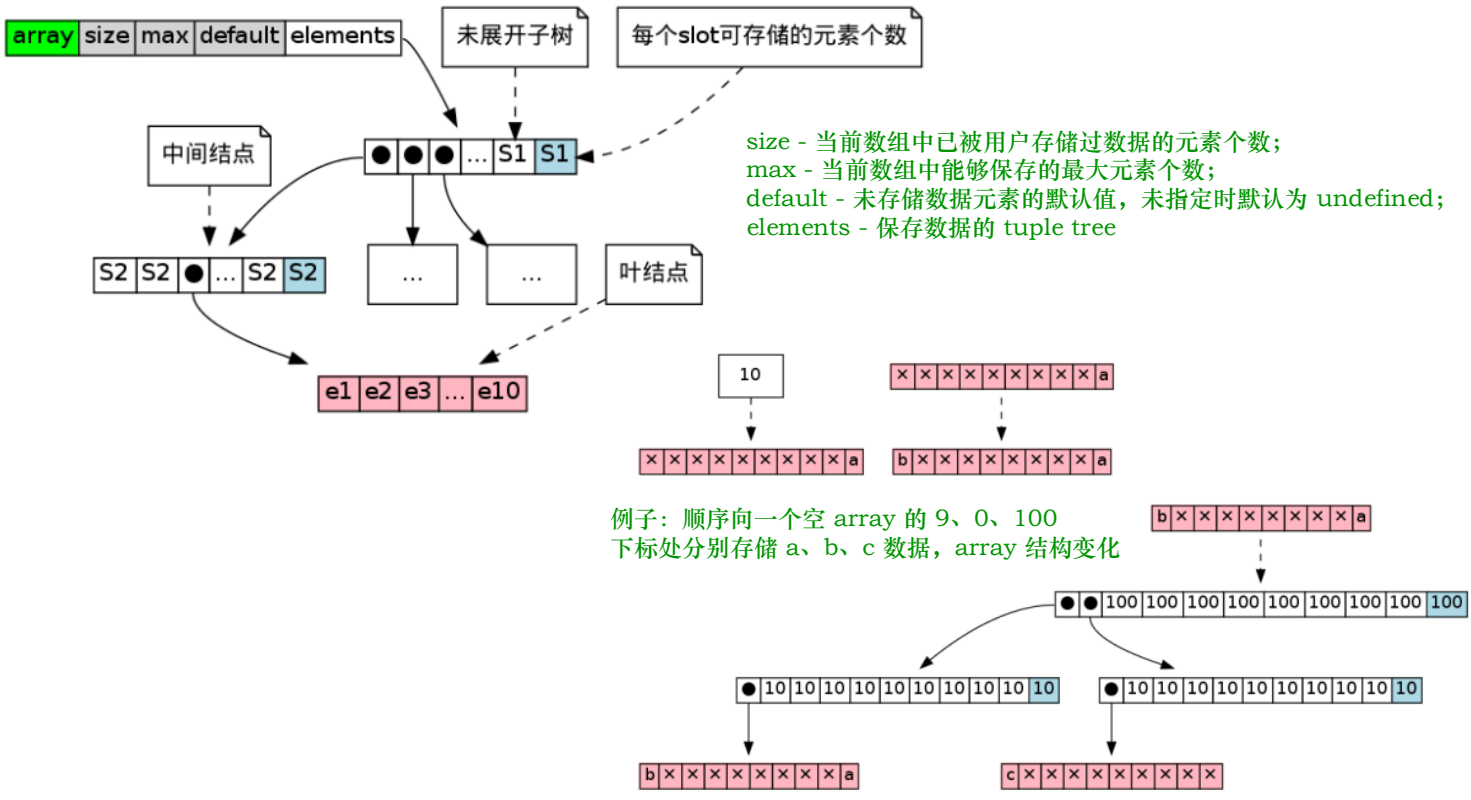
[2] 为了平衡存储空间和动态性，tuple tree 是惰性展开的，仅在某个下标被使用时才会创建出对应的 tuple 结构进行存储。

[3] 到 R14B02 版本为止的实现里，tuple tree 都是只增大不减小的，对 array 进行 resize/2 操作时仅仅是更新了其 size 属性而已；

[4] array 的下标是从 0 开始的，同 tuple、list 等结构起始下标不同！

[5] array 有两种遍历方式：普通遍历和稀疏遍历，区别在于稀疏遍历会跳过所有未定义（取值为默认值）的元素，而普通遍历则不会跳过。

- array 内部结构如下图所示：



base64

Erlang module

Implements base 64 encode and decode, see RFC2045.

Exports

`encode(Data) -> Base64`

`encode_to_string(Data) -> Base64String`

Types:

Data = string() | binary()

Base64 = binary()

Base64String = string()

Encodes a plain ASCII string into base64. The result will be 33% larger than the data.

`decode(Base64) -> Data`

`decode_to_string(Base64) -> DataString`

`mime_decode(Base64) -> Data`

`mime_decode_to_string(Base64) -> DataString`

Types:

Base64 = string() | binary()

Data = binary()

DataString = string()

Decodes a base64 encoded string to plain ASCII. See RFC4648. `mime_decode/1` and `mime_decode_to_string/1` strips away illegal characters, while `decode/1` and `decode_to_string/1` only strips away whitespace characters.

beam_lib

Erlang module

`beam_lib` provides an interface to files created by the BEAM compiler ("BEAM files"). The format used, a variant of "EA IFF 1985" Standard for Interchange Format Files, divides data into chunks.

Chunk data can be returned as binaries or as compound terms. Compound terms are returned when chunks are referenced by names (atoms) rather than identifiers (strings). The names recognized and the corresponding identifiers are:

- `abstract_code ("Abst")`
- `attributes ("Attr")`
- `compile_info ("CInf")`
- `exports ("ExpT")`
- `labeled_exports ("ExpT")`
- `imports ("ImpT")`
- `indexed_imports ("ImpT")`
- `locals ("LocT")`
- `labeled_locals ("LocT")`
- `atoms ("Atom")`

Debug Information/Abstract Code

The option `debug_info` can be given to the compiler (see *compile(3)*) in order to have debug information in the form of abstract code (see *The Abstract Format* in ERTS User's Guide) stored in the `abstract_code` chunk. Tools such as Debugger and Xref require the debug information to be included.

> ### Warning:
>
> Source code can be reconstructed from the debug information. Use encrypted debug information (see below) to prevent this.

The debug information can also be removed from BEAM files using *strip/1*, *strip_files/1* and/or *strip_release/1*.

Reconstructing source code

Here is an example of how to reconstruct source code from the debug information in a BEAM file `Beam`:

```
{ok,{_,[{abstract_code,{_,AC}}]}} = beam_lib:chunks(Beam,[abstract_code]).
io:fwrite("~s~n", [erl_prettypr:format(erl_syntax:form_list(AC))]).
```

Encrypted debug information

The debug information can be encrypted in order to keep the source code secret, but still being able to use tools such as Xref or Debugger.

To use encrypted debug information, a key must be provided to the compiler and `beam_lib`. The key is given as a string and it is recommended that it contains at least 32 characters and that both upper and lower case letters as well as digits and special characters are used.

The default type -- and currently the only type -- of crypto algorithm is `des3_cbc`, three rounds of DES. The key string will be scrambled using `erlang:md5/1` to generate the actual keys used for `des3_cbc`.

> ### Note:
>
> As far as we know by the time of writing, it is infeasible to break `des3_cbc` encryption without any knowledge of the key. Therefore, as long as the key is kept safe and is unguessable, the encrypted debug information *should* be safe from intruders.

There are two ways to provide the key:

- Use the compiler option `{debug_info,Key}`, see *compile(3)*, and the function *crypto_key_fun/1* to register a fun which returns the key whenever `beam_lib` needs to decrypt the debug information.

  If no such fun is registered, `beam_lib` will instead search for a `.erlang.crypt` file, see below.
- Store the key in a text file named `.erlang.crypt`.

  In this case, the compiler option `encrypt_debug_info` can be used, see *compile(3)*.

.erlang.crypt

`beam_lib` searches for `.erlang.crypt` in the current directory and then the home directory for the current user. If the file is found and contains a key, `beam_lib` will implicitly create a crypto key fun and register it.

The `.erlang.crypt` file should contain a single list of tuples:

```
    {debug_info, Mode, Module, Key}
```

`Mode` is the type of crypto algorithm; currently, the only allowed value thus is `des3_cbc`. `Module` is either an atom, in which case `Key` will only be used for the module `Module`, or `[]`, in which case `Key` will be used for all modules. `Key` is the non-empty key string.

The `Key` in the first tuple where both `Mode` and `Module` matches will be used.

Here is an example of an `.erlang.crypt` file that returns the same key for all modules:

```
[{debug_info, des3_cbc, [], "%>7}|pc/DM6Cga*68$Mw]L#&_Gejr]G^"}].
```

And here is a slightly more complicated example of an `.erlang.crypt` which provides one key for the module `t`, and another key for all other modules:

```
[{debug_info, des3_cbc, t, "My KEY"},
 {debug_info, des3_cbc, [], "%>7}|pc/DM6Cga*68$Mw]L#&_Gejr]G^"}].
```

> ### Note:
>
> Do not use any of the keys in these examples. Use your own keys.

DATA TYPES

```
beam() -> Module | Filename | binary()
  Module = atom()
  Filename = string() | atom()
```

Each of the functions described below accept either the module name, the filename, or a binary containing the beam module.

```
chunkdata() = {ChunkId, DataB} | {ChunkName, DataT}
  ChunkId = chunkid()
  DataB = binary()
  {ChunkName, DataT} =
        {abstract_code, AbstractCode}
      | {attributes, [{Attribute, [AttributeValue]}]}
      | {compile_info, [{InfoKey, [InfoValue]}]}
      | {exports, [{Function, Arity}]}
      | {labeled_exports, [{Function, Arity, Label}]}
      | {imports, [{Module, Function, Arity}]}
      | {indexed_imports, [{Index, Module, Function, Arity}]}
      | {locals, [{Function, Arity}]}]}
      | {labeled_locals, [{Function, Arity, Label}]}]}
      | {atoms, [{integer(), atom()}]}
  AbstractCode = {AbstVersion, Forms} | no_abstract_code
    AbstVersion = atom()
  Attribute = atom()
  AttributeValue = term()
  Module = Function = atom()
  Arity = int()
  Label = int()
```

It is not checked that the forms conform to the abstract format indicated by `AbstVersion`. `no_abstract_code` means that the `"Abst"` chunk is present, but empty.

The list of attributes is sorted on `Attribute`, and each attribute name occurs once in the list. The attribute values occur in the same order as in the file. The lists of functions are also sorted.

```
chunkid() = "Abst" | "Attr" | "CInf"
            | "ExpT" | "ImpT" | "LocT"
            | "Atom"

chunkname() = abstract_code | attributes | compile_info
            | exports | labeled_exports
            | imports | indexed_imports
            | locals | labeled_locals
            | atoms

chunkref() = chunkname() | chunkid()
```

Exports

`chunks(Beam, [ChunkRef]) -> {ok, {Module, [ChunkData]}} | {error, beam_lib, Reason}`

Types:

Beam = beam()

ChunkRef = chunkref()

Module = atom()

ChunkData = chunkdata()

Reason = {unknown_chunk, Filename, atom()}

| {key_missing_or_invalid, Filename, abstract_code}

| Reason1 -- see info/1

Filename = string()

Reads chunk data for selected chunks refs. The order of the returned list of chunk data is determined by the order of the list of chunks references.

`chunks(Beam, [ChunkRef], [Option]) -> {ok, {Module, [ChunkResult]}} | {error, beam_lib, Reason}`

Types:

Beam = beam()

ChunkRef = chunkref()

Module = atom()

Option = allow_missing_chunks

ChunkResult = {chunkref(), ChunkContents} | {chunkref(), missing_chunk}

Reason = {missing_chunk, Filename, atom()}

| {key_missing_or_invalid, Filename, abstract_code}

| Reason1 -- see info/1

Filename = string()

Reads chunk data for selected chunks refs. The order of the returned list of chunk data is determined by the order of the list of chunks references.

By default, if any requested chunk is missing in `Beam`, an `error` tuple is returned. However, if the option `allow_missing_chunks` has been given, a result will be returned even if chunks are missing. In the result list, any missing chunks will be represented as `{ChunkRef,missing_chunk}`. Note, however, that if the `"Atom"` chunk if missing, that is considered a fatal error and the return value will be an `error` tuple.

`version(Beam) -> {ok, {Module, [Version]}} | {error, beam_lib, Reason}`

Types:

Beam = beam()

Module = atom()

Version = term()

Reason -- see chunks/2

Returns the module version(s). A version is defined by the module attribute `-vsn(Vsn)`. If this attribute is not specified, the version defaults to the checksum of the module. Note that if the version `Vsn` is not a list, it is made into one, that is `{ok,{Module,[Vsn]}}` is returned. If there are several `-vsn` module attributes, the result is the concatenated list of versions. Examples:

```
1> beam_lib:version(a). % -vsn(1).
{ok,{a,[1]}}
2> beam_lib:version(b). % -vsn([1]).
{ok,{b,[1]}}
```

```
3> beam_lib:version(c). % -vsn([1]). -vsn(2).
{ok,{c,[1,2]}}
4> beam_lib:version(d). % no -vsn attribute
{ok,{d,[275613208176997377698094100858909383631]}}
```

md5(Beam) -> {ok, {Module, MD5}} | {error, beam_lib, Reason}

Types:

Beam = beam()

Module = atom()

MD5 = binary()

Reason -- see chunks/2

Calculates an MD5 redundancy check for the code of the module (compilation date and other attributes are not included).

info(Beam) -> [{Item, Info}] | {error, beam_lib, Reason1}

Types:

Beam = beam()

Item, Info -- see below

Reason1 = {chunk_too_big, Filename, ChunkId, ChunkSize, FileSize}
| {invalid_beam_file, Filename, Pos}
| {invalid_chunk, Filename, ChunkId}
| {missing_chunk, Filename, ChunkId}
| {not_a_beam_file, Filename}
| {file_error, Filename, Posix}

Filename = string()

ChunkId = chunkid()

ChunkSize = FileSize = int()

Pos = int()

Posix = posix() -- see file(3)

Returns a list containing some information about a BEAM file as tuples `{Item, Info}`:

`{file, Filename} | {binary, Binary}`

The name (string) of the BEAM file, or the binary from which the information was extracted.

`{module, Module}`

The name (atom) of the module.

`{chunks, [{ChunkId, Pos, Size}]}`

For each chunk, the identifier (string) and the position and size of the chunk data, in bytes.

cmp(Beam1, Beam2) -> ok | {error, beam_lib, Reason}

Types:

Beam1 = Beam2 = beam()

Reason = {modules_different, Module1, Module2}
| {chunks_different, ChunkId}
| different_chunks

| Reason1 -- see info/1

Module1 = Module2 = atom()

ChunkId = chunkid()

Compares the contents of two BEAM files. If the module names are the same, and all chunks except for the `"CInf"` chunk (the chunk containing the compilation information which is returned by `Module:module_info(compile)`) have the same contents in both files, `ok` is returned. Otherwise an error message is returned.

`cmp_dirs(Dir1, Dir2) -> {Only1, Only2, Different} | {error, beam_lib, Reason1}`

Types:

Dir1 = Dir2 = string() | atom()

Different = [{Filename1, Filename2}]

Only1 = Only2 = [Filename]

Filename = Filename1 = Filename2 = string()

Reason1 = {not_a_directory, term()} | -- see info/1

The `cmp_dirs/2` function compares the BEAM files in two directories. Only files with extension `".beam"` are compared. BEAM files that exist in directory `Dir1` (`Dir2`) only are returned in `Only1` (`Only2`). BEAM files that exist on both directories but are considered different by `cmp/2` are returned as pairs {`Filename1`, `Filename2`} where `Filename1` (`Filename2`) exists in directory `Dir1` (`Dir2`).

`diff_dirs(Dir1, Dir2) -> ok | {error, beam_lib, Reason1}`

Types:

Dir1 = Dir2 = string() | atom()

Reason1 = {not_a_directory, term()} | -- see info/1

The `diff_dirs/2` function compares the BEAM files in two directories the way `cmp_dirs/2` does, but names of files that exist in only one directory or are different are presented on standard output.

`strip(Beam1) -> {ok, {Module, Beam2}} | {error, beam_lib, Reason1}`

Types:

Beam1 = Beam2 = beam()

Module = atom()

Reason1 -- see info/1

The `strip/1` function removes all chunks from a BEAM file except those needed by the loader. In particular, the debug information (`abstract_code` chunk) is removed.

`strip_files(Files) -> {ok, [{Module, Beam2}]} | {error, beam_lib, Reason1}`

Types:

Files = [Beam1]

Beam1 = beam()

Module = atom()

Beam2 = beam()

Reason1 -- see info/1

The `strip_files/1` function removes all chunks except those needed by the loader from BEAM files. In particular, the debug information (`abstract_code` chunk) is removed. The returned list contains one element for each given file name, in the same order as in `Files`.

`strip_release(Dir) -> {ok, [{Module, Filename]}} | {error, beam_lib, Reason1}`

Types:

Dir = string() | atom()

Module = atom()

Filename = string()

Reason1 = {not_a_directory, term()} | -- see info/1

The `strip_release/1` function removes all chunks except those needed by the loader from the BEAM files of a release. `Dir` should be the installation root directory. For example, the current OTP release can be stripped with the call `beam_lib:strip_release(code:root_dir())`.

`format_error(Reason) -> Chars`

Types:

Reason -- see other functions

Chars = [char() | Chars]

Given the error returned by any function in this module, the function `format_error` returns a descriptive string of the error in English. For file errors, the function `file:format_error(Posix)` should be called.

`crypto_key_fun(CryptoKeyFun) -> ok | {error, Reason}`

Types:

CryptoKeyFun = fun() -- see below

Reason = badfun | exists | term()

The `crypto_key_fun/1` function registers a unary fun that will be called if `beam_lib` needs to read an `abstract_code` chunk that has been encrypted. The fun is held in a process that is started by the function.

If there already is a fun registered when attempting to register a fun, `{error, exists}` is returned.

The fun must handle the following arguments:

```
CryptoKeyFun(init) -> ok | {ok, NewCryptoKeyFun} | {error, Term}
```

Called when the fun is registered, in the process that holds the fun. Here the crypto key fun can do any necessary initializations. If `{ok, NewCryptoKeyFun}` is returned then `NewCryptoKeyFun` will be registered instead of `CryptoKeyFun`. If `{error, Term}` is returned, the registration is aborted and `crypto_key_fun/1` returns `{error, Term}` as well.

```
CryptoKeyFun({debug_info, Mode, Module, Filename}) -> Key
```

Called when the key is needed for the module `Module` in the file named `Filename`. `Mode` is the type of crypto algorithm; currently, the only possible value thus is `des3_cbc`. The call should fail (raise an exception) if there is no key available.

```
CryptoKeyFun(clear) -> term()
```

Called before the fun is unregistered. Here any cleaning up can be done. The return value is not important, but is passed back to the caller of `clear_crypto_key_fun/0` as part of its return value.

`clear_crypto_key_fun() -> {ok, Result}`

Types:

Result = undefined | term()

Unregisters the crypto key fun and terminates the process holding it, started by `crypto_key_fun/1`.

The `clear_crypto_key_fun/1` either returns `{ok, undefined}` if there was no crypto key fun registered, or `{ok, Term}`, where `Term` is the return value from `CryptoKeyFun(clear)`, see `crypto_key_fun/1`.

binary

Erlang module

This module contains functions for manipulating byte-oriented binaries. Although the majority of functions could be implemented using bit-syntax, the functions in this library are highly optimized and are expected to either execute faster or consume less memory (or both) than a counterpart written in pure Erlang.

The module is implemented according to the EEP (Erlang Enhancement Proposal) 31.

> ### Note:
>
> The library handles byte-oriented data. Bitstrings that are not binaries (does not contain whole octets of bits) will result in a `badarg` exception being thrown from any of the functions in this module.

DATA TYPES

```
cp()
 - Opaque data-type representing a compiled search-pattern. Guaranteed to be a tuple()
   to allow programs to distinguish it from non precompiled search patterns.
```

```
part() = {Start,Length}
Start = int()
Length = int()
  - A representaion of a part (or range) in a binary. Start is a
    zero-based offset into a binary() and Length is the length of
    that part. As input to functions in this module, a reverse
    part specification is allowed, constructed with a negative
    Length, so that the part of the binary begins at Start +
    Length and is -Length long. This is useful for referencing the
    last N bytes of a binary as {size(Binary), -N}. The functions
    in this module always return part()'s with positive Length.
```

Exports

`at(Subject, Pos) -> int()`

Types:

Subject = binary()

Pos = int() >= 0

Returns the byte at position `Pos` (zero-based) in the binary `Subject` as an integer. If `Pos` >= `byte_size(Subject)`, a `badarg` exception is raised.

`bin_to_list(Subject) -> list()`

Types:

Subject = binary()

The same as `bin_to_list(Subject,{0,byte_size(Subject)})`.

`bin_to_list(Subject, PosLen) -> list()`

Types:

Subject = **binary()**

PosLen = **part()**

Converts `Subject` to a list of `int()`s, each representing the value of one byte. The `part()` denotes which part of the `binary()` to convert. Example:

```
1> binary:bin_to_list(<<"erlang">>,{1,3}).
"rla"
%% or [114,108,97] in list notation.
```

If `PosLen` in any way references outside the binary, a `badarg` exception is raised.

`bin_to_list(Subject, Pos, Len) -> list()`

Types:

Subject = **binary()**

Pos = **int()**

Len = **int()**

The same as `bin_to_list(Subject,{Pos,Len})`.

`compile_pattern(Pattern) -> cp()`

Types:

Pattern = **binary() | [binary()]**

Builds an internal structure representing a compilation of a search-pattern, later to be used in the *match/3*, *matches/3*, *split/3* or *replace/4* functions. The `cp()` returned is guaranteed to be a `tuple()` to allow programs to distinguish it from non pre-compiled search patterns

When a list of binaries is given, it denotes a set of alternative binaries to search for. I.e if `[<<"functional">>,<<"programming">>]` is given as `Pattern`, this means "either `<<"functional">>` or `<<"programming">>`". The pattern is a set of alternatives; when only a single binary is given, the set has only one element. The order of alternatives in a pattern is not significant.

The list of binaries used for search alternatives shall be flat and proper.

If `Pattern` is not a binary or a flat proper list of binaries with length > 0, a `badarg` exception will be raised.

`copy(Subject) -> binary()`

Types:

Subject = **binary()**

The same as `copy(Subject, 1)`.

`copy(Subject,N) -> binary()`

Types:

Subject = **binary()**

N = **int()** >= **0**

Creates a binary with the content of `Subject` duplicated `N` times.

This function will always create a new binary, even if `N = 1`. By using `copy/1` on a binary referencing a larger binary, one might free up the larger binary for garbage collection.

> Note:
>
> By deliberately copying a single binary to avoid referencing a larger binary, one might, instead of freeing up the larger binary for later garbage collection, create much more binary data than needed. Sharing binary data is usually good. Only in special cases, when small parts reference large binaries and the large binaries are no longer used in any process, deliberate copying might be a good idea.

If `N < 0`, a `badarg` exception is raised.

`decode_unsigned(Subject) -> Unsigned`

Types:

Subject = binary()

Unsigned = int() >= 0

The same as `decode_unsigned(Subject,big)`.

`decode_unsigned(Subject, Endianess) -> Unsigned`

Types:

Subject = binary()

Endianess = big | little

Unsigned = int() >= 0

Converts the binary digit representation, in big or little endian, of a positive integer in `Subject` to an Erlang `int()`.

Example:

```
1> binary:decode_unsigned(<<169,138,199>>,big).
11111111
```

`encode_unsigned(Unsigned) -> binary()`

Types:

Unsigned = int() >= 0

The same as `encode_unsigned(Unsigned,big)`.

`encode_unsigned(Unsigned,Endianess) -> binary()`

Types:

Unsigned = int() >= 0

Endianess = big | little

Converts a positive integer to the smallest possible representation in a binary digit representation, either big or little endian.

Example:

```
1> binary:encode_unsigned(11111111,big).
<<169,138,199>>
```

`first(Subject) -> int()`

Types:

Subject = binary()

Returns the first byte of the binary `Subject` as an integer. If the size of `Subject` is zero, a `badarg` exception is raised.

`last(Subject) -> int()`

Types:

Subject = binary()

Returns the last byte of the binary `Subject` as an integer. If the size of `Subject` is zero, a `badarg` exception is raised.

`list_to_bin(ByteList) -> binary()`

Types:

ByteList = iodata() (see module erlang)

Works exactly as `erlang:list_to_binary/1`, added for completeness.

`longest_common_prefix(Binaries) -> int()`

Types:

Binaries = [binary()]

Returns the length of the longest common prefix of the binaries in the list `Binaries`. Example:

```
1> binary:longest_common_prefix([<<"erlang">>,<<"ergonomy">>]).
2
2> binary:longest_common_prefix([<<"erlang">>,<<"perl">>]).
0
```

If `Binaries` is not a flat list of binaries, a `badarg` exception is raised.

`longest_common_suffix(Binaries) -> int()`

Types:

Binaries = [binary()]

Returns the length of the longest common suffix of the binaries in the list `Binaries`. Example:

```
1> binary:longest_common_suffix([<<"erlang">>,<<"fang">>]).
3
2> binary:longest_common_suffix([<<"erlang">>,<<"perl">>]).
```

```
0
```

If `Binaries` is not a flat list of binaries, a `badarg` exception is raised.

`match(Subject, Pattern) -> Found | nomatch`

Types:

Subject = binary()

Pattern = binary() | [binary()] | cp()

Found = part()

The same as `match(Subject, Pattern, [])`.

`match(Subject,Pattern,Options) -> Found | nomatch`

Types:

Subject = binary()

Pattern = binary() | [binary()] | cp()

Found = part()

Options = [Option]

Option = {scope, part()}

Searches for the first occurrence of `Pattern` in `Subject` and returns the position and length.

The function will return `{Pos,Length}` for the binary in `Pattern` starting at the lowest position in `Subject`, Example:

```
1> binary:match(<<"abcde">>, [<<"bcde">>,<<"cd">>],[]).
{1,4}
```

Even though `<<"cd">>` ends before `<<"bcde">>`, `<<"bcde">>` begins first and is therefore the first match. If two overlapping matches begin at the same position, the longest is returned.

Summary of the options:

{scope, {Start, Length}}

Only the given part is searched. Return values still have offsets from the beginning of `Subject`. A negative `Length` is allowed as described in the TYPES section of this manual.

If none of the strings in `Pattern` is found, the atom `nomatch` is returned.

For a description of `Pattern`, see *compile_pattern/1*.

If `{scope, {Start,Length}}` is given in the options such that `Start` is larger than the size of `Subject`, `Start + Length` is less than zero or `Start + Length` is larger than the size of `Subject`, a `badarg` exception is raised.

`matches(Subject, Pattern) -> Found`

Types:

Subject = binary()

Pattern = binary() | [binary()] | cp()

Found = [part()] | []

The same as `matches(Subject, Pattern, [])`.

matches(Subject,Pattern,Options) -> Found

Types:

Subject = binary()

Pattern = binary() | [binary()] | cp()

Found = [part()] | []

Options = [Option]

Option = {scope, part()}

Works like match, but the `Subject` is searched until exhausted and a list of all non-overlapping parts matching `Pattern` is returned (in order).

The first and longest match is preferred to a shorter, which is illustrated by the following example:

```
1> binary:matches(<<"abcde">>,
                  [<<"bcde">>,<<"bc">>>,<<"de">>],[]).
[{1,4}]
```

The result shows that <<bcde">> is selected instead of the shorter match <<"bc">> (which would have given raise to one more match,<<"de">>). This corresponds to the behavior of posix regular expressions (and programs like awk), but is not consistent with alternative matches in re (and Perl), where instead lexical ordering in the search pattern selects which string matches.

If none of the strings in pattern is found, an empty list is returned.

For a description of `Pattern`, see *compile_pattern/1* and for a description of available options, see *match/3*.

If `{scope, {Start,Length}}` is given in the options such that `Start` is larger than the size of `Subject`, `Start + Length` is less than zero or `Start + Length` is larger than the size of `Subject`, a `badarg` exception is raised.

part(Subject, PosLen) -> binary()

Types:

Subject = binary()

PosLen = part()

Extracts the part of the binary `Subject` described by `PosLen`.

Negative length can be used to extract bytes at the end of a binary:

```
1> Bin = <<1,2,3,4,5,6,7,8,9,10>>.
2> binary:part(Bin,{byte_size(Bin), -5)).
<<6,7,8,9,10>>
```

> Note:
>
> *part/2*and *part/3* are also available in the `erlang` module under the names `binary_part/2` and `binary_part/3`. Those BIFs are allowed in guard tests.

If `PosLen` in any way references outside the binary, a `badarg` exception is raised.

part(Subject, Pos, Len) -> binary()

Types:

Subject = binary()

Pos = int()

Len = int()

The same as `part(Subject, {Pos, Len})`.

referenced_byte_size(binary()) -> int()

If a binary references a larger binary (often described as being a sub-binary), it can be useful to get the size of the actual referenced binary. This function can be used in a program to trigger the use of `copy/1`. By copying a binary, one might dereference the original, possibly large, binary which a smaller binary is a reference to.

Example:

```
store(Binary, GBSet) ->
  NewBin =
      case binary:referenced_byte_size(Binary) of
          Large when Large > 2 * byte_size(Binary) ->
             binary:copy(Binary);
          _ ->
             Binary
      end,
  gb_sets:insert(NewBin,GBSet).
```

In this example, we chose to copy the binary content before inserting it in the `gb_set()` if it references a binary more than twice the size of the data we're going to keep. Of course different rules for when copying will apply to different programs.

Binary sharing will occur whenever binaries are taken apart, this is the fundamental reason why binaries are fast, decomposition can always be done with O(1) complexity. In rare circumstances this data sharing is however undesirable, why this function together with `copy/1` might be useful when optimizing for memory use.

Example of binary sharing:

```
1> A = binary:copy(<<1>>,100).
<<1,1,1,1,1 ...
2> byte_size(A).
100
3> binary:referenced_byte_size(A)
100
4> <<_:10/binary,B:10/binary,_/binary>> = A.
<<1,1,1,1,1 ...
5> byte_size(B).
10
6> binary:referenced_byte_size(B)
100
```

> Note:
>
> Binary data is shared among processes. If another process still references the larger binary, copying the part this process uses only consumes more memory and will not free up the larger binary for garbage collection. Use this kind of intrusive functions with extreme care, and only if a real problem is detected.

`replace(Subject,Pattern,Replacement) -> Result`

Types:

Subject = binary()

Pattern = binary() | [binary()] | cp()

Replacement = binary()

Result = binary()

The same as `replace(Subject,Pattern,Replacement,[])`.

`replace(Subject,Pattern,Replacement,Options) -> Result`

Types:

Subject = binary()

Pattern = binary() | [binary()] | cp()

Replacement = binary()

Result = binary()

Options = [Option]

Option = global | {scope, part()} | {insert_replaced, InsPos}

InsPos = OnePos | [OnePos]

OnePos = int() =< byte_size(Replacement)

Constructs a new binary by replacing the parts in `Subject` matching `Pattern` with the content of `Replacement`.

If the matching sub-part of `Subject` giving raise to the replacement is to be inserted in the result, the option `{insert_replaced, InsPos}` will insert the matching part into `Replacement` at the given position (or positions) before actually inserting `Replacement` into the `Subject`. Example:

```
1> binary:replace(<<"abcde">>,<<"b">>,<<"[]">>,[{insert_replaced,1}]).
<<"a[b]cde">>
2> binary:replace(<<"abcde">>,[<<"b">>,<<"d">>],<<"[]">>,
                  [global,{insert_replaced,1}]).
<<"a[b]c[d]e">>
3> binary:replace(<<"abcde">>,[<<"b">>,<<"d">>],<<"[]">>,
                  [global,{insert_replaced,[1,1]}]).
<<"a[bb]c[dd]e">>
4> binary:replace(<<"abcde">>,[<<"b">>,<<"d">>],<<"[-]">>,
                  [global,{insert_replaced,[1,2]}]).
<<"a[b-b]c[d-d]e">>
```

If any position given in `InsPos` is greater than the size of the replacement binary, a `badarg` exception is raised.

The options `global` and `{scope, part()}` work as for *split/3*. The return type is always a `binary()`.

For a description of `Pattern`, see *compile_pattern/1*.

split(Subject,Pattern) -> Parts

Types:

Subject = binary()

Pattern = binary() | [binary()] | cp()

Parts = [binary()]

The same as `split(Subject, Pattern, [])`.

split(Subject,Pattern,Options) -> Parts

Types:

Subject = binary()

Pattern = binary() | [binary()] | cp()

Parts = [binary()]

Options = [Option]

Option = {scope, part()} | trim | global

Splits Binary into a list of binaries based on Pattern. If the option global is not given, only the first occurrence of Pattern in Subject will give rise to a split.

The parts of Pattern actually found in Subject are not included in the result.

Example:

```
1> binary:split(<<1,255,4,0,0,0,2,3>>, [<<0,0,0>>,<<2>>],[]).
[<<1,255,4>>, <<2,3>>]
2> binary:split(<<0,1,0,0,4,255,255,9>>, [<<0,0>>, <<255,255>>],[global]).
[<<0,1>>,<<4>>,<<9>>]
```

Summary of options:

{scope, part()}

Works as in *match/3* and *matches/3*. Note that this only defines the scope of the search for matching strings, it does not cut the binary before splitting. The bytes before and after the scope will be kept in the result. See example below.

trim

Removes trailing empty parts of the result (as does trim in `re:split/3`)

global

Repeats the split until the `Subject` is exhausted. Conceptually the global option makes split work on the positions returned by *matches/3*, while it normally works on the position returned by *match/3*.

Example of the difference between a scope and taking the binary apart before splitting:

```
1> binary:split(<<"banana">>,[<<"a">>],[{scope,{2,3}}]).
[<<"ban">>,<<"na">>]
2> binary:split(binary:part(<<"banana">>,{2,3}),[<<"a">>],[]).
[<<"n">>,<<"n">>]
```

The return type is always a list of binaries that are all referencing `Subject`. This means that the data in `Subject` is not actually copied to new binaries and that `Subject` cannot be garbage collected until the results of the split are no longer referenced.

For a description of `Pattern`, see *compile_pattern/1*.

c

Erlang module

The `c` module enables users to enter the short form of some commonly used commands.

> ### Note:
> These functions are are intended for interactive use in the Erlang shell only. The module prefix may be omitted.

Exports

`bt(Pid) -> void()`

Types:

Pid = pid()

Stack backtrace for a process. Equivalent to `erlang:process_display(Pid, backtrace)`.

`c(File) -> {ok, Module} | error`

`c(File, Options) -> {ok, Module} | error`

Types:

File = name() -- see filename(3)

Options = [Opt] -- see compile:file/2

`c/1,2` compiles and then purges and loads the code for a file. `Options` defaults to []. Compilation is equivalent to:

```
compile:file(File, Options ++ [report_errors, report_warnings])
```

Note that purging the code means that any processes lingering in old code for the module are killed without warning. See `code/3` for more information.

`cd(Dir) -> void()`

Types:

Dir = name() -- see filename(3)

Changes working directory to `Dir`, which may be a relative name, and then prints the name of the new working directory.

```
2> cd("../erlang").
/home/ron/erlang
```

`flush() -> void()`

Flushes any messages sent to the shell.

help() -> void()

Displays help information: all valid shell internal commands, and commands in this module.

i() -> void()

ni() -> void()

erlang:processes().

`i/0` displays information about the system, listing information about all processes. `ni/0` does the same, but for all nodes the network.

i(X, Y, Z) -> void()

相当于process_info(pid(X,Y,Z))

Types:

X = Y = Z = int()

Displays information about a process, Equivalent to `process_info(pid(X, Y, Z))`, but location transparent.

l(Module) -> void()

Types:

Module = atom()

Purges and loads, or reloads, a module by calling `code:purge(Module)` followed by `code:load_file(Module)`.

Note that purging the code means that any processes lingering in old code for the module are killed without warning. See `code/3` for more information.

lc(Files) -> ok

Types:

Files = [File]

File = name() -- see filename(3)

Compiles a list of files by calling `compile:file(File, [report_errors, report_warnings])` for each `File` in `Files`.

ls() -> void()

Lists files in the current directory.

ls(Dir) -> void()

Types:

Dir = name() -- see filename(3)

Lists files in directory `Dir`.

m() -> void()

Displays information about the loaded modules, including the files from which they have been loaded.

m(Module) -> void()

Types:

Module = atom()

Displays information about `Module`.

`memory() -> [{Type, Size}]`

Types:

Type, Size -- see erlang:memory/0

Memory allocation information. Equivalent to `erlang:memory/0`.

`memory(Type) -> Size`

`memory([Type]) -> [{Type, Size}]`

Types:

Type, Size -- see erlang:memory/0

Memory allocation information. Equivalent to `erlang:memory/1`.

`nc(File) -> {ok, Module} | error`

`nc(File, Options) -> {ok, Module} | error`

Types:

File = name() -- see filename(3)

Options = [Opt] -- see compile:file/2

Compiles and then loads the code for a file on all nodes. `Options` defaults to []. Compilation is equivalent to:

```
compile:file(File, Opts ++ [report_errors, report_warnings])
```

`nl(Module) -> void()`

Types:

Module = atom()

Loads `Module` on all nodes.

`pid(X, Y, Z) -> pid()`

Types:

X = Y = Z = int()

Converts `X`, `Y`, `Z` to the pid `<X.Y.Z>`. This function should only be used when debugging.

`pwd() -> void()`

Prints the name of the working directory.

`q() -> void()`

This function is shorthand for `init:stop()`, that is, it causes the node to stop in a controlled fashion.

regs() -> void()

nregs() -> void()

`regs/0` displays information about all registered processes. `nregs/0` does the same, but for all nodes in the network.

xm(ModSpec) -> void()

Types:

ModSpec = Module | Filename

Module = atom()

Filename = string()

This function finds undefined functions, unused functions, and calls to deprecated functions in a module by calling `xref:m/1`.

y(File) -> YeccRet

Types:

File = name() -- see filename(3)

YeccRet = -- see yecc:file/2

Generates an LALR-1 parser. Equivalent to:

```
yecc:file(File)
```

y(File, Options) -> YeccRet

Types:

File = name() -- see filename(3)

Options, YeccRet = -- see yecc:file/2

Generates an LALR-1 parser. Equivalent to:

```
yecc:file(File, Options)
```

See Also

compile(3), *filename(3)*, *erlang(3)*, *yecc(3)*, *xref(3)*

calendar

Erlang module

Local Time本地时间是考虑了时区和夏时令的,Universal Time世界时间是计算的0经度的时间也没有考虑夏令时

This module provides computation of local and universal time, day-of-the-week, and several time conversion functions.

0经度

Time is local when it is adjusted in accordance with the current time zone and daylight saving. Time is universal when it reflects the time at longitude zero, without any adjustment for daylight saving. Universal Coordinated Time (UTC) time is also called Greenwich Mean Time (GMT). daylight saving time n. 夏令时

The time functions `local_time/0` and `universal_time/0` provided in this module both return date and time. The reason for this is that separate functions for date and time may result in a date/time combination which is displaced by 24 hours. This happens if one of the functions is called before midnight, and the other after midnight. This problem also applies to the Erlang BIFs `date/0` and `time/0`, and their use is strongly discouraged if a reliable date/time stamp is required.

All dates conform to the Gregorian calendar. This calendar was introduced by Pope Gregory XIII in 1582 and was used in all Catholic countries from this year. Protestant parts of Germany and the Netherlands adopted it in 1698, England followed in 1752, and Russia in 1918 (the October revolution of 1917 took place in November according to the Gregorian calendar).

The Gregorian calendar in this module is extended back to year 0. For a given date, the *gregorian days* is the number of days up to and including the date specified. Similarly, the *gregorian seconds* for a given date and time, is the the number of seconds up to and including the specified date and time.

For computing differences between epochs in time, use the functions counting gregorian days or seconds. If epochs are given as local time, they must be converted to universal time, in order to get the correct value of the elapsed time between epochs. Use of the function `time_difference/2` is discouraged.

计算与公园纪元的时间差可以计算公立的天和秒,如果时间是本地时间需要转换成世界时间才能得到结果;不建议使用time_difference

DATA TYPES

```
date() = {Year, Month, Day}
  Year = int()
  Month = 1..12
  Day = 1..31
Year cannot be abbreviated. Example: 93 denotes year 93, not 1993.
Valid range depends on the underlying OS.
The date tuple must denote a valid date.

time() = {Hour, Minute, Second}
  Hour = 0..23
  Minute = Second = 0..59
```

Exports

`date_to_gregorian_days(Date) -> Days`

`date_to_gregorian_days(Year, Month, Day) -> Days`

Types:

Date = date()

Days = int()

This function computes the number of gregorian days starting with year 0 and ending at the given date.

```
5> calendar:date_to_gregorian_days({1,1,2}).
367
6> calendar:date_to_gregorian_days({0,1,2}).
1
```

datetime_to_gregorian_seconds({Date, Time}) -> Seconds

Types:

Date = date()

Time = time()

Seconds = int()

This function computes the number of gregorian seconds starting with year 0 and ending at the given date and time.

day_of_the_week(Date) -> DayNumber

day_of_the_week(Year, Month, Day) -> DayNumber

Types:

Date = date()

DayNumber = 1..7

```
13> calendar:day_of_the_week(1984,5,15).
2
14> calendar:day_of_the_week(2012,5,17).
4
```

This function computes the day of the week given `Year`, `Month` and `Day`. The return value denotes the day of the week as `1`: Monday, `2`: Tuesday, and so on.

gregorian_days_to_date(Days) -> Date

Types:

Days = int()

Date = date()

This function computes the date given the number of gregorian days.

gregorian_seconds_to_datetime(Seconds) -> {Date, Time}

Types:

Seconds = int()

Date = date()

Time = time()

This function computes the date and time from the given number of gregorian seconds.

is_leap_year(Year) -> bool()

```
%%是否闰年
```

This function checks if a year is a leap year.

last_day_of_the_month(Year, Month) -> int()

This function computes the number of days in a month.

```
%%某年某月有多少天
```

local_time() -> {Date, Time}

Types:

Date = date()

Time = time()

This function returns the local time reported by the underlying operating system.

`local_time_to_universal_time({Date1, Time1}) -> {Date2, Time2}`

This function converts from local time to Universal Coordinated Time (UTC). `Date1` must refer to a local date after Jan 1, 1970.

> Warning:
>
> This function is deprecated. Use `local_time_to_universal_time_dst/1` instead, as it gives a more correct and complete result. Especially for the period that does not exist since it gets skipped during the switch *to* daylight saving time, this function still returns a result.

`local_time_to_universal_time_dst({Date1, Time1}) -> [{Date, Time}]`

Types:

Date1 = Date = date()

Time1 = Time = time()

```
18> calendar:local_time_to_universal_time_dst(calendar:local_time()).
[{{2012,5,17},{5,18,8}}]
```

This function converts from local time to Universal Coordinated Time (UTC). `Date1` must refer to a local date after Jan 1, 1970.

The return value is a list of 0, 1 or 2 possible UTC times:

`[]`

For a local `{Date1, Time1}` during the period that is skipped when switching *to* daylight saving time, there is no corresponding UTC since the local time is illegal - it has never happened.

`[DstDateTimeUTC, DateTimeUTC]`

For a local `{Date1, Time1}` during the period that is repeated when switching *from* daylight saving time, there are two corresponding UTCs. One for the first instance of the period when daylight saving time is still active, and one for the second instance.

`[DateTimeUTC]`

For all other local times there is only one corresponding UTC.

`now_to_local_time(Now) -> {Date, Time}`

Types:

Now -- see erlang:now/0

Date = date()

Time = time()

```
calendar:now_to_local_time(erlang:now()).
{{2012,5,17},{13,34,2}}

2> erlang:now().
{1337,233146,582368}
```

This function returns local date and time converted from the return value from `erlang:now()`.

`now_to_universal_time(Now) -> {Date, Time}`

`now_to_datetime(Now) -> {Date, Time}`

Types:

Now -- see erlang:now/0

Date = date()

Time = time()

```
calendar:now_to_universal_time(erlang:now()).
{{2012,5,17},{5,39,38}}
```

This function returns Universal Coordinated Time (UTC) converted from the return value from `erlang:now()`.

```
%% 秒->天 时间
```

`seconds_to_daystime(Seconds) -> {Days, Time}`

Types:

Seconds = Days = int()

Time = time()

This function transforms a given number of seconds into days, hours, minutes, and seconds. The `Time` part is always non-negative, but `Days` is negative if the argument `Seconds` is.

`seconds_to_time(Seconds) -> Time`

Types:

Seconds = int() < 86400

Time = time()

This function computes the time from the given number of seconds. `Seconds` must be less than the number of seconds per day (86400).

`time_difference(T1, T2) -> {Days, Time}`

This function returns the difference between two `{Date, Time}` tuples. `T2` should refer to an epoch later than `T1`.

> Warning:
>
> This function is obsolete. Use the conversion functions for gregorian days and seconds instead.

`time_to_seconds(Time) -> Seconds`

Types:

Time = time()

Seconds = int()

This function computes the number of seconds since midnight up to the specified time.

`universal_time() -> {Date, Time}`

Types:

Date = date()

Time = time()

```
16> calendar:universal_time().
{{2012,5,17},{5,5,19}}
17> calendar:local_time().
{{2012,5,17},{13,5,44}}
```

This function returns the Universal Coordinated Time (UTC) reported by the underlying operating system. Local time is returned if universal time is not available.

`universal_time_to_local_time({Date1, Time1}) -> {Date2, Time2}`

Types:

Date1 = Date2 = date()

Time1 = Time2 = time()

This function converts from Universal Coordinated Time (UTC) to local time. `Date1` must refer to a date after Jan 1, 1970.

valid_date(Date) -> bool()

valid_date(Year, Month, Day) -> bool()

Types:

Date = date()

This function checks if a date is a valid.

Leap Years

The notion that every fourth year is a leap year is not completely true. By the Gregorian rule, a year Y is a leap year if either of the following rules is valid:

- Y is divisible by 4, but not by 100; or
- Y is divisible by 400.

Accordingly, 1996 is a leap year, 1900 is not, but 2000 is.

Date and Time Source

Local time is obtained from the Erlang BIF `localtime/0`. Universal time is computed from the BIF `universaltime/0`.

The following facts apply:

- there are 86400 seconds in a day
- there are 365 days in an ordinary year
- there are 366 days in a leap year
- there are 1461 days in a 4 year period
- there are 36524 days in a 100 year period
- there are 146097 days in a 400 year period
- there are 719528 days between Jan 1, 0 and Jan 1, 1970.

dets

Erlang module

The module `dets` provides a term storage on file. The stored terms, in this module called *objects*, are tuples such that one element is defined to be the key. A Dets *table* is a collection of objects with the key at the same position stored on a file.

Dets is used by the Mnesia application, and is provided as is for users who are interested in an efficient storage of Erlang terms on disk only. Many applications just need to store some terms in a file. Mnesia adds transactions, queries, and distribution. The size of Dets files cannot exceed 2 GB. If larger tables are needed, Mnesia's table fragmentation can be used.

There are three types of Dets tables: set, bag and duplicate_bag. A table of type *set* has at most one object with a given key. If an object with a key already present in the table is inserted, the existing object is overwritten by the new object. A table of type *bag* has zero or more different objects with a given key. A table of type *duplicate_bag* has zero or more possibly matching objects with a given key.

Dets tables must be opened before they can be updated or read, and when finished they must be properly closed. If a table has not been properly closed, Dets will automatically repair the table. This can take a substantial time if the table is large. A Dets table is closed when the process which opened the table terminates. If several Erlang processes (users) open the same Dets table, they will share the table. The table is properly closed when all users have either terminated or closed the table. Dets tables are not properly closed if the Erlang runtime system is terminated abnormally.

> ### Note:
>
> A ^C command abnormally terminates an Erlang runtime system in a Unix environment with a break-handler.

Since all operations performed by Dets are disk operations, it is important to realize that a single look-up operation involves a series of disk seek and read operations. For this reason, the Dets functions are much slower than the corresponding Ets functions, although Dets exports a similar interface.

Dets organizes data as a linear hash list and the hash list grows gracefully as more data is inserted into the table. Space management on the file is performed by what is called a buddy system. The current implementation keeps the entire buddy system in RAM, which implies that if the table gets heavily fragmented, quite some memory can be used up. The only way to defragment a table is to close it and then open it again with the `repair` option set to `force`.

It is worth noting that the ordered_set type present in Ets is not yet implemented by Dets, neither is the limited support for concurrent updates which makes a sequence of `first` and `next` calls safe to use on fixed Ets tables. Both these features will be implemented by Dets in a future release of Erlang/OTP. Until then, the Mnesia application (or some user implemented method for locking) has to be used to implement safe concurrency. Currently, no library of Erlang/OTP has support for ordered disk based term storage.

Two versions of the format used for storing objects on file are supported by Dets. The first version, 8, is the format always used for tables created by OTP R7 and earlier. The second version, 9, is the default version of tables created by OTP R8 (and later OTP releases). OTP R8 can create version 8 tables, and convert version 8 tables to version 9, and vice versa, upon request.

All Dets functions return `{error, Reason}` if an error occurs (`first/1` and `next/2` are exceptions, they exit the process with the error tuple). If given badly formed arguments, all functions exit the process with a `badarg` message.

Types

```
access() = read | read_write
auto_save() = infinity | int()
bindings_cont() = tuple()
bool() = true | false
file() = string()
int() = integer() >= 0
keypos() = integer() >= 1
name() = atom() | reference()
no_slots() = integer() >= 0 | default
object() = tuple()
object_cont() = tuple()
select_cont() = tuple()
type() = bag | duplicate_bag | set
version() = 8 | 9 | default
```

Exports

`all() -> [Name]`

Types:

Name = name()

Returns a list of the names of all open tables on this node.

`bchunk(Name, Continuation) -> {Continuation2, Data} | '$end_of_table' | {error, Reason}`

Types:

Name = name()

Continuation = start | cont()

Continuation2 = cont()

Data = binary() | tuple()

Returns a list of objects stored in a table. The exact representation of the returned objects is not public. The lists of data can be used for initializing a table by giving the value `bchunk` to the `format` option of the `init_table/3` function. The Mnesia application uses this function for copying open tables.

Unless the table is protected using `safe_fixtable/2`, calls to `bchunk/2` may not work as expected if concurrent updates are made to the table.

The first time `bchunk/2` is called, an initial continuation, the atom `start`, must be provided.

The `bchunk/2` function returns a tuple `{Continuation2, Data}`, where `Data` is a list of objects. `Continuation2` is another continuation which is to be passed on to a subsequent call to `bchunk/2`. With a series of calls to `bchunk/2` it is possible to extract all objects of the table.

`bchunk/2` returns `'$end_of_table'` when all objects have been returned, or `{error, Reason}` if an error occurs.

`close(Name) -> ok | {error, Reason}`

Types:

Name = name()

Closes a table. Only processes that have opened a table are allowed to close it.

All open tables must be closed before the system is stopped. If an attempt is made to open a table which has not been properly closed, Dets automatically tries to repair the table.

delete(Name, Key) -> ok | {error, Reason}

Types:

Name = name()

Deletes all objects with the key `Key` from the table `Name`.

delete_all_objects(Name) -> ok | {error, Reason}

Types:

Name = name()

Deletes all objects from a table in almost constant time. However, if the table if fixed, `delete_all_objects(T)` is equivalent to `match_delete(T, '_')`.

delete_object(Name, Object) -> ok | {error, Reason}

Types:

Name = name()

Object = object()

Deletes all instances of a given object from a table. If a table is of type `bag` or `duplicate_bag`, the `delete/2` function cannot be used to delete only some of the objects with a given key. This function makes this possible.

first(Name) -> Key | '$end_of_table'

Types:

Key = term()

Name = name()

Returns the first key stored in the table `Name` according to the table's internal order, or `'$end_of_table'` if the table is empty.

Unless the table is protected using `safe_fixtable/2`, subsequent calls to `next/2` may not work as expected if concurrent updates are made to the table.

Should an error occur, the process is exited with an error tuple `{error, Reason}`. The reason for not returning the error tuple is that it cannot be distinguished from a key.

There are two reasons why `first/1` and `next/2` should not be used: they are not very efficient, and they prevent the use of the key `'$end_of_table'` since this atom is used to indicate the end of the table. If possible, the `match`, `match_object`, and `select` functions should be used for traversing tables.

foldl(Function, Acc0, Name) -> Acc1 | {error, Reason}

Types:

Function = fun(Object, AccIn) -> AccOut

Acc0 = Acc1 = AccIn = AccOut = term()

Name = name()

Object = object()

Calls `Function` on successive elements of the table `Name` together with an extra argument `AccIn`. The order in which the elements of the table are traversed is unspecified. `Function` must return a new accumulator which is passed to the next call. `Acc0` is returned if the table is empty.

foldr(Function, Acc0, Name) -> Acc1 | {error, Reason}

Types:

Function = fun(Object, AccIn) -> AccOut

Acc0 = Acc1 = AccIn = AccOut = term()

Name = name()

Object = object()

Calls `Function` on successive elements of the table `Name` together with an extra argument `AccIn`. The order in which the elements of the table are traversed is unspecified. `Function` must return a new accumulator which is passed to the next call. `Acc0` is returned if the table is empty.

`from_ets(Name, EtsTab) -> ok | {error, Reason}`

Types:

Name = name()

EtsTab = - see ets(3) -

Deletes all objects of the table `Name` and then inserts all the objects of the Ets table `EtsTab`. The order in which the objects are inserted is not specified. Since `ets:safe_fixtable/2` is called the Ets table must be public or owned by the calling process.

`info(Name) -> InfoList | undefined`

Types:

Name = name()

InfoList = [{Item, Value}]

Returns information about the table `Name` as a list of `{Item, Value}` tuples:

- `{file_size, int()}`, the size of the file in bytes.
- `{filename, file()}`, the name of the file where objects are stored.
- `{keypos, keypos()}`, the position of the key.
- `{size, int()}`, the number of objects stored in the table.
- `{type, type()}`, the type of the table.

`info(Name, Item) -> Value | undefined`

Types:

Name = name()

Returns the information associated with `Item` for the table `Name`. In addition to the `{Item, Value}` pairs defined for `info/1`, the following items are allowed:

- `{access, access()}`, the access mode.
- `{auto_save, auto_save()}`, the auto save interval.
- `{bchunk_format, binary()}`, an opaque binary describing the format of the objects returned by `bchunk/2`. The binary can be used as argument to `is_compatible_chunk_format/2`. Only available for version 9 tables.
- `{hash, Hash}`. Describes which BIF is used to calculate the hash values of the objects stored in the Dets table. Possible values of Hash are `hash`, which implies that the `erlang:hash/2` BIF is used, `phash`, which implies that the `erlang:phash/2` BIF is used, and `phash2`, which implies that the `erlang:phash2/1` BIF is used.
- `{memory, int()}`, the size of the file in bytes. The same value is associated with the item `file_size`.
- `{no_keys, int()}`, the number of different keys stored in the table. Only available for version 9 tables.
- `{no_objects, int()}`, the number of objects stored in the table.

- `{no_slots, {Min, Used, Max}}`, the number of slots of the table. `Min` is the minimum number of slots, `Used` is the number of currently used slots, and `Max` is the maximum number of slots. Only available for version 9 tables.
- `{owner, pid()}`, the pid of the process that handles requests to the Dets table.
- `{ram_file, bool()}`, whether the table is kept in RAM.
- `{safe_fixed,` SafeFixed`}`. If the table is fixed, SafeFixed is a tuple `{FixedAtTime, [{Pid,RefCount}]}`. `FixedAtTime` is the time when the table was first fixed, and `Pid` is the pid of the process that fixes the table `RefCount` times. There may be any number of processes in the list. If the table is not fixed, SafeFixed is the atom `false`.
- `{version, int()}`, the version of the format of the table.

`init_table(Name, InitFun [, Options]) -> ok | {error, Reason}`

Types:

Name = atom()

InitFun = fun(Arg) -> Res

Arg = read | close

Res = end_of_input | {[object()], InitFun} | {Data, InitFun} | term()

Data = binary() | tuple()

Replaces the existing objects of the table `Name` with objects created by calling the input function `InitFun`, see below. The reason for using this function rather than calling `insert/2` is that of efficiency. It should be noted that the input functions are called by the process that handles requests to the Dets table, not by the calling process.

When called with the argument `read` the function `InitFun` is assumed to return `end_of_input` when there is no more input, or `{Objects, Fun}`, where `Objects` is a list of objects and `Fun` is a new input function. Any other value Value is returned as an error `{error, {init_fun, Value}}`. Each input function will be called exactly once, and should an error occur, the last function is called with the argument `close`, the reply of which is ignored.

If the type of the table is `set` and there is more than one object with a given key, one of the objects is chosen. This is not necessarily the last object with the given key in the sequence of objects returned by the input functions. Duplicate keys should be avoided, or the file will be unnecessarily fragmented. This holds also for duplicated objects stored in tables of type `bag`.

It is important that the table has a sufficient number of slots for the objects. If not, the hash list will start to grow when `init_table/2` returns which will significantly slow down access to the table for a period of time. The minimum number of slots is set by the `open_file/2` option `min_no_slots` and returned by the `info/2` item `no_slots`. See also the `min_no_slots` option below.

The `Options` argument is a list of `{Key, Val}` tuples where the following values are allowed:

- `{min_no_slots, no_slots()}`. Specifies the estimated number of different keys that will be stored in the table. The `open_file` option with the same name is ignored unless the table is created, and in that case performance can be enhanced by supplying an estimate when initializing the table.
- `{format, Format}`. Specifies the format of the objects returned by the function `InitFun`. If `Format` is `term` (the default), `InitFun` is assumed to return a list of tuples. If `Format` is `bchunk`, `InitFun` is assumed to return `Data` as returned by `bchunk/2`. This option overrides the `min_no_slots` option.

`insert(Name, Objects) -> ok | {error, Reason}`

Types:

Name = name()

Objects = object() | [object()]

Inserts one or more objects into the table `Name`. If there already exists an object with a key matching the key of some of the given objects and the table type is `set`, the old object will be replaced.

`insert_new(Name, Objects) -> Bool`

Types:

Name = name()

Objects = object() | [object()]

Bool = bool()

Inserts one or more objects into the table `Name`. If there already exists some object with a key matching the key of any of the given objects the table is not updated and `false` is returned, otherwise the objects are inserted and `true` returned.

`is_compatible_bchunk_format(Name, BchunkFormat) -> Bool`

Types:

Name = name()

BchunkFormat = binary()

Bool = bool()

Returns `true` if it would be possible to initialize the table `Name`, using `init_table/3` with the option `{format, bchunk}`, with objects read with `bchunk/2` from some table `T` such that calling `info(T, bchunk_format)` returns `BchunkFormat`.

`is_dets_file(FileName) -> Bool | {error, Reason}`

Types:

FileName = file()

Bool = bool()

Returns `true` if the file `FileName` is a Dets table, `false` otherwise.

`lookup(Name, Key) -> [Object] | {error, Reason}`

Types:

Key = term()

Name = name()

Object = object()

Returns a list of all objects with the key `Key` stored in the table `Name`. For example:

```
2> dets:open_file(abc, [{type, bag}]).
{ok,abc}
3> dets:insert(abc, {1,2,3}).
ok
4> dets:insert(abc, {1,3,4}).
ok
5> dets:lookup(abc, 1).
[{1,2,3},{1,3,4}]
```

If the table is of type `set`, the function returns either the empty list or a list with one object, as there cannot be more than one object with a given key. If the table is of type `bag` or `duplicate_bag`, the function returns a list of arbitrary length.

Note that the order of objects returned is unspecified. In particular, the order in which objects were inserted is not reflected.

match(Continuation) -> {[Match], Continuation2} | '$end_of_table' | {error, Reason}

Types:

Continuation = Continuation2 = bindings_cont()

Match = [term()]

Matches some objects stored in a table and returns a non-empty list of the bindings that match a given pattern in some unspecified order. The table, the pattern, and the number of objects that are matched are all defined by `Continuation`, which has been returned by a prior call to `match/1` or `match/3`.

When all objects of the table have been matched, `'$end_of_table'` is returned.

match(Name, Pattern) -> [Match] | {error, Reason}

Types:

Name = name()

Pattern = tuple()

Match = [term()]

Returns for each object of the table `Name` that matches `Pattern` a list of bindings in some unspecified order. See *ets(3)* for a description of patterns. If the keypos'th element of `Pattern` is unbound, all objects of the table are matched. If the keypos'th element is bound, only the objects with the right key are matched.

match(Name, Pattern, N) -> {[Match], Continuation} | '$end_of_table' | {error, Reason}

Types:

Name = name()

Pattern = tuple()

N = default | int()

Match = [term()]

Continuation = bindings_cont()

Matches some or all objects of the table `Name` and returns a non-empty list of the bindings that match `Pattern` in some unspecified order. See *ets(3)* for a description of patterns.

A tuple of the bindings and a continuation is returned, unless the table is empty, in which case `'$end_of_table'` is returned. The continuation is to be used when matching further objects by calling `match/1`.

If the keypos'th element of `Pattern` is bound, all objects of the table are matched. If the keypos'th element is unbound, all objects of the table are matched, `N` objects at a time, until at least one object matches or the end of the table has been reached. The default, indicated by giving `N` the value `default`, is to let the number of objects vary depending on the sizes of the objects. If `Name` is a version 9 table, all objects with the same key are always matched at the same time which implies that more than N objects may sometimes be matched.

The table should always be protected using `safe_fixtable/2` before calling `match/3`, or errors may occur when calling `match/1`.

match_delete(Name, Pattern) -> ok | {error, Reason}

Types:

Name = name()

Pattern = tuple()

Deletes all objects that match `Pattern` from the table `Name`. See *ets:match/2* for a description of patterns.

If the keypos'th element of `Pattern` is bound, only the objects with the right key are matched.

`match_object(Continuation) -> {[Object], Continuation2} | '$end_of_table' | {error, Reason}`

Types:

Continuation = Continuation2 = object_cont()

Object = object()

Returns a non-empty list of some objects stored in a table that match a given pattern in some unspecified order. The table, the pattern, and the number of objects that are matched are all defined by `Continuation`, which has been returned by a prior call to `match_object/1` or `match_object/3`.

When all objects of the table have been matched, `'$end_of_table'` is returned.

`match_object(Name, Pattern) -> [Object] | {error, Reason}`

Types:

Name = name()

Pattern = tuple()

Object = object()

Returns a list of all objects of the table `Name` that match `Pattern` in some unspecified order. See *ets(3)* for a description of patterns.

If the keypos'th element of `Pattern` is unbound, all objects of the table are matched. If the keypos'th element of `Pattern` is bound, only the objects with the right key are matched.

Using the `match_object` functions for traversing all objects of a table is more efficient than calling `first/1` and `next/2` or `slot/2`.

`match_object(Name, Pattern, N) -> {[Object], Continuation} | '$end_of_table' | {error, Reason}`

Types:

Name = name()

Pattern = tuple()

N = default | int()

Object = object()

Continuation = object_cont()

Matches some or all objects stored in the table `Name` and returns a non-empty list of the objects that match `Pattern` in some unspecified order. See *ets(3)* for a description of patterns.

A list of objects and a continuation is returned, unless the table is empty, in which case `'$end_of_table'` is returned. The continuation is to be used when matching further objects by calling `match_object/1`.

If the keypos'th element of `Pattern` is bound, all objects of the table are matched. If the keypos'th element is unbound, all objects of the table are matched, `N` objects at a time, until at least one object matches or the end of the table has been reached. The default, indicated by giving `N` the value `default`, is to let the number of objects vary depending on the sizes of the objects. If `Name` is a version 9 table, all matching objects with the same key are always returned in the same reply which implies that more than N objects may sometimes be returned.

The table should always be protected using `safe_fixtable/2` before calling `match_object/3`, or errors may occur when calling `match_object/1`.

`member(Name, Key) -> Bool | {error, Reason}`

Types:

Name = **name()**

Key = **term()**

Bool = **bool()**

Works like `lookup/2`, but does not return the objects. The function returns `true` if one or more elements of the table has the key `Key`, `false` otherwise.

`next(Name, Key1) -> Key2 | '$end_of_table'`

Types:

Name = **name()**

Key1 = **Key2** = **term()**

Returns the key following `Key1` in the table `Name` according to the table's internal order, or `'$end_of_table'` if there is no next key.

Should an error occur, the process is exited with an error tuple `{error, Reason}`.

Use `first/1` to find the first key in the table.

`open_file(Filename) -> {ok, Reference} | {error, Reason}`

Types:

FileName = **file()**

Reference = **reference()**

Opens an existing table. If the table has not been properly closed, it will be repaired. The returned reference is to be used as the name of the table. This function is most useful for debugging purposes.

`open_file(Name, Args) -> {ok, Name} | {error, Reason}`

Types:

Name = **atom()**

Opens a table. An empty Dets table is created if no file exists.

The atom `Name` is the name of the table. The table name must be provided in all subsequent operations on the table. The name can be used by other processes as well, and several process can share one table.

If two processes open the same table by giving the same name and arguments, then the table will have two users. If one user closes the table, it still remains open until the second user closes the table.

The `Args` argument is a list of `{Key, Val}` tuples where the following values are allowed:

- `{access, access()}`. It is possible to open existing tables in read-only mode. A table which is opened in read-only mode is not subjected to the automatic file reparation algorithm if it is later opened after a crash. The default value is `read_write`.
- `{auto_save, auto_save()}`, the auto save interval. If the interval is an integer `Time`, the table is flushed to disk whenever it is not accessed for `Time` milliseconds. A table that has been flushed will require no reparation when reopened after an uncontrolled emulator halt. If the interval is the atom `infinity`, auto save is disabled. The default value is 180000 (3 minutes).
- `{estimated_no_objects, int()}`. Equivalent to the `min_no_slots` option.

- `{file, file()}`, the name of the file to be opened. The default value is the name of the table.
- `{max_no_slots, no_slots()}`, the maximum number of slots that will be used. The default value as well as the maximal value is 32 M. Note that a higher value may increase the fragmentation of the table, and conversely, that a smaller value may decrease the fragmentation, at the expense of execution time. Only available for version 9 tables.
- `{min_no_slots, no_slots()}`. Application performance can be enhanced with this flag by specifying, when the table is created, the estimated number of different keys that will be stored in the table. The default value as well as the minimum value is 256.
- `{keypos, keypos()}`, the position of the element of each object to be used as key. The default value is 1. The ability to explicitly state the key position is most convenient when we want to store Erlang records in which the first position of the record is the name of the record type.
- `{ram_file, bool()}`, whether the table is to be kept in RAM. Keeping the table in RAM may sound like an anomaly, but can enhance the performance of applications which open a table, insert a set of objects, and then close the table. When the table is closed, its contents are written to the disk file. The default value is `false`.
- `{repair, Value}`. `Value` can be either a `bool()` or the atom `force`. The flag specifies whether the Dets server should invoke the automatic file reparation algorithm. The default is `true`. If `false` is specified, there is no attempt to repair the file and `{error, {needs_repair, FileName}}` is returned if the table needs to be repaired.

  The value `force` means that a reparation will take place even if the table has been properly closed. This is how to convert tables created by older versions of STDLIB. An example is tables hashed with the deprecated `erlang:hash/2` BIF. Tables created with Dets from a STDLIB version of 1.8.2 and later use the `erlang:phash/2` function or the `erlang:phash2/1` function, which is preferred.

  The `repair` option is ignored if the table is already open.
- `{type, type()}`, the type of the table. The default value is `set`.
- `{version, version()}`, the version of the format used for the table. The default value is `9`. Tables on the format used before OTP R8 can be created by giving the value `8`. A version 8 table can be converted to a version 9 table by giving the options `{version,9}` and `{repair,force}`.

`pid2name(Pid) -> {ok, Name} | undefined`

Types:

Name = name()

Pid = pid()

Returns the name of the table given the pid of a process that handles requests to a table, or `undefined` if there is no such table.

This function is meant to be used for debugging only.

`repair_continuation(Continuation, MatchSpec) -> Continuation2`

Types:

Continuation = Continuation2 = select_cont()

MatchSpec = match_spec()

This function can be used to restore an opaque continuation returned by `select/3` or `select/1` if the continuation has passed through external term format (been sent between nodes or stored on disk).

The reason for this function is that continuation terms contain compiled match specifications and therefore will be invalidated if converted to external term format. Given that the original match specification is kept intact, the continuation can be restored, meaning it can once again be used in subsequent `select/1` calls even though it has been stored on disk or on another node.

See also `ets(3)` for further explanations and examples.

> Note:
>
> This function is very rarely needed in application code. It is used by Mnesia to implement distributed `select/3` and `select/1` sequences. A normal application would either use Mnesia or keep the continuation from being converted to external format.
>
> The reason for not having an external representation of compiled match specifications is performance. It may be subject to change in future releases, while this interface will remain for backward compatibility.

`safe_fixtable(Name, Fix)`

Types:

Name = name()

Fix = bool()

If `Fix` is `true`, the table `Name` is fixed (once more) by the calling process, otherwise the table is released. The table is also released when a fixing process terminates.

If several processes fix a table, the table will remain fixed until all processes have released it or terminated. A reference counter is kept on a per process basis, and N consecutive fixes require N releases to release the table.

It is not guaranteed that calls to `first/1`, `next/2`, select and match functions work as expected even if the table has been fixed; the limited support for concurrency implemented in Ets has not yet been implemented in Dets. Fixing a table currently only disables resizing of the hash list of the table.

If objects have been added while the table was fixed, the hash list will start to grow when the table is released which will significantly slow down access to the table for a period of time.

`select(Continuation) -> {Selection, Continuation2} | '$end_of_table' | {error, Reason}`

Types:

Continuation = Continuation2 = select_cont()

Selection = [term()]

Applies a match specification to some objects stored in a table and returns a non-empty list of the results. The table, the match specification, and the number of objects that are matched are all defined by `Continuation`, which has been returned by a prior call to `select/1` or `select/3`.

When all objects of the table have been matched, `'$end_of_table'` is returned.

`select(Name, MatchSpec) -> Selection | {error, Reason}`

Types:

Name = name()

MatchSpec = match_spec()

Selection = [term()]

Returns the results of applying the match specification `MatchSpec` to all or some objects stored in the table `Name`. The order of the objects is not specified. See the ERTS User's Guide for a description of match specifications.

If the keypos'th element of `MatchSpec` is unbound, the match specification is applied to all objects of the table. If the keypos'th element is bound, the match specification is applied to the objects with the right key(s) only.

Using the `select` functions for traversing all objects of a table is more efficient than calling `first/1` and `next/2` or `slot/2`.

`select(Name, MatchSpec, N) -> {Selection, Continuation} | '$end_of_table' | {error, Reason}`

Types:

Name = name()

MatchSpec = match_spec()

N = default | int()

Selection = [term()]

Continuation = select_cont()

Returns the results of applying the match specification `MatchSpec` to some or all objects stored in the table `Name`. The order of the objects is not specified. See the ERTS User's Guide for a description of match specifications.

A tuple of the results of applying the match specification and a continuation is returned, unless the table is empty, in which case `'$end_of_table'` is returned. The continuation is to be used when matching further objects by calling `select/1`.

If the keypos'th element of `MatchSpec` is bound, the match specification is applied to all objects of the table with the right key(s). If the keypos'th element of `MatchSpec` is unbound, the match specification is applied to all objects of the table, `N` objects at a time, until at least one object matches or the end of the table has been reached. The default, indicated by giving `N` the value `default`, is to let the number of objects vary depending on the sizes of the objects. If `Name` is a version 9 table, all objects with the same key are always handled at the same time which implies that the match specification may be applied to more than N objects.

The table should always be protected using `safe_fixtable/2` before calling `select/3`, or errors may occur when calling `select/1`.

`select_delete(Name, MatchSpec) -> N | {error, Reason}`

Types:

Name = name()

MatchSpec = match_spec()

N = int()

Deletes each object from the table `Name` such that applying the match specification `MatchSpec` to the object returns the value `true`. See the ERTS User's Guide for a description of match specifications. Returns the number of deleted objects.

If the keypos'th element of `MatchSpec` is bound, the match specification is applied to the objects with the right key(s) only.

`slot(Name, I) -> '$end_of_table' | [Object] | {error, Reason}`

Types:

Name = name()

I = int()

Object = object()

The objects of a table are distributed among slots, starting with slot `0` and ending with slot n. This function returns the list of objects associated with slot `I`. If `I` is greater than n `'$end_of_table'` is returned.

sync(Name) -> ok | {error, Reason}

Types:

Name = name()

Ensures that all updates made to the table `Name` are written to disk. This also applies to tables which have been opened with the `ram_file` flag set to `true`. In this case, the contents of the RAM file are flushed to disk.

Note that the space management data structures kept in RAM, the buddy system, is also written to the disk. This may take some time if the table is fragmented.

table(Name [, Options]) -> QueryHandle

Types:

Name = name()

QueryHandle = - a query handle, see qlc(3) -

Options = [Option] | Option

Option = {n_objects, Limit} | {traverse, TraverseMethod}

Limit = default | integer() >= 1

TraverseMethod = first_next | select | {select, MatchSpec}

MatchSpec = match_spec()

Returns a QLC (Query List Comprehension) query handle. The module `qlc` implements a query language aimed mainly at Mnesia but Ets tables, Dets tables, and lists are also recognized by `qlc` as sources of data. Calling `dets:table/1,2` is the means to make the Dets table `Name` usable to `qlc`.

When there are only simple restrictions on the key position `qlc` uses `dets:lookup/2` to look up the keys, but when that is not possible the whole table is traversed. The option `traverse` determines how this is done:

- `first_next`. The table is traversed one key at a time by calling `dets:first/1` and `dets:next/2`.
- `select`. The table is traversed by calling `dets:select/3` and `dets:select/1`. The option `n_objects` determines the number of objects returned (the third argument of `select/3`). The match specification (the second argument of `select/3`) is assembled by `qlc`: simple filters are translated into equivalent match specifications while more complicated filters have to be applied to all objects returned by `select/3` given a match specification that matches all objects.
- `{select, MatchSpec}`. As for `select` the table is traversed by calling `dets:select/3` and `dets:select/1`. The difference is that the match specification is explicitly given. This is how to state match specifications that cannot easily be expressed within the syntax provided by `qlc`.

The following example uses an explicit match specification to traverse the table:

```
1> dets:open_file(t, []),
ok = dets:insert(t, [{1,a},{2,b},{3,c},{4,d}]),
MS = ets:fun2ms(fun({X,Y}) when (X > 1) or (X < 5) -> {Y} end),
QH1 = dets:table(t, [{traverse, {select, MS}}]).
```

An example with implicit match specification:

```
2> QH2 = qlc:q([{Y} || {X,Y} <- dets:table(t), (X > 1) or (X < 5)]).
```

The latter example is in fact equivalent to the former which can be verified using the function `qlc:info/1`:

```
3> qlc:info(QH1) =:= qlc:info(QH2).
true
```

`qlc:info/1` returns information about a query handle, and in this case identical information is returned for the two query handles.

`to_ets(Name, EtsTab) -> EtsTab | {error, Reason}`

Types:

Name = name()

EtsTab = - see ets(3) -

Inserts the objects of the Dets table `Name` into the Ets table `EtsTab`. The order in which the objects are inserted is not specified. The existing objects of the Ets table are kept unless overwritten.

`traverse(Name, Fun) -> Return | {error, Reason}`

Types:

Fun = fun(Object) -> FunReturn

FunReturn = continue | {continue, Val} | {done, Value}

Val = Value = term()

Name = name()

Object = object()

Return = [term()]

Applies `Fun` to each object stored in the table `Name` in some unspecified order. Different actions are taken depending on the return value of `Fun`. The following `Fun` return values are allowed:

`continue`

Continue to perform the traversal. For example, the following function can be used to print out the contents of a table:

```
fun(X) -> io:format("~p~n", [X]), continue end.
```

`{continue, Val}`

Continue the traversal and accumulate `Val`. The following function is supplied in order to collect all objects of a table in a list:

```
fun(X) -> {continue, X} end.
```

`{done, Value}`

Terminate the traversal and return `[Value | Acc]`.

Any other value returned by `Fun` terminates the traversal and is immediately returned.

`update_counter(Name, Key, Increment) -> Result`

Types:

Name = name()

Key = term()

Increment = {Pos, Incr} | Incr

Pos = Incr = Result = integer()

Updates the object with key `Key` stored in the table `Name` of type `set` by adding `Incr` to the element at the `Pos`:th position. The new counter value is returned. If no position is specified, the element directly following the key is updated.

This functions provides a way of updating a counter, without having to look up an object, update the object by incrementing an element and insert the resulting object into the table again.

See Also

ets(3), mnesia(3), *qlc(3)*

dict

Erlang module

`Dict` implements a `Key` - `Value` dictionary. The representation of a dictionary is not defined.

This module provides exactly the same interface as the module `orddict`. One difference is that while this module considers two keys as different if they do not match (`=:=`), `orddict` considers two keys as different if and only if they do not compare equal (`==`). match操作符=:= 相等操作符==

DATA TYPES

```
dictionary()
  as returned by new/0
```

Exports

`append(Key, Value, Dict1) -> Dict2`

Types:

Key = Value = term()

Dict1 = Dict2 = dictionary()

```
1> dict:append(k,v,dict:new()).
{dict,1,16,16,8,80,48,
      {[],[],[],[],[],[],[],[],[],[],[],[],[],[],[],[]},
      {{[],[],[],[],[],[],[],[],[],[],[],[[k,v]],[],[],[],[]}}}
2> dict:append(k,v,v(1)).
{dict,1,16,16,8,80,48,
      {[],[],[],[],[],[],[],[],[],[],[],[],[],[],[],[]},
      {{[],[],[],[],[],[],[],[],[],[],[],[[k,v,v]],[],[],[],[]}}}
3>
```

This function appends a new `Value` to the current list of values associated with `Key`. An exception is generated if the initial value associated with `Key` is not a list of values.

`append_list(Key, ValList, Dict1) -> Dict2`

Types:

ValList = [Value]

Key = Value = term()

Dict1 = Dict2 = dictionary()

```
3> dict:append(k,[a,b,c,d,e,f,g,h],v(1)).
{dict,1,16,16,8,80,48,
      {[],[],[],[],[],[],[],[],[],[],[],[],[],[],[],[]},
      {{[],[],[],[],[],[],[],[],[],[],[],
        [[k,v,[a,b,c,d,e,f,g,h]]],
        [],[],[],[]}}}
ok
```

This function appends a list of values `ValList` to the current list of values associated with `Key`. An exception is generated if the initial value associated with `Key` is not a list of values.

`erase(Key, Dict1) -> Dict2`

Types:

Key = term()

Dict1 = Dict2 = dictionary()

```
5> dict:erase(k,v(3)).
{dict,0,16,16,8,80,48,
      {[],[],[],[],[],[],[],[],[],[],[],[],[],[],[],[]},
      {{[],[],[],[],[],[],[],[],[],[],[],[],[],[],[],[]}}}
```

This function erases all items with a given key from a dictionary.

`fetch(Key, Dict) -> Value`

Types:

Key = Value = term()

Dict = dictionary()

```
6> dict:fetch(k,v(3)).
[v,[a,b,c,d,e,f,g,h]]
```

This function returns the value associated with `Key` in the dictionary `Dict`. `fetch` assumes that the `Key` is present in the dictionary and an exception is generated if `Key` is not in the dictionary.

fetch_keys(Dict) -> Keys

Types:

Dict = dictionary()

Keys = [term()]

```
7> dict:fetch_keys(v(3)).
[k]
```

This function returns a list of all keys in the dictionary.

filter(Pred, Dict1) -> Dict2

Types:

Pred = fun(Key, Value) -> bool()

Key = Value = term()

Dict1 = Dict2 = dictionary()

```
8> dict:filter(fun(K,V) -> length(V)>0 end,v(3)).
{dict,1,16,16,8,80,48,
      {[],[],[],[],[],[],[],[],[],[],[],[],[],[],[],[]},
      {{[],[],[],[],[],[],[],[],[],[],[],[],
        [[k,v,[a,b,c|...]]],
        [],[],[],[]}}}
9> rp(v(8)).
{dict,1,16,16,8,80,48,
      {[],[],[],[],[],[],[],[],[],[],[],[],[],[],[],[]},
      {{[],[],[],[],[],[],[],[],[],[],[],[],
        [[k,v,[a,b,c,d,e,f,g,h]]],
        [],[],[],[]}}}
ok
10> dict:filter(fun(K,V) -> length(V)>9 end,v(3)).
{dict,0,16,16,8,80,48,
      {[],[],[],[],[],[],[],[],[],[],[],[],[],[],[],[]},
      {{[],[],[],[],[],[],[],[],[],[],[],[],[],[],[],[]}}}
11>
```

`Dict2` is a dictionary of all keys and values in `Dict1` for which `Pred(Key, Value)` is `true`.

find(Key, Dict) -> {ok, Value} | error

Types:

Key = Value = term()

Dict = dictionary()

This function searches for a key in a dictionary. Returns `{ok, Value}` where `Value` is the value associated with `Key`, or `error` if the key is not present in the dictionary.

fold(Fun, Acc0, Dict) -> Acc1

Types:

Fun = fun(Key, Value, AccIn) -> AccOut

Key = Value = term()

Acc0 = Acc1 = AccIn = AccOut = term()

Dict = dictionary()

Calls `Fun` on successive keys and values of `Dict` together with an extra argument `Acc` (short for accumulator). `Fun` must return a new accumulator which is passed to the next call. `Acc0` is returned if the list is empty. The evaluation order is undefined.

from_list(List) -> Dict

Types:

List = [{Key, Value}]

Dict = dictionary()

This function converts the key/value list `List` to a dictionary.

is_key(Key, Dict) -> bool()

Types:

Key = term()

Dict = dictionary()

This function tests if `Key` is contained in the dictionary `Dict`.

map(Fun, Dict1) -> Dict2

Types:

Fun = fun(Key, Value1) -> Value2

Key = Value1 = Value2 = term()

Dict1 = Dict2 = dictionary()

`map` calls `Func` on successive keys and values of `Dict` to return a new value for each key. The evaluation order is undefined.

merge(Fun, Dict1, Dict2) -> Dict3

Types:

Fun = fun(Key, Value1, Value2) -> Value

Key = Value1 = Value2 = Value3 = term()

Dict1 = Dict2 = Dict3 = dictionary()

`merge` merges two dictionaries, `Dict1` and `Dict2`, to create a new dictionary. All the `Key - Value` pairs from both dictionaries are included in the new dictionary. If a key occurs in both dictionaries then `Fun` is called with the key and both values to return a new value. `merge` could be defined as:

```
merge(Fun, D1, D2) ->
    fold(fun (K, V1, D) ->
                update(K, fun (V2) -> Fun(K, V1, V2) end, V1, D)
         end, D2, D1).
```

but is faster.

new() -> dictionary()

This function creates a new dictionary.

size(Dict) -> int()

Types:

Dict = dictionary()

Returns the number of elements in a `Dict`.

store(Key, Value, Dict1) -> Dict2

Types:

Key = Value = term()

Dict1 = Dict2 = dictionary()

This function stores a `Key - Value` pair in a dictionary. If the `Key` already exists in `Dict1`, the associated value is replaced by `Value`.

to_list(Dict) -> List

Types:

Dict = dictionary()

List = [{Key, Value}]

This function converts the dictionary to a list representation.

update(Key, Fun, Dict1) -> Dict2

Types:

Key = term()

Fun = fun(Value1) -> Value2

Value1 = Value2 = term()

Dict1 = Dict2 = dictionary()

Update the a value in a dictionary by calling `Fun` on the value to get a new value. An exception is generated if `Key` is not present in the dictionary.

update(Key, Fun, Initial, Dict1) -> Dict2

Types:

Key = Initial = term()

Fun = fun(Value1) -> Value2

Value1 = Value2 = term()

Dict1 = Dict2 = dictionary()

Update the a value in a dictionary by calling `Fun` on the value to get a new value. If `Key` is not present in the dictionary then `Initial` will be stored as the first value. For example `append/3` could be defined as:

```
append(Key, Val, D) ->
    update(Key, fun (Old) -> Old ++ [Val] end, [Val], D).
```

update_counter(Key, Increment, Dict1) -> Dict2

Types:

Key = term()

Increment = number()

Dict1 = Dict2 = dictionary()

Add `Increment` to the value associated with `Key` and store this value. If `Key` is not present in the dictionary then `Increment` will be stored as the first value.

This could be defined as:

```
update_counter(Key, Incr, D) ->
    update(Key, fun (Old) -> Old + Incr end, Incr, D).
```

but is faster.

Notes

The functions `append` and `append_list` are included so we can store keyed values in a list *accumulator*. For example:

```
> D0 = dict:new(),
  D1 = dict:store(files, [], D0),
  D2 = dict:append(files, f1, D1),
  D3 = dict:append(files, f2, D2),
```

```
  D4 = dict:append(files, f3, D3),
  dict:fetch(files, D4).
[f1,f2,f3]
```

This saves the trouble of first fetching a keyed value, appending a new value to the list of stored values, and storing the result.

The function `fetch` should be used if the key is known to be in the dictionary, otherwise `find`.

See Also

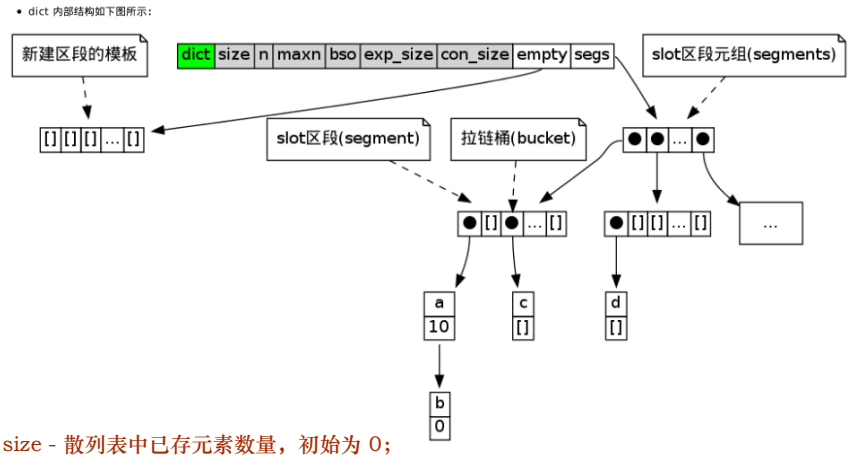
gb_trees(3), orddict(3)

[1] dict 基于论文 The Design and Implementation of Dynamic Hashing for Sets and Tables in Icon
论文中提出的动态散列技术构建，通过在散列 bucket 数量同存储元素数量之间维持一定比例，实现较高的内存使用率及访问效率。

[2] 同 array 结构采用的策略相似，整个散列 bucket 数组也被划分为若干个固定长度的区段（segments），
以便在更新特定 bucket 内容时不需要整体复制数据。目前每个segment 的大小设定为 16，是反复测试后选择的最佳值。

[3] dict 的扩张和紧缩都依赖于一定的元素数量阈值触发，该阈值基于当前活动的 bucket 数量按比例计算得出，目前设定为元素数量超过 bucket_no*5 时触发扩张操作，低于bucket_no*3 时触发紧缩操作。可能导致元素数量减少的操作都会在操作完成后尝试紧缩 bucket 区段；可能导致元素数量增加的操作都会在操作进行前尝试扩张 bucket 区段；

[4] bucket 中的多个键值对以无序列表形式存储，对 dict 的访问操作最终都会对特定的 bucket 进行遍历。只读操作的遍历是尾递归形式，bucket 较大时附加开销很小；但更新操作的遍历是普通递归形式，附加开销较大，所幸通常情况下 dict 的算法可以保证 bucket 不会很长。
dict 中的键值对本身以 improper list 形式保存，主要是希望以相同的代码高效处理 K-V（1对1） 和 K-Bag（1对多） 两种数据形式。



size - 散列表中已存元素数量，初始为 0；
n - 散列表中活动 bucket 数量，初始为 16；
maxn - 散列表中当前允许的最大 bucket 数量，扩张操作需要据此判断是否要增加新的 bucket 区段，初始为 16；
bso - 动态散列算法中计算一个 bucket 及其伙伴 bucket 的分隔下标，以此为界将 bucket 数组划分为对称的两半，初始为 8；
exp_size - 触发扩张操作的元素数量阈值，根据活动 bucket 数量和加载因子计算得出，初始为 16*5=80；
con_size - 触发紧缩操作的元素数量阈值，根据活动 bucket 数量和加载因子计算得出，初始为 16*3=48；
empty - 用于快速创建新 bucket 区段的区段模板，为 16 元 tuple，每个元素都是空列表；
segs - 散列 bucket 区段组，其下保存了所有的元素数据；

- 紧缩发生条件：
  - 操作完成后元素数量小于 dict 紧缩阈值 dict.con_size
  - 满足前者的条件下，dict 活动 bucket 数量 dict.n 大于 16
- 紧缩过程：每次紧缩掉一个 bucket，bucket 数量足够少时减半 bucket 区段
  - 计算出最后一个 bucket (dict.n) X 的伙伴 bucket Y 的位置 (dict.n-dict.bso) ；
  - 将 X 和 Y 的列表合并，替换掉 Y 列表，并清空 X 列表；
  - 活动 bucket 数量 dict.n 减 1，同时根据加载因子重新计算紧缩与扩充阈值；
  - 若活动 bucket 数量等于分隔下标 dict.bso，则将 bucket 区段、最大 bucket 数量 dict.maxn 和分隔下标 dict.bso 都减半。

- 扩张发生条件：
  - 操作进行前后元素数量大于 dict 扩张阈值 dict.exp_size
- 扩张过程：每次增加一个 bucket，bucket 数量增长到最大值时倍增 bucket 区段
  - 若活动 bucket 数量 dict.n 等于最大 bucket 数量 dict.maxn，则将 bucket 区段、最大 bucket 数量 dict.maxn 和分隔下标 dict.bso 都倍增；
  - 活动 bucket 数量 dict.n 加 1，作为新增 bucket X 的下标，同时根据加载因子重新计算紧缩与扩充阈值；
  - 计算出新增 bucket X 的伙伴 bucket Y 的位置 (dict.n-dict.bso) ；
  - 将 Y 的列表重新散列，分配到 X 和 Y 两个 bucket 里；

sets 同 dict 所用实现方法一致，只是存储元素从键值对变为了键本身，这里不再赘述。

digraph

Erlang module

The `digraph` module implements a version of labeled directed graphs. What makes the graphs implemented here non-proper directed graphs is that multiple edges between vertices are allowed. However, the customary definition of directed graphs will be used in the text that follows.

A *directed graph* (or just "digraph") is a pair (V, E) of a finite set V of *vertices* and a finite set E of *directed edges* (or just "edges"). The set of edges E is a subset of V × V (the Cartesian product of V with itself). In this module, V is allowed to be empty; the so obtained unique digraph is called the *empty digraph*. Both vertices and edges are represented by unique Erlang terms.

Digraphs can be annotated with additional information. Such information may be attached to the vertices and to the edges of the digraph. A digraph which has been annotated is called a *labeled digraph*, and the information attached to a vertex or an edge is called a *label*. Labels are Erlang terms.

An edge e = (v, w) is said to *emanate* from vertex v and to be *incident* on vertex w. The *out-degree* of a vertex is the number of edges emanating from that vertex. The *in-degree* of a vertex is the number of edges incident on that vertex. If there is an edge emanating from v and incident on w, then w is said to be an *out-neighbour* of v, and v is said to be an *in-neighbour* of w. A *path* P from v[1] to v[k] in a digraph (V, E) is a non-empty sequence v[1], v[2], ..., v[k] of vertices in V such that there is an edge (v[i],v[i+1]) in E for 1 <= i < k. The *length* of the path P is k-1. P is *simple* if all vertices are distinct, except that the first and the last vertices may be the same. P is a *cycle* if the length of P is not zero and v[1] = v[k]. A *loop* is a cycle of length one. A *simple cycle* is a path that is both a cycle and simple. An *acyclic digraph* is a digraph that has no cycles.

Exports

`add_edge(G, E, V1, V2, Label) -> edge() | {error, Reason}`

`add_edge(G, V1, V2, Label) -> edge() | {error, Reason}`

`add_edge(G, V1, V2) -> edge() | {error, Reason}`

Types:

G = digraph()

E = edge()

V1 = V2 = vertex()

Label = label()

Reason = {bad_edge, Path} | {bad_vertex, V}

Path = [vertex()]

`add_edge/5` creates (or modifies) the edge `E` of the digraph `G`, using `Label` as the (new) *label* of the edge. The edge is *emanating* from `V1` and *incident* on `V2`. Returns `E`.

`add_edge(G, V1, V2, Label)` is equivalent to `add_edge(G, E, V1, V2, Label)`, where `E` is a created edge. The created edge is represented by the term `['$e' | N]`, where N is an integer >= 0.

`add_edge(G, V1, V2)` is equivalent to `add_edge(G, V1, V2, [])`.

If the edge would create a cycle in an *acyclic digraph*, then `{error, {bad_edge, Path}}` is returned. If either of `V1` or `V2` is not a vertex of the digraph `G`, then `{error, {bad_vertex, V}}` is returned, V = `V1` or V = `V2`.

add_vertex(G, V, Label) -> vertex()

add_vertex(G, V) -> vertex()

add_vertex(G) -> vertex()

Types:

G = digraph()

V = vertex()

Label = label()

`add_vertex/3` creates (or modifies) the vertex `V` of the digraph `G`, using `Label` as the (new) *label* of the vertex. Returns `V`.

`add_vertex(G, V)` is equivalent to `add_vertex(G, V, [])`.

`add_vertex/1` creates a vertex using the empty list as label, and returns the created vertex. The created vertex is represented by the term `['$v' | N]`, where N is an integer >= 0.

del_edge(G, E) -> true

Types:

G = digraph()

E = edge()

Deletes the edge `E` from the digraph `G`.

del_edges(G, Edges) -> true

Types:

G = digraph()

Edges = [edge()]

Deletes the edges in the list `Edges` from the digraph `G`.

del_path(G, V1, V2) -> true

Types:

G = digraph()

V1 = V2 = vertex()

Deletes edges from the digraph `G` until there are no *paths* from the vertex `V1` to the vertex `V2`.

A sketch of the procedure employed: Find an arbitrary *simple path* v[1], v[2], ..., v[k] from `V1` to `V2` in `G`. Remove all edges of `G` *emanating* from v[i] and *incident* to v[i+1] for 1 <= i < k (including multiple edges). Repeat until there is no path between `V1` and `V2`.

del_vertex(G, V) -> true

Types:

G = digraph()

V = vertex()

Deletes the vertex `V` from the digraph `G`. Any edges *emanating* from `V` or *incident* on `V` are also deleted.

`del_vertices(G, Vertices) -> true`

Types:

G = digraph()

Vertices = [vertex()]

Deletes the vertices in the list `Vertices` from the digraph `G`.

`delete(G) -> true`

Types:

G = digraph()

Deletes the digraph `G`. This call is important because digraphs are implemented with `Ets`. There is no garbage collection of `Ets` tables. The digraph will, however, be deleted if the process that created the digraph terminates.

`edge(G, E) -> {E, V1, V2, Label} | false`

Types:

G = digraph()

E = edge()

V1 = V2 = vertex()

Label = label()

Returns `{E, V1, V2, Label}` where `Label` is the *label* of the edge `E` *emanating* from `V1` and *incident* on `V2` of the digraph `G`. If there is no edge `E` of the digraph `G`, then `false` is returned.

`edges(G) -> Edges`

Types:

G = digraph()

Edges = [edge()]

Returns a list of all edges of the digraph `G`, in some unspecified order.

`edges(G, V) -> Edges`

Types:

G = digraph()

V = vertex()

Edges = [edge()]

Returns a list of all edges *emanating* from or *incident* on `V` of the digraph `G`, in some unspecified order.

`get_cycle(G, V) -> Vertices | false`

Types:

G = digraph()

V1 = V2 = vertex()

Vertices = [vertex()]

If there is a *simple cycle* of length two or more through the vertex `V`, then the cycle is returned as a list `[V, ..., V]` of vertices, otherwise if there is a *loop* through `V`, then the loop is returned as a list `[V]`. If there are no cycles through `V`, then `false` is returned.

`get_path/3` is used for finding a simple cycle through `V`.

`get_path(G, V1, V2) -> Vertices | false`

Types:

G = digraph()

V1 = V2 = vertex()

Vertices = [vertex()]

Tries to find a *simple path* from the vertex `V1` to the vertex `V2` of the digraph `G`. Returns the path as a list `[V1, ..., V2]` of vertices, or `false` if no simple path from `V1` to `V2` of length one or more exists.

The digraph `G` is traversed in a depth-first manner, and the first path found is returned.

`get_short_cycle(G, V) -> Vertices | false`

Types:

G = digraph()

V1 = V2 = vertex()

Vertices = [vertex()]

Tries to find an as short as possible *simple cycle* through the vertex `V` of the digraph `G`. Returns the cycle as a list `[V, ..., V]` of vertices, or `false` if no simple cycle through `V` exists. Note that a *loop* through `V` is returned as the list `[V, V]`.

`get_short_path/3` is used for finding a simple cycle through `V`.

`get_short_path(G, V1, V2) -> Vertices | false`

Types:

G = digraph()

V1 = V2 = vertex()

Vertices = [vertex()]

Tries to find an as short as possible *simple path* from the vertex `V1` to the vertex `V2` of the digraph `G`. Returns the path as a list `[V1, ..., V2]` of vertices, or `false` if no simple path from `V1` to `V2` of length one or more exists.

The digraph `G` is traversed in a breadth-first manner, and the first path found is returned.

`in_degree(G, V) -> integer()`

Types:

G= digraph()

V = vertex()

Returns the *in-degree* of the vertex `V` of the digraph `G`.

`in_edges(G, V) -> Edges`

Types:

G = digraph()

V = vertex()

Edges = [edge()]

Returns a list of all edges *incident* on `V` of the digraph `G`, in some unspecified order.

`in_neighbours(G, V) -> Vertices`

Types:

G = digraph()

V = vertex()

Vertices = [vertex()]

Returns a list of all *in-neighbours* of `V` of the digraph `G`, in some unspecified order.

`info(G) -> InfoList`

Types:

G = digraph()

InfoList = [{cyclicity, Cyclicity}, {memory, NoWords}, {protection, Protection}]

Cyclicity = cyclic | acyclic

Protection = protected | private

NoWords = integer() >= 0

Returns a list of `{Tag, Value}` pairs describing the digraph `G`. The following pairs are returned:

- `{cyclicity, Cyclicity}`, where `Cyclicity` is `cyclic` or `acyclic`, according to the options given to `new`.
- `{memory, NoWords}`, where `NoWords` is the number of words allocated to the `ets` tables.
- `{protection, Protection}`, where `Protection` is `protected` or `private`, according to the options given to `new`.

`new() -> digraph()`

Equivalent to `new([])`.

`new(Type) -> digraph()`

Types:

Type = [cyclic | acyclic | private | protected]

Returns an *empty digraph* with properties according to the options in `Type`:

`cyclic`
: Allow *cycles* in the digraph (default).

`acyclic`
: The digraph is to be kept *acyclic*.

`protected`
: Other processes can read the digraph (default).

`private`
: The digraph can be read and modified by the creating process only.

If an unrecognized type option `T` is given or `Type` is not a proper list, there will be a `badarg` exception.

`no_edges(G) -> integer() >= 0`

Types:

G = digraph()

Returns the number of edges of the digraph `G`.

`no_vertices(G) -> integer() >= 0`

Types:

G = digraph()

Returns the number of vertices of the digraph G.

out_degree(G, V) -> integer()

Types:

G = digraph()

V = vertex()

Returns the *out-degree* of the vertex V of the digraph G.

out_edges(G, V) -> Edges

Types:

G = digraph()

V = vertex()

Edges = [edge()]

Returns a list of all edges *emanating* from V of the digraph G, in some unspecified order.

out_neighbours(G, V) -> Vertices

Types:

G = digraph()

V = vertex()

Vertices = [vertex()]

Returns a list of all *out-neighbours* of V of the digraph G, in some unspecified order.

vertex(G, V) -> {V, Label} | false

Types:

G = digraph()

V = vertex()

Label = label()

Returns {V, Label} where Label is the *label* of the vertex V of the digraph G, or false if there is no vertex V of the digraph G.

vertices(G) -> Vertices

Types:

G = digraph()

Vertices = [vertex()]

Returns a list of all vertices of the digraph G, in some unspecified order.

See Also

digraph_utils(3), ets(3)

digraph_utils

Erlang module

The `digraph_utils` module implements some algorithms based on depth-first traversal of directed graphs. See the `digraph` module for basic functions on directed graphs.

A *directed graph* (or just "digraph") is a pair (V, E) of a finite set V of *vertices* and a finite set E of *directed edges* (or just "edges"). The set of edges E is a subset of V × V (the Cartesian product of V with itself).

Digraphs can be annotated with additional information. Such information may be attached to the vertices and to the edges of the digraph. A digraph which has been annotated is called a *labeled digraph*, and the information attached to a vertex or an edge is called a *label*.

An edge e = (v, w) is said to *emanate* from vertex v and to be *incident* on vertex w. If there is an edge emanating from v and incident on w, then w is said to be an *out-neighbour* of v, and v is said to be an *in-neighbour* of w. A *path* P from v[1] to v[k] in a digraph (V, E) is a non-empty sequence v[1], v[2], ..., v[k] of vertices in V such that there is an edge (v[i],v[i+1]) in E for 1 <= i < k. The *length* of the path P is k-1. P is a *cycle* if the length of P is not zero and v[1] = v[k]. A *loop* is a cycle of length one. An *acyclic digraph* is a digraph that has no cycles.

A *depth-first traversal* of a directed digraph can be viewed as a process that visits all vertices of the digraph. Initially, all vertices are marked as unvisited. The traversal starts with an arbitrarily chosen vertex, which is marked as visited, and follows an edge to an unmarked vertex, marking that vertex. The search then proceeds from that vertex in the same fashion, until there is no edge leading to an unvisited vertex. At that point the process backtracks, and the traversal continues as long as there are unexamined edges. If there remain unvisited vertices when all edges from the first vertex have been examined, some hitherto unvisited vertex is chosen, and the process is repeated.

A *partial ordering* of a set S is a transitive, antisymmetric and reflexive relation between the objects of S. The problem of *topological sorting* is to find a total ordering of S that is a superset of the partial ordering. A digraph G = (V, E) is equivalent to a relation E on V (we neglect the fact that the version of directed graphs implemented in the `digraph` module allows multiple edges between vertices). If the digraph has no cycles of length two or more, then the reflexive and transitive closure of E is a partial ordering.

A *subgraph* G' of G is a digraph whose vertices and edges form subsets of the vertices and edges of G. G' is *maximal* with respect to a property P if all other subgraphs that include the vertices of G' do not have the property P. A *strongly connected component* is a maximal subgraph such that there is a path between each pair of vertices. A *connected component* is a maximal subgraph such that there is a path between each pair of vertices, considering all edges undirected. An *arborescence* is an acyclic digraph with a vertex V, the *root*, such that there is a unique path from V to every other vertex of G. A *tree* is an acyclic non-empty digraph such that there is a unique path between every pair of vertices, considering all edges undirected.

Exports

`arborescence_root(Digraph) -> no | {yes, Root}`

Types:

Digraph = digraph()

Root = vertex()

Returns `{yes, Root}` if `Root` is the *root* of the arborescence `Digraph`, `no` otherwise.

`components(Digraph) -> [Component]`

Types:

Digraph = digraph()

Component = [vertex()]

Returns a list of *connected components*. Each component is represented by its vertices. The order of the vertices and the order of the components are arbitrary. Each vertex of the digraph `Digraph` occurs in exactly one component.

`condensation(Digraph) -> CondensedDigraph`

Types:

Digraph = CondensedDigraph = digraph()

Creates a digraph where the vertices are the *strongly connected components* of `Digraph` as returned by `strong_components/1`. If X and Y are strongly connected components, and there exist vertices x and y in X and Y respectively such that there is an edge *emanating* from x and *incident* on y, then an edge emanating from X and incident on Y is created.

The created digraph has the same type as `Digraph`. All vertices and edges have the default *label* `[]`.

Each and every *cycle* is included in some strongly connected component, which implies that there always exists a *topological ordering* of the created digraph.

`cyclic_strong_components(Digraph) -> [StrongComponent]`

Types:

Digraph = digraph()

StrongComponent = [vertex()]

Returns a list of *strongly connected components*. Each strongly component is represented by its vertices. The order of the vertices and the order of the components are arbitrary. Only vertices that are included in some *cycle* in `Digraph` are returned, otherwise the returned list is equal to that returned by `strong_components/1`.

`is_acyclic(Digraph) -> bool()`

Types:

Digraph = digraph()

Returns `true` if and only if the digraph `Digraph` is *acyclic*.

`is_arborescence(Digraph) -> bool()`

Types:

Digraph = digraph()

Returns `true` if and only if the digraph `Digraph` is an *arborescence*.

`is_tree(Digraph) -> bool()`

Types:

Digraph = digraph()

Returns `true` if and only if the digraph `Digraph` is a *tree*.

`loop_vertices(Digraph) -> Vertices`

Types:

Digraph = digraph()

Vertices = [vertex()]

Returns a list of all vertices of `Digraph` that are included in some *loop*.

postorder(Digraph) -> Vertices

Types:

Digraph = digraph()

Vertices = [vertex()]

Returns all vertices of the digraph `Digraph`. The order is given by a *depth-first traversal* of the digraph, collecting visited vertices in postorder. More precisely, the vertices visited while searching from an arbitrarily chosen vertex are collected in postorder, and all those collected vertices are placed before the subsequently visited vertices.

preorder(Digraph) -> Vertices

Types:

Digraph = digraph()

Vertices = [vertex()]

Returns all vertices of the digraph `Digraph`. The order is given by a *depth-first traversal* of the digraph, collecting visited vertices in pre-order.

reachable(Vertices, Digraph) -> Vertices

Types:

Digraph = digraph()

Vertices = [vertex()]

Returns an unsorted list of digraph vertices such that for each vertex in the list, there is a *path* in `Digraph` from some vertex of `Vertices` to the vertex. In particular, since paths may have length zero, the vertices of `Vertices` are included in the returned list.

reachable_neighbours(Vertices, Digraph) -> Vertices

Types:

Digraph = digraph()

Vertices = [vertex()]

Returns an unsorted list of digraph vertices such that for each vertex in the list, there is a *path* in `Digraph` of length one or more from some vertex of `Vertices` to the vertex. As a consequence, only those vertices of `Vertices` that are included in some *cycle* are returned.

reaching(Vertices, Digraph) -> Vertices

Types:

Digraph = digraph()

Vertices = [vertex()]

Returns an unsorted list of digraph vertices such that for each vertex in the list, there is a *path* from the vertex to some vertex of `Vertices`. In particular, since paths may have length zero, the vertices of `Vertices` are included in the returned list.

reaching_neighbours(Vertices, Digraph) -> Vertices

Types:

Digraph = digraph()

Vertices = [vertex()]

Returns an unsorted list of digraph vertices such that for each vertex in the list, there is a *path* of length one or more from the vertex to some vertex of `Vertices`. As a consequence, only those vertices of `Vertices` that are included in some *cycle* are returned.

`strong_components(Digraph) -> [StrongComponent]`

Types:

Digraph = digraph()

StrongComponent = [vertex()]

Returns a list of *strongly connected components*. Each strongly component is represented by its vertices. The order of the vertices and the order of the components are arbitrary. Each vertex of the digraph `Digraph` occurs in exactly one strong component.

`subgraph(Digraph, Vertices [, Options]) -> Subgraph`

Types:

Digraph = Subgraph = digraph()

Options = [{type, SubgraphType}, {keep_labels, bool()}]

SubgraphType = inherit | type()

Vertices = [vertex()]

Creates a maximal *subgraph* of `Digraph` having as vertices those vertices of `Digraph` that are mentioned in `Vertices`.

If the value of the option `type` is `inherit`, which is the default, then the type of `Digraph` is used for the subgraph as well. Otherwise the option value of `type` is used as argument to `digraph:new/1`.

If the value of the option `keep_labels` is `true`, which is the default, then the *labels* of vertices and edges of `Digraph` are used for the subgraph as well. If the value is `false`, then the default label, `[]`, is used for the subgraph's vertices and edges.

`subgraph(Digraph, Vertices)` is equivalent to `subgraph(Digraph, Vertices, [])`.

There will be a `badarg` exception if any of the arguments are invalid.

`topsort(Digraph) -> Vertices | false`

Types:

Digraph = digraph()

Vertices = [vertex()]

Returns a *topological ordering* of the vertices of the digraph `Digraph` if such an ordering exists, `false` otherwise. For each vertex in the returned list, there are no *out-neighbours* that occur earlier in the list.

See Also

digraph(3)

epp

Erlang module

The Erlang code preprocessor includes functions which are used by `compile` to preprocess macros and include files before the actual parsing takes place.

Exports

`open(FileName, IncludePath) -> {ok,Epp} | {error, ErrorDescriptor}`

`open(FileName, IncludePath, PredefMacros) -> {ok,Epp} | {error, ErrorDescriptor}`

Types:

FileName = atom() | string()

IncludePath = [DirectoryName]

DirectoryName = atom() | string()

PredefMacros = [{atom(),term()}]

Epp = pid() -- handle to the epp server

ErrorDescriptor = term()

Opens a file for preprocessing.

`close(Epp) -> ok`

Types:

Epp = pid() -- handle to the epp server

Closes the preprocessing of a file.

`parse_erl_form(Epp) -> {ok, AbsForm} | {eof, Line} | {error, ErrorInfo}`

Types:

Epp = pid()

AbsForm = term()

Line = integer()

ErrorInfo = see separate description below.

Returns the next Erlang form from the opened Erlang source file. The tuple `{eof, Line}` is returned at end-of-file. The first form corresponds to an implicit attribute `-file(File,1).`, where `File` is the name of the file.

`parse_file(FileName,IncludePath,PredefMacro) -> {ok,[Form]} | {error,OpenError}`

Types:

FileName = atom() | string()

IncludePath = [DirectoryName]

DirectoryName = atom() | string()

PredefMacros = [{atom(),term()}]

Form = term() -- same as returned by erl_parse:parse_form

Preprocesses and parses an Erlang source file. Note that the tuple `{eof, Line}` returned at end-of-file is included as a "form".

Error Information

The `ErrorInfo` mentioned above is the standard `ErrorInfo` structure which is returned from all IO modules. It has the following format:

```
{ErrorLine, Module, ErrorDescriptor}
```

A string which describes the error is obtained with the following call:

```
Module:format_error(ErrorDescriptor)
```

See Also

erl_parse(3)

erl_eval

Erlang module

This module provides an interpreter for Erlang expressions. The expressions are in the abstract syntax as returned by `erl_parse`, the Erlang parser, or a call to `io:parse_erl_exprs/2`.

Exports

`exprs(Expressions, Bindings) -> {value, Value, NewBindings}`

`exprs(Expressions, Bindings, LocalFunctionHandler) -> {value, Value, NewBindings}`

`exprs(Expressions, Bindings, LocalFunctionHandler, NonlocalFunctionHandler) -> {value, Value, NewBindings}`

Types:

Expressions = as returned by erl_parse or io:parse_erl_exprs/2

Bindings = as returned by bindings/1

LocalFunctionHandler = {value, Func} | {eval, Func} | none

NonlocalFunctionHandler = {value, Func} | none

Evaluates `Expressions` with the set of bindings `Bindings`, where `Expressions` is a sequence of expressions (in abstract syntax) of a type which may be returned by `io:parse_erl_exprs/2`. See below for an explanation of how and when to use the arguments `LocalFunctionHandler` and `NonlocalFunctionHandler`.

Returns `{value, Value, NewBindings}`

`expr(Expression, Bindings) -> { value, Value, NewBindings }`

`expr(Expression, Bindings, LocalFunctionHandler) -> { value, Value, NewBindings }`

`expr(Expression, Bindings, LocalFunctionHandler, NonlocalFunctionHandler) -> { value, Value, NewBindings }`

`expr(Expression, Bindings, LocalFunctionHandler, NonlocalFunctionHandler, ReturnFormat) -> { value, Value, NewBindings } | Value`

Types:

Expression = as returned by io:parse_erl_form/2, for example

Bindings = as returned by bindings/1

LocalFunctionHandler = {value, Func} | {eval, Func} | none

NonlocalFunctionHandler = {value, Func} | none

ReturnFormat = value | none

Evaluates `Expression` with the set of bindings `Bindings`. `Expression` is an expression (in abstract syntax) of a type which may be returned by `io:parse_erl_form/2`. See below for an explanation of how and when to use the arguments `LocalFunctionHandler` and `NonlocalFunctionHandler`.

Returns `{value, Value, NewBindings}` by default. But if the `ReturnFormat` is `value` only the `Value` is returned.

`expr_list(ExpressionList, Bindings) -> {ValueList, NewBindings}`

`expr_list(ExpressionList, Bindings, LocalFunctionHandler) -> {ValueList, NewBindings}`

`expr_list(ExpressionList, Bindings, LocalFunctionHandler, NonlocalFunctionHandler) -> {ValueList, NewBindings}`

Evaluates a list of expressions in parallel, using the same initial bindings for each expression. Attempts are made to merge the bindings returned from each evaluation. This function is useful in the `LocalFunctionHandler`. See below.

Returns `{ValueList, NewBindings}`.

`new_bindings() -> BindingStruct`

Returns an empty binding structure.

`bindings(BindingStruct) -> Bindings`

Returns the list of bindings contained in the binding structure.

`binding(Name, BindingStruct) -> Binding`

Returns the binding of `Name` in `BindingStruct`.

`add_binding(Name, Value, Bindings) -> BindingStruct`

Adds the binding `Name = Value` to `Bindings`. Returns an updated binding structure.

`del_binding(Name, Bindings) -> BindingStruct`

Removes the binding of `Name` in `Bindings`. Returns an updated binding structure.

Local Function Handler

During evaluation of a function, no calls can be made to local functions. An undefined function error would be generated. However, the optional argument `LocalFunctionHandler` may be used to define a function which is called when there is a call to a local function. The argument can have the following formats:

`{value,Func}`

This defines a local function handler which is called with:

```
Func(Name, Arguments)
```

`Name` is the name of the local function (an atom) and `Arguments` is a list of the *evaluated* arguments. The function handler returns the value of the local function. In this case, it is not possible to access the current bindings. To signal an error, the function handler just calls `exit/1` with a suitable exit value.

`{eval,Func}`

This defines a local function handler which is called with:

```
Func(Name, Arguments, Bindings)
```

`Name` is the name of the local function (an atom), `Arguments` is a list of the *unevaluated* arguments, and `Bindings` are the current variable bindings. The function handler returns:

```
{value,Value,NewBindings}
```

`Value` is the value of the local function and `NewBindings` are the updated variable bindings. In this case, the function handler must itself evaluate all the function arguments and manage the bindings. To signal an error, the function handler just calls `exit/1` with a suitable exit value.

`none`

There is no local function handler.

Non-local Function Handler

The optional argument `NonlocalFunctionHandler` may be used to define a function which is called in the following cases: a functional object (fun) is called; a built-in function is called; a function is called using the M:F syntax, where M and F are atoms or expressions; an operator Op/A is called (this is handled as a call to the function `erlang:Op/A`). Exceptions are calls to `erlang:apply/2,3`; neither of the function handlers will be called for such calls. The argument can have the following formats:

`{value,Func}`

This defines an nonlocal function handler which is called with:

```
Func(FuncSpec, Arguments)
```

`FuncSpec` is the name of the function on the form `{Module,Function}` or a fun, and `Arguments` is a list of the *evaluated* arguments. The function handler returns the value of the function. To signal an error, the function handler just calls `exit/1` with a suitable exit value.

`none`

There is no nonlocal function handler.

> Note:
>
> For calls such as `erlang:apply(Fun, Args)` or `erlang:apply(Module, Function, Args)` the call of the non-local function handler corresponding to the call to `erlang:apply/2,3` itself--`Func({erlang, apply}, [Fun, Args])` or `Func({erlang, apply}, [Module, Function, Args])`--will never take place. The non-local function handler *will* however be called with the evaluated arguments of the call to `erlang:apply/2,3`: `Func(Fun, Args)` or `Func({Module, Function}, Args)` (assuming that `{Module, Function}` is not `{erlang, apply}`).
>
> Calls to functions defined by evaluating fun expressions `"fun ... end"` are also hidden from non-local function handlers.

The nonlocal function handler argument is probably not used as frequently as the local function handler argument. A possible use is to call `exit/1` on calls to functions that for some reason are not allowed to be called.

Bugs

The evaluator is not complete. `receive` cannot be handled properly.

Any undocumented functions in `erl_eval` should not be used.

erl_expand_records

Erlang module

Exports

`module(AbsForms, CompileOptions) -> AbsForms`

Types:

AbsForms = [term()]

CompileOptions = [term()]

Expands all records in a module. The returned module has no references to records, neither attributes nor code.

See Also

The *abstract format* documentation in ERTS User's Guide

erl_id_trans

Erlang module

This module performs an identity parse transformation of Erlang code. It is included as an example for users who may wish to write their own parse transformers. If the option `{parse_transform,Module}` is passed to the compiler, a user written function `parse_transform/2` is called by the compiler before the code is checked for errors.

Exports

`parse_transform(Forms, Options) -> Forms`

Types:

Forms = [erlang_form()]

Options = [compiler_options()]

Performs an identity transformation on Erlang forms, as an example.

Parse Transformations

Parse transformations are used if a programmer wants to use Erlang syntax, but with different semantics. The original Erlang code is then transformed into other Erlang code.

> ### Note:
>
> Programmers are strongly advised not to engage in parse transformations and no support is offered for problems encountered.

See Also

erl_parse(3), compile(3).

erl_internal

Erlang module

This module defines Erlang BIFs, guard tests and operators. This module is only of interest to programmers who manipulate Erlang code.

Exports

`bif(Name, Arity) -> bool()`

Types:

Name = **atom()**

Arity = **integer()**

Returns `true` if `Name/Arity` is an Erlang BIF which is automatically recognized by the compiler, otherwise `false`.

`guard_bif(Name, Arity) -> bool()`

Types:

Name = **atom()**

Arity = **integer()**

Returns `true` if `Name/Arity` is an Erlang BIF which is allowed in guards, otherwise `false`.

`type_test(Name, Arity) -> bool()`

Types:

Name = **atom()**

Arity = **integer()**

Returns `true` if `Name/Arity` is a valid Erlang type test, otherwise `false`.

`arith_op(OpName, Arity) -> bool()`

Types:

OpName = **atom()**

Arity = **integer()**

Returns `true` if `OpName/Arity` is an arithmetic operator, otherwise `false`.

`bool_op(OpName, Arity) -> bool()`

Types:

OpName = **atom()**

Arity = **integer()**

Returns `true` if `OpName/Arity` is a Boolean operator, otherwise `false`.

`comp_op(OpName, Arity) -> bool()`

Types:

OpName = **atom()**

Arity = **integer()**

Returns `true` if `OpName/Arity` is a comparison operator, otherwise `false`.

`list_op(OpName, Arity) -> bool()`

Types:

OpName = **atom()**

Arity = **integer()**

Returns `true` if `OpName/Arity` is a list operator, otherwise `false`.

`send_op(OpName, Arity) -> bool()`

Types:

OpName = **atom()**

Arity = **integer()**

Returns `true` if `OpName/Arity` is a send operator, otherwise `false`.

`op_type(OpName, Arity) -> Type`

Types:

OpName = **atom()**

Arity = **integer()**

Type = **arith | bool | comp | list | send**

Returns the `Type` of operator that `OpName/Arity` belongs to, or generates a `function_clause` error if it is not an operator at all.

erl_lint

Erlang module

This module is used to check Erlang code for illegal syntax and other bugs. It also warns against coding practices which are not recommended.

The errors detected include:

- redefined and undefined functions
- unbound and unsafe variables
- illegal record usage.

Warnings include:

- unused functions and imports
- unused variables
- variables imported into matches
- variables exported from `if`/`case`/`receive`
- variables shadowed in lambdas and list comprehensions.

Some of the warnings are optional, and can be turned on by giving the appropriate option, described below.

The functions in this module are invoked automatically by the Erlang compiler and there is no reason to invoke these functions separately unless you have written your own Erlang compiler.

Exports

```
module(AbsForms) -> {ok,Warnings} | {error,Errors,Warnings}

module(AbsForms, FileName) -> {ok,Warnings} | {error,Errors,Warnings}

module(AbsForms, FileName, CompileOptions) -> {ok,Warnings} |
{error,Errors,Warnings}
```

Types:

AbsForms = [term()]

FileName = FileName2 = atom() | string()

Warnings = Errors = [{Filename2,[ErrorInfo]}]

ErrorInfo = see separate description below.

CompileOptions = [term()]

This function checks all the forms in a module for errors. It returns:

`{ok,Warnings}`

There were no errors in the module.

`{error,Errors,Warnings}`

There were errors in the module.

Since this module is of interest only to the maintainers of the compiler, and to avoid having the same description in two places to avoid the usual maintenance nightmare, the elements of `Options` that control the warnings are only described in *compile(3)*.

The `AbsForms` of a module which comes from a file that is read through `epp`, the Erlang pre-processor, can come from many files. This means that any references to errors must include the file name (see *epp(3)*, or parser *erl_parse(3)*). The warnings and errors returned have the following format:

```
[{FileName2,[ErrorInfo]}]
```

The errors and warnings are listed in the order in which they are encountered in the forms. This means that the errors from one file may be split into different entries in the list of errors.

`is_guard_test(Expr) -> bool()`

Types:

Expr = term()

This function tests if `Expr` is a legal guard test. `Expr` is an Erlang term representing the abstract form for the expression. `erl_parse:parse_exprs(Tokens)` can be used to generate a list of `Expr`.

`format_error(ErrorDescriptor) -> Chars`

Types:

ErrorDescriptor = errordesc()

Chars = [char() | Chars]

Takes an `ErrorDescriptor` and returns a string which describes the error or warning. This function is usually called implicitly when processing an `ErrorInfo` structure (see below).

Error Information

The `ErrorInfo` mentioned above is the standard `ErrorInfo` structure which is returned from all IO modules. It has the following format:

```
{ErrorLine, Module, ErrorDescriptor}
```

A string which describes the error is obtained with the following call:

```
Module:format_error(ErrorDescriptor)
```

See Also

erl_parse(3), *epp(3)*

erl_parse

Erlang module

This module is the basic Erlang parser which converts tokens into the abstract form of either forms (i.e., top-level constructs), expressions, or terms. The Abstract Format is described in the ERTS User's Guide. Note that a token list must end with the *dot* token in order to be acceptable to the parse functions (see *erl_scan(3)*).

Exports

`parse_form(Tokens) -> {ok, AbsForm} | {error, ErrorInfo}`

Types:

Tokens = [Token]

Token = {Tag,Line} | {Tag,Line,term()}

Tag = atom()

AbsForm = term()

ErrorInfo = see section Error Information below.

This function parses `Tokens` as if it were a form. It returns:

`{ok, AbsForm}`

The parsing was successful. `AbsForm` is the abstract form of the parsed form.

`{error, ErrorInfo}`

An error occurred.

`parse_exprs(Tokens) -> {ok, Expr_list} | {error, ErrorInfo}`

Types:

Tokens = [Token]

Token = {Tag,Line} | {Tag,Line,term()}

Tag = atom()

Expr_list = [AbsExpr]

AbsExpr = term()

ErrorInfo = see section Error Information below.

This function parses `Tokens` as if it were a list of expressions. It returns:

`{ok, Expr_list}`

The parsing was successful. `Expr_list` is a list of the abstract forms of the parsed expressions.

`{error, ErrorInfo}`

An error occurred.

`parse_term(Tokens) -> {ok, Term} | {error, ErrorInfo}`

Types:

Tokens = [Token]

Token = {Tag,Line} | {Tag,Line,term()}

Tag = atom()

Term = term()

ErrorInfo = see section Error Information below.

This function parses `Tokens` as if it were a term. It returns:

`{ok, Term}`

The parsing was successful. `Term` is the Erlang term corresponding to the token list.

`{error, ErrorInfo}`

An error occurred.

`format_error(ErrorDescriptor) -> Chars`

Types:

ErrorDescriptor = errordesc()

Chars = [char() | Chars]

Uses an `ErrorDescriptor` and returns a string which describes the error. This function is usually called implicitly when an `ErrorInfo` structure is processed (see below).

`tokens(AbsTerm) -> Tokens`

`tokens(AbsTerm, MoreTokens) -> Tokens`

Types:

Tokens = MoreTokens = [Token]

Token = {Tag,Line} | {Tag,Line,term()}

Tag = atom()

AbsTerm = term()

ErrorInfo = see section Error Information below.

This function generates a list of tokens representing the abstract form `AbsTerm` of an expression. Optionally, it appends `Moretokens`.

`normalise(AbsTerm) -> Data`

Types:

AbsTerm = Data = term()

Converts the abstract form `AbsTerm` of a term into a conventional Erlang data structure (i.e., the term itself). This is the inverse of `abstract/1`.

`abstract(Data) -> AbsTerm`

Types:

Data = AbsTerm = term()

Converts the Erlang data structure `Data` into an abstract form of type `AbsTerm`. This is the inverse of `normalise/1`.

Error Information

The `ErrorInfo` mentioned above is the standard `ErrorInfo` structure which is returned from all IO modules. It has the format:

```
{ErrorLine, Module, ErrorDescriptor}
```

A string which describes the error is obtained with the following call:

```
Module:format_error(ErrorDescriptor)
```

See Also

io(3), *erl_scan(3)*, ERTS User's Guide

erl_pp

Erlang module

The functions in this module are used to generate aesthetically attractive representations of abstract forms, which are suitable for printing. All functions return (possibly deep) lists of characters and generate an error if the form is wrong.

All functions can have an optional argument which specifies a hook that is called if an attempt is made to print an unknown form.

Exports

`form(Form) -> DeepCharList`

`form(Form, HookFunction) -> DeepCharList`

Types:

Form = term()

HookFunction = see separate description below.

DeepCharList = [char()|DeepCharList]

Pretty prints a `Form` which is an abstract form of a type which is returned by `erl_parse:parse_form`.

`attribute(Attribute) -> DeepCharList`

`attribute(Attribute, HookFunction) -> DeepCharList`

Types:

Attribute = term()

HookFunction = see separate description below.

DeepCharList = [char()|DeepCharList]

The same as `form`, but only for the attribute `Attribute`.

`function(Function) -> DeepCharList`

`function(Function, HookFunction) -> DeepCharList`

Types:

Function = term()

HookFunction = see separate description below.

DeepCharList = [char()|DeepCharList]

The same as `form`, but only for the function `Function`.

`guard(Guard) -> DeepCharList`

`guard(Guard, HookFunction) -> DeepCharList`

Types:

Form = term()

HookFunction = see separate description below.

DeepCharList = [char()|DeepCharList]

The same as `form`, but only for the guard test `Guard`.

`exprs(Expressions) -> DeepCharList`

`exprs(Expressions, HookFunction) -> DeepCharList`

`exprs(Expressions, Indent, HookFunction) -> DeepCharList`

Types:

Expressions = term()

HookFunction = see separate description below.

Indent = integer()

DeepCharList = [char()|DeepCharList]

The same as `form`, but only for the sequence of expressions in `Expressions`.

`expr(Expression) -> DeepCharList`

`expr(Expression, HookFunction) -> DeepCharList`

`expr(Expression, Indent, HookFunction) -> DeepCharList`

`expr(Expression, Indent, Precedence, HookFunction) ->-> DeepCharList`

Types:

Expression = term()

HookFunction = see separate description below.

Indent = integer()

Precedence =

DeepCharList = [char()|DeepCharList]

This function prints one expression. It is useful for implementing hooks (see below).

Unknown Expression Hooks

The optional argument `HookFunction`, shown in the functions described above, defines a function which is called when an unknown form occurs where there should be a valid expression. It can have the following formats:

`Function`

The hook function is called by:

```
Function(Expr,
         CurrentIndentation,
         CurrentPrecedence,
         HookFunction)
```

`none`

There is no hook function

The called hook function should return a (possibly deep) list of characters. `expr/4` is useful in a hook.

If `CurrentIndentation` is negative, there will be no line breaks and only a space is used as a separator.

Bugs

It should be possible to have hook functions for unknown forms at places other than expressions.

See Also

io(3), erl_parse(3), erl_eval(3)

erl_scan

Erlang module

This module contains functions for tokenizing characters into Erlang tokens.

DATA TYPES

```
category() = atom()
column() = integer() > 0
line() = integer()
location() = line() | {line(), column()}
reserved_word_fun() -> fun(atom()) -> bool()
set_attribute_fun() -> fun(term()) -> term()
symbol() = atom() | float() | integer() | string()
token() = {category(), attributes()} | {category(), attributes(), symbol()}
attributes() = line() | list() | tuple()
```

Exports

```
1> erl_scan:string("a,c,{b,d}").
{ok,[{atom,1,a},
     {',',1},
     {atom,1,c},
     {',',1},
     {'{',1},
     {atom,1,b},
     {',',1},
     {atom,1,d},
     {'}',1}],
    1}
2> erl_scan:string("a,D,{b,d}").
{ok,[{atom,1,a},
     {',',1},
     {var,1,'D'},
     {',',1},
     {'{',1},
     {atom,1,b},
     {',',1},
     {atom,1,d},
     {'}',1}],
    1}
```

`string(String) -> Return`

`string(String, StartLocation) -> Return`

`string(String, StartLocation, Options) -> Return`

Types:

String = string()

Return = {ok, Tokens, EndLocation} | Error

Tokens = [token()]

Error = {error, ErrorInfo, EndLocation}

StartLocation = EndLocation = location()

Options = Option | [Option]

Option = {reserved_word_fun,reserved_word_fun()} | return_comments | return_white_spaces | return | text

Takes the list of characters `String` and tries to scan (tokenize) them. Returns `{ok, Tokens, EndLocation}`, where `Tokens` are the Erlang tokens from `String`. `EndLocation` is the first location after the last token.

`{error, ErrorInfo, EndLocation}` is returned if an error occurs. `EndLocation` is the first location after the erroneous token.

`string(String)` is equivalent to `string(String, 1)`, and `string(String, StartLocation)` is equivalent to `string(String, StartLocation, [])`.

`StartLocation` indicates the initial location when scanning starts. If `StartLocation` is a line `attributes()` as well as `EndLocation` and `ErrorLocation` will be lines. If `StartLocation` is a pair of a line and a column `attributes()` takes the form of an opaque compound data type, and `EndLocation` and `ErrorLocation` will be pairs of a line and a column. The *token attributes* contain information about the column and the line where the token begins, as well as the text of the token (if the `text` option is given), all of which can be accessed by calling *token_info/1,2* or *attributes_info/1,2*.

A *token* is a tuple containing information about syntactic category, the token attributes, and the actual terminal symbol. For punctuation characters (e.g. `;`, `|`) and reserved words, the category and the symbol coincide, and the token is represented by a two-tuple. Three-tuples have one of the following forms: `{atom, Info, atom()}`, `{char, Info, integer()}`, `{comment, Info, string()}`, `{float, Info, float()}`, `{integer, Info, integer()}`, `{var, Info, atom()}`, and `{white_space, Info, string()}`.

The valid options are:

`{reserved_word_fun, reserved_word_fun()}`

A callback function that is called when the scanner has found an unquoted atom. If the function returns `true`, the unquoted atom itself will be the category of the token; if the function returns `false`, `atom` will be the category of the unquoted atom.

`return_comments`

Return comment tokens.

`return_white_spaces`

Return white space tokens. By convention, if there is a newline character, it is always the first character of the text (there cannot be more than one newline in a white space token).

`return`

Short for `[return_comments, return_white_spaces]`.

`text`

Include the token's text in the token attributes. The text is the part of the input corresponding to the token.

`tokens(Continuation, CharSpec, StartLocation) -> Return`

`tokens(Continuation, CharSpec, StartLocation, Options) -> Return`

Types:

Continuation = [] | Continuation1

Return = {done, Result, LeftOverChars} | {more, Continuation1}

LeftOverChars = CharSpec

CharSpec = string() | eof

Continuation1 = tuple()

Result = {ok, Tokens, EndLocation} | {eof, EndLocation} | Error

Tokens = [token()]

Error = {error, ErrorInfo, EndLocation}

StartLocation = EndLocation = location()

Options = Option | [Option]

Option = {reserved_word_fun,reserved_word_fun()} | return_comments | return_white_spaces | return | text

This is the re-entrant scanner which scans characters until a *dot* ('.' followed by a white space) or `eof` has been reached. It returns:

`{done, Result, LeftOverChars}`

This return indicates that there is sufficient input data to get a result. `Result` is:

`{ok, Tokens, EndLocation}`

The scanning was successful. `Tokens` is the list of tokens including *dot*.

`{eof, EndLocation}`

End of file was encountered before any more tokens.

`{error, ErrorInfo, EndLocation}`

An error occurred. `LeftOverChars` is the remaining characters of the input data, starting from `EndLocation`.

`{more, Continuation1}`

More data is required for building a term. `Continuation1` must be passed in a new call to `tokens/3,4` when more data is available.

The `CharSpec` `eof` signals end of file. `LeftOverChars` will then take the value `eof` as well.

`tokens(Continuation, CharSpec, StartLocation)` is equivalent to `tokens(Continuation, CharSpec, StartLocation, [])`.

See *string/3* for a description of the various options.

`reserved_word(Atom) -> bool()`

Types:

Atom = atom()

Returns `true` if `Atom` is an Erlang reserved word, otherwise `false`.

`token_info(Token) -> TokenInfo`

Types:

Token = token()

TokenInfo = [TokenInfoTuple]

TokenInfoTuple = {TokenItem, Info}

TokenItem = atom()

Info = term()

Returns a list containing information about the token `Token`. The order of the `TokenInfoTuple`s is not defined. The following `TokenItem`s are returned: `category`, `column`, `length`, `line`, `symbol`, and `text`. See *token_info/2* for information about specific `TokenInfoTuple`s.

Note that if `token_info(Token, TokenItem)` returns `undefined` for some `TokenItem` in the list above, the item is not included in `TokenInfo`.

`token_info(Token, TokenItemSpec) -> TokenInfo`

Types:

Token = token()

TokenItemSpec = TokenItem | [TokenItem]

TokenInfo = TokenInfoTuple | undefined | [TokenInfoTuple]

TokenInfoTuple = {TokenItem, Info}

TokenItem = atom()

Info = term()

Returns a list containing information about the token `Token`. If `TokenItemSpec` is a single `TokenItem`, the returned value is the corresponding `TokenInfoTuple`, or `undefined` if the `TokenItem` has no value. If `TokenItemSpec` is a list of `TokenItem`, the result is a list of `TokenInfoTuple`. The `TokenInfoTuple`s

will appear with the corresponding `TokenItem`s in the same order as the `TokenItem`s appeared in the list of `TokenItem`s. `TokenItem`s with no value are not included in the list of `TokenInfoTuple`.

The following `TokenInfoTuple`s with corresponding `TokenItem`s are valid:

`{category, category()}`

The category of the token.

`{column, column()}`

The column where the token begins.

`{length, integer() > 0}`

The length of the token's text.

`{line, line()}`

The line where the token begins.

`{location, location()}`

The line and column where the token begins, or just the line if the column unknown.

`{symbol, symbol()}`

The token's symbol.

`{text, string()}`

The token's text.

`attributes_info(Attributes) -> AttributesInfo`

Types:

Attributes = attributes()

AttributesInfo = [AttributeInfoTuple]

AttributeInfoTuple = {AttributeItem, Info}

AttributeItem = atom()

Info = term()

Returns a list containing information about the token attributes `Attributes`. The order of the `AttributeInfoTuple`s is not defined. The following `AttributeItem`s are returned: `column`, `length`, `line`, and `text`. See *attributes_info/2* for information about specific `AttributeInfoTuple`s.

Note that if `attributes_info(Token, AttributeItem)` returns `undefined` for some `AttributeItem` in the list above, the item is not included in `AttributesInfo`.

`attributes_info(Attributes, AttributeItemSpec) -> AttributesInfo`

Types:

Attributes = attributes()

AttributeItemSpec = AttributeItem | [AttributeItem]

AttributesInfo = AttributeInfoTuple | undefined | [AttributeInfoTuple]

AttributeInfoTuple = {AttributeItem, Info}

AttributeItem = atom()

Info = term()

Returns a list containing information about the token attributes `Attributes`. If `AttributeItemSpec` is a single `AttributeItem`, the returned value is the corresponding `AttributeInfoTuple`, or `undefined` if the

`AttributeItem` has no value. If `AttributeItemSpec` is a list of `AttributeItem`, the result is a list of `AttributeInfoTuple`. The `AttributeInfoTuple`s will appear with the corresponding `AttributeItem`s in the same order as the `AttributeItem`s appeared in the list of `AttributeItem`s. `AttributeItem`s with no value are not included in the list of `AttributeInfoTuple`.

The following `AttributeInfoTuple`s with corresponding `AttributeItem`s are valid:

`{column, column()}`

The column where the token begins.

`{length, integer() > 0}`

The length of the token's text.

`{line, line()}`

The line where the token begins.

`{location, location()}`

The line and column where the token begins, or just the line if the column unknown.

`{text, string()}`

The token's text.

`set_attribute(AttributeItem, Attributes, SetAttributeFun) -> AttributesInfo`

Types:

AttributeItem = line

Attributes = attributes()

SetAttributeFun = set_attribute_fun()

Sets the value of the `line` attribute of the token attributes `Attributes`.

The `SetAttributeFun` is called with the value of the `line` attribute, and is to return the new value of the `line` attribute.

`format_error(ErrorDescriptor) -> string()`

Types:

ErrorDescriptor = errordesc()

Takes an `ErrorDescriptor` and returns a string which describes the error or warning. This function is usually called implicitly when processing an `ErrorInfo` structure (see below).

Error Information

The `ErrorInfo` mentioned above is the standard `ErrorInfo` structure which is returned from all IO modules. It has the following format:

```
{ErrorLocation, Module, ErrorDescriptor}
```

A string which describes the error is obtained with the following call:

```
Module:format_error(ErrorDescriptor)
```

Notes

The continuation of the first call to the re-entrant input functions must be `[]`. Refer to Armstrong, Virding and Williams, 'Concurrent Programming in Erlang', Chapter 13, for a complete description of how the re-entrant input scheme works.

See Also

io(3), erl_parse(3)

erl_tar

Erlang module

The `erl_tar` module archives and extract files to and from a tar file. The tar file format is the POSIX extended tar file format specified in IEEE Std 1003.1 and ISO/IEC 9945-1. That is the same format as used by `tar` program on Solaris, but is not the same as used by the GNU tar program.

By convention, the name of a tar file should end in "`.tar`". To abide to the convention, you'll need to add "`.tar`" yourself to the name.

Tar files can be created in one operation using the *create/2* or *create/3* function.

Alternatively, for more control, the *open*, *add/3,4*, and *close/1* functions can be used.

To extract all files from a tar file, use the *extract/1* function. To extract only some files or to be able to specify some more options, use the *extract/2* function.

To return a list of the files in a tar file, use either the *table/1* or *table/2* function. To print a list of files to the Erlang shell, use either the *t/1* or *tt/1* function.

To convert an error term returned from one of the functions above to a readable message, use the *format_error/1* function.

LIMITATIONS

For maximum compatibility, it is safe to archive files with names up to 100 characters in length. Such tar files can generally be extracted by any `tar` program.

If filenames exceed 100 characters in length, the resulting tar file can only be correctly extracted by a POSIX-compatible `tar` program (such as Solaris `tar`), not by GNU tar.

File have longer names than 256 bytes cannot be stored at all.

The filename of the file a symbolic link points is always limited to 100 characters.

Exports

`add(TarDescriptor, Filename, Options) -> RetValue`

Types:

TarDescriptor = term()

Filename = filename()

Options = [Option]

Option = dereference|verbose

RetValue = ok|{error,{Filename,Reason}}

Reason = term()

The `add/3` function adds a file to a tar file that has been opened for writing by *open/1*.

`dereference`

By default, symbolic links will be stored as symbolic links in the tar file. Use the `dereference` option to override the default and store the file that the symbolic link points to into the tar file.

`verbose`

Print an informational message about the file being added.

add(TarDescriptor, FilenameOrBin, NameInArchive, Options) -> RetValue

Types:

TarDescriptor = term()

FilenameOrBin = Filename()|binary()

Filename = filename()()

NameInArchive = filename()

Options = [Option]

Option = dereference|verbose

RetValue = ok|{error,{Filename,Reason}}

Reason = term()

The `add/4` function adds a file to a tar file that has been opened for writing by *open/1*. It accepts the same options as *add/3*.

`NameInArchive` is the name under which the file will be stored in the tar file. That is the name that the file will get when it will be extracted from the tar file.

close(TarDescriptor)

Types:

TarDescriptor = term()

The `close/1` function closes a tar file opened by *open/1*.

create(Name, FileList) ->RetValue

Types:

Name = filename()

FileList = [Filename|{NameInArchive, binary()},{NameInArchive, Filename}]

Filename = filename()

NameInArchive = filename()

RetValue = ok|{error,{Name,Reason}} <V>Reason = term()

The `create/2` function creates a tar file and archives the files whose names are given in `FileList` into it. The files may either be read from disk or given as binaries.

create(Name, FileList, OptionList)

Types:

Name = filename()

FileList = [Filename|{NameInArchive, binary()},{NameInArchive, Filename}]

Filename = filename()

NameInArchive = filename()

OptionList = [Option]

Option = compressed|cooked|dereference|verbose

RetValue = ok|{error,{Name,Reason}} <V>Reason = term()

The `create/3` function creates a tar file and archives the files whose names are given in `FileList` into it. The files may either be read from disk or given as binaries.

The options in `OptionList` modify the defaults as follows.

`compressed`

The entire tar file will be compressed, as if it has been run through the `gzip` program. To abide to the convention that a compressed tar file should end in "`.tar.gz`" or "`.tgz`", you'll need to add the appropriate extension yourself.

`cooked`

By default, the `open/2` function will open the tar file in `raw` mode, which is faster but does not allow a remote (erlang) file server to be used. Adding `cooked` to the mode list will override the default and open the tar file without the `raw` option.

`dereference`

By default, symbolic links will be stored as symbolic links in the tar file. Use the `dereference` option to override the default and store the file that the symbolic link points to into the tar file.

`verbose`

Print an informational message about each file being added.

`extract(Name) -> RetValue`

Types:

Name = filename()

RetValue = ok|{error,{Name,Reason}}

Reason = term()

The `extract/1` function extracts all files from a tar archive.

If the `Name` argument is given as "`{binary,Binary}`", the contents of the binary is assumed to be a tar archive.

If the `Name` argument is given as "`{file,Fd}`", `Fd` is assumed to be a file descriptor returned from the `file:open/2` function.

Otherwise, `Name` should be a filename.

`extract(Name, OptionList)`

Types:

Name = filename() | {binary,Binary} | {file,Fd}

Binary = binary()

Fd = file_descriptor()

OptionList = [Option]

Option = {cwd,Cwd}|{files,FileList}|keep_old_files|verbose|memory

Cwd = [dirname()]

FileList = [filename()]

RetValue = ok|MemoryRetValue|{error,{Name,Reason}}

MemoryRetValue = {ok, [{NameInArchive,binary()}]}

NameInArchive = filename()

Reason = term()

The `extract/2` function extracts files from a tar archive.

If the `Name` argument is given as "`{binary,Binary}`", the contents of the binary is assumed to be a tar archive.

If the `Name` argument is given as "`{file,Fd}`", `Fd` is assumed to be a file descriptor returned from the `file:open/2` function.

Otherwise, `Name` should be a filename.

The following options modify the defaults for the extraction as follows.

`{cwd,Cwd}`

Files with relative filenames will by default be extracted to the current working directory. Given the `{cwd,Cwd}` option, the `extract/2` function will extract into the directory `Cwd` instead of to the current working directory.

`{files,FileList}`

By default, all files will be extracted from the tar file. Given the `{files,Files}` option, the `extract/2` function will only extract the files whose names are included in `FileList`.

`compressed`

Given the `compressed` option, the `extract/2` function will uncompress the file while extracting If the tar file is not actually compressed, the `compressed` will effectively be ignored.

`cooked`

By default, the `open/2` function will open the tar file in `raw` mode, which is faster but does not allow a remote (erlang) file server to be used. Adding `cooked` to the mode list will override the default and open the tar file without the `raw` option.

`memory`

Instead of extracting to a directory, the memory option will give the result as a list of tuples {Filename, Binary}, where Binary is a binary containing the extracted data of the file named Filename in the tar file.

`keep_old_files`

By default, all existing files with the same name as file in the tar file will be overwritten Given the `keep_old_files` option, the `extract/2` function will not overwrite any existing files.

`verbose`

Print an informational message as each file is being extracted.

`format_error(Reason) -> string()`

Types:

Reason = term()

The `format_error/1` converts an error reason term to a human-readable error message string.

`open(Name, OpenModeList) -> RetValue`

Types:

Name = filename()

OpenModeList = [OpenMode]

Mode = write|compressed|cooked

RetValue = {ok,TarDescriptor}|{error,{Name,Reason}}

TarDescriptor = term()

Reason = term()

The `open/2` function creates a tar file for writing. (Any existing file with the same name will be truncated.)

By convention, the name of a tar file should end in "`.tar`". To abide to the convention, you'll need to add "`.tar`" yourself to the name.

Except for the `write` atom the following atoms may be added to `OpenModeList`:

`compressed`

The entire tar file will be compressed, as if it has been run through the `gzip` program. To abide to the convention that a compressed tar file should end in "`.tar.gz`" or "`.tgz`", you'll need to add the appropriate extension yourself.

`cooked`

By default, the `open/2` function will open the tar file in `raw` mode, which is faster but does not allow a remote (erlang) file server to be used. Adding `cooked` to the mode list will override the default and open the tar file without the `raw` option.

Use the *add/3,4* functions to add one file at the time into an opened tar file. When you are finished adding files, use the *close* function to close the tar file.

> Warning:
>
> The `TarDescriptor` term is not a file descriptor. You should not rely on the specific contents of the `TarDescriptor` term, as it may change in future versions as more features are added to the `erl_tar` module.

`table(Name) -> RetValue`

Types:

Name = filename()

RetValue = {ok,[string()]}|{error,{Name,Reason}}

Reason = term()

The `table/1` function retrieves the names of all files in the tar file `Name`.

`table(Name, Options)`

Types:

Name = filename()

The `table/2` function retrieves the names of all files in the tar file `Name`.

`t(Name)`

Types:

Name = filename()

The `t/1` function prints the names of all files in the tar file `Name` to the Erlang shell. (Similar to "`tar t`".)

`tt(Name)`

Types:

Name = filename()

The `tt/1` function prints names and information about all files in the tar file `Name` to the Erlang shell. (Similar to "`tar tv`".)

ets

Erlang module

数据的组织方式是一些列的动态表存储方式是元组.

读取时间

This module is an interface to the Erlang built-in term storage BIFs. These provide the ability to store very large quantities of data in an Erlang runtime system, and to have constant access time to the data. (In the case of `ordered_set`, see below, access time is proportional to the logarithm of the number of objects stored).

储存格式

Data is organized as a set of dynamic tables, which can store tuples. Each table is created by a process. When the process terminates, the table is automatically destroyed. Every table has access rights set at creation.

set与Key 是一对一的关系 bag是一对多的关系

Tables are divided into four different types, `set`, `ordered_set`, `bag` and `duplicate_bag`. A `set` or `ordered_set` table can only have one object associated with each key. A `bag` or `duplicate_bag` can have many objects associated with each key.

The number of tables stored at one Erlang node is limited. The current default limit is approximately 1400 tables. The upper limit can be increased by setting the environment variable `ERL_MAX_ETS_TABLES` before starting the Erlang runtime system (i.e. with the `-env` option to `erl`/`werl`). The actual limit may be slightly higher than the one specified, but never lower.

一个节点的ETS Table数量限制1400，可以调高

Note that there is no automatic garbage collection for tables. Even if there are no references to a table from any process, it will not automatically be destroyed unless the owner process terminates. It can be destroyed explicitly by using `delete/1`. The default owner is the process that created the table. Table ownership can be transferred at process termination by using the *heir* option or explicitly by calling *give_away/3*.

ETS表不在GC的管理范围内，除非拥有它的进程死掉它才会终止；可以通过delete删除数据 Table的拥有关系可以give_away 转交

Some implementation details:

- In the current implementation, every object insert and look-up operation results in a copy of the object.
- `'$end_of_table'` should not be used as a key since this atom is used to mark the end of the table when using `first`/`next`.

得到的值都是数据拷贝副本

Also worth noting is the subtle difference between *matching* and *comparing equal*, which is demonstrated by the different table types `set` and `ordered_set`. Two Erlang terms `match` if they are of the same type and have the same value, so that `1` matches `1`, but not `1.0` (as `1.0` is a `float()` and not an `integer()`). Two Erlang terms *compare equal* if they either are of the same type and value, or if both are numeric types and extend to the same value, so that `1` compares equal to both `1` and `1.0`. The `ordered_set` works on the *Erlang term order* and there is no defined order between an `integer()` and a `float()` that extends to the same value, hence the key `1` and the key `1.0` are regarded as equal in an `ordered_set` table.

这里可能是个误区，小心

匹配是值和类型相同就可以 compare equal类型相同 数值型可以认为都相同;ordered_set按照Erlang数据项排序规则进行排序,但是由于没有定义整形和浮点型之间的顺序所以1 和1.0 认为是相同的值.可以认为match是强条件匹配.

Failure

In general, the functions below will exit with reason `badarg` if any argument is of the wrong format, if the table identifier is invalid or if the operation is denied due to table access rights (*protected* or *private*).

Concurrency

并发控制：所有针对一个对象的更新都被保证是原子的、隔离的；修改要么全部成功要么失败。也没有其它的中间结果被其它的进程使用。有些方法可以在处理多个对象的时候保证这种原子性和隔离性。在数据库术语中隔离级别被称作序列化，就好像所有隔离的操作一个接一个严格按照顺序执行。

This module provides some limited support for concurrent access. All updates to single objects are guaranteed to be both *atomic* and *isolated*. This means that an updating operation towards a single object will either succeed or fail completely without any effect at all (atomicy). Nor can any intermediate results of the update be seen by other processes (isolation). Some functions that update several objects state that they even guarantee atomicy and isolation for the entire operation. In database terms the isolation level can be seen as "serializable", as if all isolated operations were carried out serially, one after the other in a strict order.

No other support is available within ETS that would guarantee consistency between objects. However, the `safe_fixtable/2` function can be used to guarantee that a sequence of `first/1` and `next/2` calls will traverse the table without errors and that each existing object in the table is visited exactly once, even if another process (or

除此之外没有其它的机制保证数据的一致性。

用到逐一遍历的场景就很少，使用safe_fixtable的情景就更少。不过这个机制是非常有用的，还记得之前.net版本中很麻烦的一件事情就是遍历在线玩家用户列表.由于玩家登录退出的变化,这里的异常几乎是不可避免的.

the same process) simultaneously deletes or inserts objects into the table. Nothing more is guaranteed: in particular objects that are inserted or deleted during such a traversal may be visited once or not at all. Functions that internally traverse over a table, like `select` and `match`, will give the same guarantee as `safe_fixtable`.

看到答案了吧:select和match的内部实现都是使用了safe_fixtable,所以会出现锁表的情况.

Match Specifications

Some of the functions uses a *match specification*, match_spec. A brief explanation is given in *select/2*. For a detailed description, see the chapter "Match specifications in Erlang" in *ERTS User's Guide*.

DATA TYPES

```
match_spec()
  a match specification, see above

tid()
  a table identifier, as returned by new/2
```

Exports

`all() -> [Tab]`

Types:

Tab = tid() | atom()

Returns a list of all tables at the node. Named tables are given by their names, unnamed tables are given by their table identifiers.

`delete(Tab) -> true`

1.获取所有ETS表
2.删除表
3.删除表中所有键值为Key的数据
4.清空表:原子 隔离

Types:

Tab = tid() | atom()

Deletes the entire table `Tab`.

`delete(Tab, Key) -> true`

Types:

Tab = tid() | atom()

Key = term()

Deletes all objects with the key `Key` from the table `Tab`.

`delete_all_objects(Tab) -> true`

Types:

Tab = tid() | atom()

Delete all objects in the ETS table `Tab`. The operation is guaranteed to be *atomic and isolated*.

`delete_object(Tab,Object) -> true`

Types:

Tab = tid() | atom()

Object = tuple()

从ETS中删除一条数据,通过数据本身去匹配,如果在bag的情况下,所有匹配的数据都会被删除.

Delete the exact object `Object` from the ETS table, leaving objects with the same key but other differences (useful for type `bag`). In a `duplicate_bag`, all instances of the object will be deleted.

`file2tab(Filename) -> {ok,Tab} | {error,Reason}`

Types:

Filename = string() | atom()

Tab = tid() | atom()

Reason = term()

Reads a file produced by *tab2file/2* or *tab2file/3* and creates the corresponding table `Tab`.

Equivalent to `file2tab(Filename,[])`.

`file2tab(Filename,Options) -> {ok,Tab} | {error,Reason}`

Types:

Filename = string() | atom()

Tab = tid() | atom()

Options = [Option]

Option = {verify, bool()}

Reason = term()

Reads a file produced by *tab2file/2* or *tab2file/3* and creates the corresponding table `Tab`.

The currently only supported option is `{verify,bool()}`. If verification is turned on (by means of specifying `{verify,true}`), the function utilizes whatever information is present in the file to assert that the information is not damaged. How this is done depends on which `extended_info` was written using *tab2file/3*.

If no `extended_info` is present in the file and `{verify,true}` is specified, the number of objects written is compared to the size of the original table when the dump was started. This might make verification fail if the table was `public` and objects were added or removed while the table was dumped to file. To avoid this type of problems, either do not verify files dumped while updated simultaneously or use the `{extended_info, [object_count]}` option to *tab2file/3*, which extends the information in the file with the number of objects actually written.

If verification is turned on and the file was written with the option `{extended_info, [md5sum]}`, reading the file is slower and consumes radically more CPU time than otherwise.

`{verify,false}` is the default.

`first(Tab) -> Key | '$end_of_table'`

Types:

Tab = tid() | atom()

Key = term()

Returns the first key `Key` in the table `Tab`. If the table is of the `ordered_set` type, the first key in Erlang term order will be returned. If the table is of any other type, the first key according to the table's internal order will be returned. If the table is empty, `'$end_of_table'` will be returned.

Use `next/2` to find subsequent keys in the table.

`foldl(Function, Acc0, Tab) -> Acc1`

针对表中的每一条数据都执行同样的操作

Types:

Function = fun(A, AccIn) -> AccOut

Tab = tid() | atom()

```
ets:foldl(fun(A,AC)-> io:format("Data:~p~n",[A]) end ,0,zen_ets).
```

Acc0 = Acc1 = AccIn = AccOut = term()

`Acc0` is returned if the table is empty. This function is similar to `lists:foldl/3`. The order in which the elements of the table are traversed is unspecified, except for tables of type `ordered_set`, for which they are traversed first to last.

这个表是没有锁表的 所以有可能导致数据不一致

If `Function` inserts objects into the table, or another process inserts objects into the table, those objects *may* (depending on key ordering) be included in the traversal.

`foldr(Function, Acc0, Tab) -> Acc1`

Types:

当我们要遍历ets的时候 可以用first/next来遍历 也可以用foldr foldl来看ets看成list来使用。
但是这样使用的时候有效率问题 数据要从erts内部搬到process 当ets很大的时候就效率低。

Function = fun(A, AccIn) -> AccOut

Tab = tid() | atom()

Acc0 = Acc1 = AccIn = AccOut = term()

`Acc0` is returned if the table is empty. This function is similar to `lists:foldr/3`. The order in which the elements of the table are traversed is unspecified, except for tables of type `ordered_set`, for which they are traversed last to first.

If `Function` inserts objects into the table, or another process inserts objects into the table, those objects *may* (depending on key ordering) be included in the traversal.

`from_dets(Tab, DetsTab) -> true`

Types:

Tab = tid() | atom()

DetsTab = atom()

Fills an already created ETS table with the objects in the already opened Dets table named `DetsTab`. The existing objects of the ETS table are kept unless overwritten.

Throws a badarg error if any of the tables does not exist or the dets table is not open.

`fun2ms(LiteralFun) -> MatchSpec`

把文本描述的规则转化成匹配的规范

Types:

LiteralFun -- see below

MatchSpec = match_spec()

Pseudo function that by means of a `parse_transform` translates `LiteralFun` typed as parameter in the function call to a *match_spec*. With "literal" is meant that the fun needs to textually be written as the parameter of the function, it cannot be held in a variable which in turn is passed to the function).

实现细节: The parse transform is implemented in the module `ms_transform` and the source *must* include the file `ms_transform.hrl` in `stdlib` for this pseudo function to work. Failing to include the hrl file in the source will result in a runtime error, not a compile time ditto. The include file is easiest included by adding the line `-include_lib("stdlib/include/ms_transform.hrl").` to the source file.

要求: The fun is very restricted, it can take only a single parameter (the object to match): a sole variable or a tuple. It needs to use the `is_XXX` guard tests. Language constructs that have no representation in a match_spec (like `if`, `case`, `receive` etc) are not allowed.

只有一个参数,需要对数据类型的guard检查 无法表达结构化的语义(不支持if case receive)

The return value is the resulting match_spec.

Example:

```
1> ets:fun2ms(fun({M,N}) when N > 3 -> M end).
[{{'$1','$2'},[{'>','$2',3}],['$1']}]
```

Variables from the environment can be imported, so that this works:

使用外部变量的情况:

```
2> X=3.
3
3> ets:fun2ms(fun({M,N}) when N > X -> M end).
[{{'$1','$2'},[{'>','$2',{const,3}}],['$1']}]
```

注意:引入的变量会一次性被替换编译到表达式中

The imported variables will be replaced by match_spec `const` expressions, which is consistent with the static scoping for Erlang funs. Local or global function calls can not be in the guard or body of the fun however. Calls to builtin match_spec functions of course is allowed: 除了BIF 本地或者全局的方法调用都不能出现在guard中

```
4> ets:fun2ms(fun({M,N}) when N > X, is_atomm(M) -> M end).
Error: fun containing local Erlang function calls
('is_atomm' called in guard) cannot be translated into match_spec
{error,transform_error}
5> ets:fun2ms(fun({M,N}) when N > X, is_atom(M) -> M end).
[{{'$1','$2'},[{'>','$2',{const,3}},{is_atom,'$1'}],['$1']}]
```

As can be seen by the example, the function can be called from the shell too. The fun needs to be literally in the call when used from the shell as well. Other means than the parse_transform are used in the shell case, but more or less the same restrictions apply (the exception being records, as they are not handled by the shell).

> Warning: parse_transform执行是其逻辑前提
>
> If the parse_transform is not applied to a module which calls this pseudo function, the call will fail in runtime (with a `badarg`). The module `ets` actually exports a function with this name, but it should never really be called except for when using the function in the shell. If the `parse_transform` is properly applied by including the `ms_transform.hrl` header file, compiled code will never call the function, but the function call is replaced by a literal match_spec.

For more information, see *ms_transform(3)*.

`give_away(Tab, Pid, GiftData) -> true`

Types: ETS拥有权移交

Tab = tid() | atom()

Pid = pid()

GiftData

GiftData = term()

Make process `Pid` the new owner of table `Tab`. If successful, the message `{'ETS-TRANSFER',Tab,FromPid,GiftData}` will be sent to the new owner.

The process `Pid` must be alive, local and not already the owner of the table. The calling process must be the table owner.

Note that `give_away` does not at all affect the *heir* option of the table. A table owner can for example set the `heir` to itself, give the table away and then get it back in case the receiver terminates.

1.查看ETS整体信息
2.查看具体的Table信息

i() -> ok

Displays information about all ETS tables on tty.

i(Tab) -> ok

Types:

Tab = tid() | atom()

Browses the table `Tab` on tty.

如果要查看ETS表的物理信息可以使用这个方法

info(Tab) -> [{Item, Value}] | undefined

Types:

Tab = tid() | atom()

Item = atom(), see below

Value = term(), see below

Returns information about the table `Tab` as a list of `{Item, Value}` tuples. If `Tab` has the correct type for a table identifier, but does not refer to an existing ETS table, `undefined` is returned. If `Tab` is not of the correct type, this function fails with reason `badarg`.

- `Item=memory, Value=int()`
  The number of words allocated to the table.
- `Item=owner, Value=pid()`
  The pid of the owner of the table.
- `Item=heir, Value=pid()|none`
  The pid of the heir of the table, or `none` if no heir is set.
- `Item=name, Value=atom()`
  The name of the table.
- `Item=size, Value=int()`
  The number of objects inserted in the table.
- `Item=node, Value=atom()`
  The node where the table is stored. This field is no longer meaningful as tables cannot be accessed from other nodes.
- `Item=named_table, Value=true|false`
  Indicates if the table is named or not.
- `Item=type, Value=set|ordered_set|bag|duplicate_bag`
  The table type.
- `Item=keypos, Value=int()`
  The key position.
- `Item=protection, Value=public|protected|private`
  The table access rights.
- `Item=compressed, Value=true|false`
  Indicates if the table is compressed or not.

什么时候会启用压缩呢?

查看单项信息

info(Tab, Item) -> Value | undefined

Types:

Tab = tid() | atom()

Item, Value - see below

Returns the information associated with `Item` for the table `Tab`, or returns `undefined` if `Tab` does not refer an existing ETS table. If `Tab` is not of the correct type, or if `Item` is not one of the allowed values, this function fails with reason `badarg`.

> **Warning:**
>
> In R11B and earlier, this function would not fail but return `undefined` for invalid values for `Item`.

In addition to the `{Item,Value}` pairs defined for `info/1`, the following items are allowed:

- `Item=fixed, Value=true|false`
  Indicates if the table is fixed by any process or not.
- `Item=safe_fixed, Value={FirstFixed,Info}|false`

  If the table has been fixed using `safe_fixtable/2`, the call returns a tuple where `FirstFixed` is the time when the table was first fixed by a process, which may or may not be one of the processes it is fixed by right now.

  `Info` is a possibly empty lists of tuples `{Pid,RefCount}`, one tuple for every process the table is fixed by right now. `RefCount` is the value of the reference counter, keeping track of how many times the table has been fixed by the process.

  If the table never has been fixed, the call returns `false`.

如果怀疑一个ETS表被锁住那么可以使用这个方法检测运行时状态
FristFixedTime是第一次被锁的时间,不一定是现在锁着的进程的时间.
那这个参数的意义就不大了,运行时我们更关注某一个表已经锁了多久了

Info参数是非常有用的,在运行期就要检查每一个ETS表,如果它们当前被锁着那么就可以见到这个表已经被当前的进程锁住多少次了,这样比较容易定位是哪个模块出了问题.

`init_table(Name, InitFun) -> true`

Types:

Name = atom()

InitFun = fun(Arg) -> Res

Arg = read | close

Res = end_of_input | {[object()], InitFun} | term()

Replaces the existing objects of the table `Tab` with objects created by calling the input function `InitFun`, see below. This function is provided for compatibility with the `dets` module, it is not more efficient than filling a table by using `ets:insert/2`.

这是提供了一个初始化入口和insert相比并没有性能上的提升

When called with the argument `read` the function `InitFun` is assumed to return `end_of_input` when there is no more input, or `{Objects, Fun}`, where `Objects` is a list of objects and `Fun` is a new input function. Any other value Value is returned as an error `{error, {init_fun, Value}}`. Each input function will be called exactly once, and should an error occur, the last function is called with the argument `close`, the reply of which is ignored.

If the type of the table is `set` and there is more than one object with a given key, one of the objects is chosen. This is not necessarily the last object with the given key in the sequence of objects returned by the input functions. This holds also for duplicated objects stored in tables of type `bag`.

`insert(Tab, ObjectOrObjects) -> true`

Types:

Tab = tid() | atom()

ObjectOrObjects = tuple() | [tuple()]

Inserts the object or all of the objects in the list `ObjectOrObjects` into the table `Tab`. If the table is a `set` and the key of the inserted objects *matches* the key of any object in the table, the old object will be replaced. If the table is an `ordered_set` and the key of the inserted object *compares equal* to the key of any object in the table, the old object is

如果表类型是set插入一条key已经存在的数据 行为相当于更新这条数据,ordered_set也是同样的处理只不过它比较方式是compares equal

also replaced. If the list contains more than one object with *matching* keys and the table is a `set`, one will be inserted, which one is not defined. The same thing holds for `ordered_set`, but will also happen if the keys *compare equal*.

The entire operation is guaranteed to be *atomic and isolated*, even when a list of objects is inserted.

`insert_new(Tab, ObjectOrObjects) -> bool()`

Types:

Tab = tid() | atom()

ObjectOrObjects = tuple() | [tuple()]

只要存在就不再进行 直接返回false

This function works exactly like `insert/2`, with the exception that instead of overwriting objects with the same key (in the case of `set` or `ordered_set`) or adding more objects with keys already existing in the table (in the case of `bag` and `duplicate_bag`), it simply returns `false`. If `ObjectOrObjects` is a list, the function checks *every* key prior to inserting anything. Nothing will be inserted if not *all* keys present in the list are absent from the table. Like `insert/2`, the entire operation is guaranteed to be *atomic and isolated*.

`is_compiled_ms(Term) -> bool()`

Types:

Term = term()

This function is used to check if a term is a valid compiled *match_spec*. The compiled match_spec is an opaque datatype which can *not* be sent between Erlang nodes nor be stored on disk. Any attempt to create an external representation of a compiled match_spec will result in an empty binary (<<>>). As an example, the following expression:

```
ets:is_compiled_ms(ets:match_spec_compile([{'_',[],[true]}])).
```

will yield `true`, while the following expressions:

```
MS = ets:match_spec_compile([{'_',[],[true]}]),
Broken = binary_to_term(term_to_binary(MS)),
ets:is_compiled_ms(Broken).
```

will yield false, as the variable `Broken` will contain a compiled match_spec that has passed through external representation.

> **Note:**
>
> The fact that compiled match_specs has no external representation is for performance reasons. It may be subject to change in future releases, while this interface will still remain for backward compatibility reasons.

`last(Tab) -> Key | '$end_of_table'`

Types:

Tab = tid() | atom()

Key = term()

Returns the last key `Key` according to Erlang term order in the table `Tab` of the `ordered_set` type. If the table is of any other type, the function is synonymous to `first/2`. If the table is empty, `'$end_of_table'` is returned.

对于没有顺序概念的表,直接返回first/2 这个真是!@#$%#%$

Use `prev/2` to find preceding keys in the table.

`lookup(Tab, Key) -> [Object]`

Types:

Tab = tid() | atom()

Key = term()

Object = tuple()

Returns a list of all objects with the key `Key` in the table `Tab`.

In the case of `set, bag and duplicate_bag`, an object is returned only if the given key *matches* the key of the object in the table. If the table is an `ordered_set` however, an object is returned if the key given *compares equal* to the key of an object in the table. The difference being the same as between `=:=` and `==`. As an example, one might insert an object with the `integer()1` as a key in an `ordered_set` and get the object returned as a result of doing a `lookup/2` with the `float()1.0` as the key to search for.

If the table is of type `set` or `ordered_set`, the function returns either the empty list or a list with one element, as there cannot be more than one object with the same key. If the table is of type `bag` or `duplicate_bag`, the function returns a list of arbitrary length.

Note that the time order of object insertions is preserved; The first object inserted with the given key will be first in the resulting list, and so on. 注意:如果允许重复,那么查询出来的结果是按照插入顺序排列

Insert and look-up times in tables of type `set`, `bag` and `duplicate_bag` are constant, regardless of the size of the table. For the `ordered_set` data-type, time is proportional to the (binary) logarithm of the number of objects.
这个结论非常重要:查询时间 插入时间对于set bag duplicate_bag都是常量值与表大小无关.

`lookup_element(Tab, Key, Pos) -> Elem`

Types:

Tab = tid() | atom()

Key = term()

Pos = int()

Elem = term() | [term()]

这个方法很纠结key还比较容易得到,Pos是指元素的位置,这个对于只提取一个字段值是非常有用的.

注意这里使用的是位置 也就是说依赖元素之间的相对位置 这个慎用

If the table `Tab` is of type `set` or `ordered_set`, the function returns the `Pos`:th element of the object with the key `Key`.

If the table is of type `bag` or `duplicate_bag`, the functions returns a list with the `Pos`:th element of every object with the key `Key`.

If no object with the key `Key` exists, the function will exit with reason `badarg`.

The difference between `set`, `bag` and `duplicate_bag` on one hand, and `ordered_set` on the other, regarding the fact that `ordered_set`'s view keys as equal when they *compare equal* whereas the other table types only regard them equal when they *match*, naturally holds for `lookup_element` as well as for `lookup`.
其它只有是不是相等的问题 没有大小前后的问题

`match(Tab, Pattern) -> [Match]`

Types:

Tab = tid() | atom()

Pattern = tuple()

Match = [term()]

Matches the objects in the table `Tab` against the pattern `Pattern`.

A pattern is a term that may contain:

- bound parts (Erlang terms),
- `'_'` which matches any Erlang term, and
- pattern variables: `'$N'` where `N`=0,1,... 占位符+新数据结构的相对位置

The function returns a list with one element for each matching object, where each element is an ordered list of pattern variable bindings. An example:

```
6> ets:match(T, '$1'). % Matches every object in the table
[[{rufsen,dog,7}],[{brunte,horse,5}],[{ludde,dog,5}]]
7> ets:match(T, {'_',dog,'$1'}).
[[7],[5]]
8> ets:match(T, {'_',cow,'$1'}).
[]
```

性能:如果能够指定Key 匹配效率是非常高的.如果只是变量和通配符,整个表内进行检索.如果表特别大,这个检索时间就非常可观了.

If the key is specified in the pattern, the match is very efficient. If the key is not specified, i.e. if it is a variable or an underscore, the entire table must be searched. The search time can be substantial if the table is very large.

On tables of the `ordered_set` type, the result is in the same order as in a `first/next` traversal.

`match(Tab, Pattern, Limit) -> {[Match],Continuation} | '$end_of_table'`

Types:

Tab = tid() | atom()

Pattern = tuple()

Match = [term()]

Continuation = term()

性能优化的途径:返回一定数量的匹配对象,返回值Continuation可以用来继续提取下一批匹配的值.这在空间上获得了优化,比遍历整个表还是要好很多. 对于只可能存在一条数据在ETS的查询都可以使用这种方式获得优化.

Works like `ets:match/2` but only returns a limited (`Limit`) number of matching objects. The `Continuation` term can then be used in subsequent calls to `ets:match/1` to get the next chunk of matching objects. This is a space efficient way to work on objects in a table which is still faster than traversing the table object by object using `ets:first/1` and `ets:next/1`.

`'$end_of_table'` is returned if the table is empty.

`match(Continuation) -> {[Match],Continuation} | '$end_of_table'`

Types:

Match = [term()]

Continuation = term()

Continues a match started with `ets:match/3`. The next chunk of the size given in the initial `ets:match/3` call is returned together with a new `Continuation` that can be used in subsequent calls to this function.

`'$end_of_table'` is returned when there are no more objects in the table.

`match_delete(Tab, Pattern) -> true`

Types:

Tab = tid() | atom()

Pattern = tuple()

Deletes all objects which match the pattern `Pattern` from the table `Tab`. See `match/2` for a description of patterns.

`match_object(Tab, Pattern) -> [Object]`

Types:

Tab = tid() | atom()

Pattern = Object = tuple()

Matches the objects in the table `Tab` against the pattern `Pattern`. See `match/2` for a description of patterns. The function returns a list of all objects which match the pattern.

If the key is specified in the pattern, the match is very efficient. If the key is not specified, i.e. if it is a variable or an underscore, the entire table must be searched. The search time can be substantial if the table is very large.

On tables of the `ordered_set` type, the result is in the same order as in a `first/next` traversal.

`match_object(Tab, Pattern, Limit) -> {[Match],Continuation} | '$end_of_table'`

Types:

Tab = tid() | atom()

Pattern = tuple()

Match = [term()]

Continuation = term()

Works like `ets:match_object/2` but only returns a limited (`Limit`) number of matching objects. The `Continuation` term can then be used in subsequent calls to `ets:match_object/1` to get the next chunk of matching objects. This is a space efficient way to work on objects in a table which is still faster than traversing the table object by object using `ets:first/1` and `ets:next/1`.

`'$end_of_table'` is returned if the table is empty.

`match_object(Continuation) -> {[Match],Continuation} | '$end_of_table'`

Types:

Match = [term()]

Continuation = term()

Continues a match started with `ets:match_object/3`. The next chunk of the size given in the initial `ets:match_object/3` call is returned together with a new `Continuation` that can be used in subsequent calls to this function.

`'$end_of_table'` is returned when there are no more objects in the table.

`match_spec_compile(MatchSpec) -> CompiledMatchSpec`

Types:

MatchSpec = match_spec()

CompiledMatchSpec = comp_match_spec()

This function transforms a *match_spec* into an internal representation that can be used in subsequent calls to `ets:match_spec_run/2`. The internal representation is opaque and can not be converted to external term format and then back again without losing its properties (meaning it can not be sent to a process on another node and still remain a valid compiled match_spec, nor can it be stored on disk). The validity of a compiled match_spec can be checked using `ets:is_compiled_ms/1`.

If the term `MatchSpec` can not be compiled (does not represent a valid match_spec), a `badarg` fault is thrown.

> Note:
>
> This function has limited use in normal code, it is used by Dets to perform the `dets:select` operations.

match_spec_run(List,CompiledMatchSpec) -> list()

Types:

List = [tuple()]

CompiledMatchSpec = comp_match_spec()

This function executes the matching specified in a compiled *match_spec* on a list of tuples. The `CompiledMatchSpec` term should be the result of a call to `ets:match_spec_compile/1` and is hence the internal representation of the match_spec one wants to use.

The matching will be executed on each element in `List` and the function returns a list containing all results. If an element in `List` does not match, nothing is returned for that element. The length of the result list is therefore equal or less than the the length of the parameter `List`. The two calls in the following example will give the same result (but certainly not the same execution time...):

```
Table = ets:new...
MatchSpec = ....
% The following call...
ets:match_spec_run(ets:tab2list(Table),
ets:match_spec_compile(MatchSpec)),
% ...will give the same result as the more common (and more efficient)
ets:select(Table,MatchSpec),
```

> Note:
>
> This function has limited use in normal code, it is used by Dets to perform the `dets:select` operations and by Mnesia during transactions.

member(Tab, Key) -> true | false

ContainsKey(Tab,Key)

Types:

Tab = tid() | atom()

Key = term()

Works like `lookup/2`, but does not return the objects. The function returns `true` if one or more elements in the table has the key `Key`, `false` otherwise.

new(Name, Options) -> tid() | atom()

Types:

Name = atom()

Options = [Option]

Option = Type | Access | named_table | {keypos,Pos} | {heir,pid(),HeirData} | {heir,none} | Tweaks

Type = set | ordered_set | bag | duplicate_bag

Access = public | protected | private

Tweaks = {write_concurrency,bool()} | {read_concurrency,bool()} | compressed

Pos = int()

HeirData = term()

Creates a new table and returns a table identifier which can be used in subsequent operations. The table identifier can be sent to other processes so that a table can be shared between different processes within a node.

The parameter `Options` is a list of atoms which specifies table type, access rights, key position and if the table is named or not. If one or more options are left out, the default values are used. This means that not specifying any options (`[]`) is the same as specifying `[set,protected,{keypos,1},{heir,none},{write_concurrency,false},{read_concurrency,false}]`.

- `set` The table is a `set` table - one key, one object, no order among objects. This is the default table type.
- `ordered_set` The table is a `ordered_set` table - one key, one object, ordered in Erlang term order, which is the order implied by the < and > operators. Tables of this type have a somewhat different behavior in some situations than tables of the other types. Most notably the `ordered_set` tables regard keys as equal when they *compare equal*, not only when they match. This means that to an `ordered_set`, the `integer()1` and the `float()1.0` are regarded as equal. This also means that the key used to lookup an element not necessarily *matches* the key in the elements returned, if `float()`'s and `integer()`'s are mixed in keys of a table.
- `bag` The table is a `bag` table which can have many objects, but only one instance of each object, per key.
- `duplicate_bag` The table is a `duplicate_bag` table which can have many objects, including multiple copies of the same object, per key.
- `public` Any process may read or write to the table.
- `protected` The owner process can read and write to the table. Other processes can only read the table. This is the default setting for the access rights.
- `private` Only the owner process can read or write to the table.
- `named_table` If this option is present, the name `Name` is associated with the table identifier. The name can then be used instead of the table identifier in subsequent operations.
- `{keypos,Pos}` Specfies which element in the stored tuples should be used as key. By default, it is the first element, i.e. `Pos=1`. However, this is not always appropriate. In particular, we do not want the first element to be the key if we want to store Erlang records in a table.

  Note that any tuple stored in the table must have at least `Pos` number of elements.
- `{heir,Pid,HeirData} | {heir,none}` 继承者:如果当前的主人死掉之后移交给谁
  Set a process as heir. The heir will inherit the table if the owner terminates. The message `{'ETS-TRANSFER',tid(),FromPid,HeirData}` will be sent to the heir when that happens. The heir must be a local process. Default heir is `none`, which will destroy the table when the owner terminates.
- `{write_concurrency,bool()}` Performance tuning. Default is `false`, in which case an operation that mutates (writes to) the table will obtain exclusive access, blocking any concurrent access of the same table until finished. If set to `true`, the table is optimized towards concurrent write access. Different objects of the same table can be mutated (and read) by concurrent processes. This is achieved to some degree at the expense of sequential access and concurrent reader performance. The `write_concurrency` option can be combined with the *read_concurrency* option. You typically want to combine these when large concurrent read bursts and large concurrent write bursts are common (see the documentation of the *read_concurrency* option for more information). Note that this option does not change any guarantees about *atomicy and isolation*. Functions that makes such promises over several objects (like `insert/2`) will gain less (or nothing) from this option.

  Table type `ordered_set` is not affected by this option in current implementation.
- `{read_concurrency,bool()}` Performance tuning. Default is `false`. When set to `true`, the table is optimized for concurrent read operations. When this option is enabled on a runtime system with SMP support, read operations become much cheaper; especially on systems with multiple physical processors. However, switching between read and write operations becomes more expensive. You typically want to enable this option when concurrent read operations are much more frequent than write operations, or when concurrent reads and writes comes in large read and write bursts (i.e., lots of reads not interrupted by writes, and lots of writes not interrupted by reads). You typically do *not* want to enable this option when the common access pattern is a few read operations interleaved with a few write operations repeatedly. In this case you will get a performance degradation by enabling this option. The `read_concurrency` option can be combined with the *write_concurrency* option.

写并发:false值如果写的行为会获得独占权限,阻塞其它的操作者对同一个表进行操作直到操作完成.如果设置为true表会针对并发写作特殊处理
表的不同对象还可以是对并发访问互斥的,

You typically want to combine these when large concurrent read bursts and large concurrent write bursts are common.

- `compressed` If this option is present, the table data will be stored in a more compact format to consume less memory. The downside is that it will make table operations slower. Especially operations that need to inspect entire objects, such as `match` and `select`, will get much slower. The key element is not compressed in current implementation.

肯定是拿空间换时间

`next(Tab, Key1) -> Key2 | '$end_of_table'`

Types:

Tab = tid() | atom()

Key1 = Key2 = term()

Returns the next key `Key2`, following the key `Key1` in the table `Tab`. If the table is of the `ordered_set` type, the next key in Erlang term order is returned. If the table is of any other type, the next key according to the table's internal order is returned. If there is no next key, `'$end_of_table'` is returned.

Use `first/1` to find the first key in the table.

Unless a table of type `set`, `bag` or `duplicate_bag` is protected using `safe_fixtable/2`, see below, a traversal may fail if concurrent updates are made to the table. If the table is of type `ordered_set`, the function returns the next key in order, even if the object does no longer exist.

`prev(Tab, Key1) -> Key2 | '$end_of_table'`

Types:

Tab = tid() | atom()

Key1 = Key2 = term()

Returns the previous key `Key2`, preceding the key `Key1` according the Erlang term order in the table `Tab` of the `ordered_set` type. If the table is of any other type, the function is synonymous to `next/2`. If there is no previous key, `'$end_of_table'` is returned.

这么怪的处理逻辑,非Ordered_set就不要这么用了

Use `last/1` to find the last key in the table.

`rename(Tab, Name) -> Name`

1.没有命名的表不受影响 相当于多了一种访问方式而已
2.已经命名的表就不能使用原有的名字进行访问
3.不要使用这个方法 太危险了 时机太难控制了

Types:

Tab = Name = atom()

Renames the named table `Tab` to the new name `Name`. Afterwards, the old name can not be used to access the table. Renaming an unnamed table has no effect.

`repair_continuation(Continuation, MatchSpec) -> Continuation`

Types:

Continuation = term()

MatchSpec = match_spec()

This function can be used to restore an opaque continuation returned by `ets:select/3` or `ets:select/1` if the continuation has passed through external term format (been sent between nodes or stored on disk).

The reason for this function is that continuation terms contain compiled match_specs and therefore will be invalidated if converted to external term format. Given that the original match_spec is kept intact, the continuation can be restored, meaning it can once again be used in subsequent `ets:select/1` calls even though it has been stored on disk or on another node.

As an example, the following sequence of calls will fail:

```
T=ets:new(x,[]),
...
{_,C} = ets:select(T,ets:fun2ms(fun({N,_}=A)
when (N rem 10) =:= 0 ->
A
end),10),
Broken = binary_to_term(term_to_binary(C)),
ets:select(Broken).
```

...while the following sequence will work:

```
T=ets:new(x,[]),
...
MS = ets:fun2ms(fun({N,_}=A)
when (N rem 10) =:= 0 ->
A
end),
{_,C} = ets:select(T,MS,10),
Broken = binary_to_term(term_to_binary(C)),
ets:select(ets:repair_continuation(Broken,MS)).
```

...as the call to `ets:repair_continuation/2` will reestablish the (deliberately) invalidated continuation `Broken`.

> ## Note:
>
> This function is very rarely needed in application code. It is used by Mnesia to implement distributed `select/3` and `select/1` sequences. A normal application would either use Mnesia or keep the continuation from being converted to external format.
>
> The reason for not having an external representation of a compiled match_spec is performance. It may be subject to change in future releases, while this interface will remain for backward compatibility.

`safe_fixtable(Tab, true|false) -> true`

Types:

Tab = tid() | atom()

Fixes a table of the `set`, `bag` or `duplicate_bag` table type for safe traversal.

A process fixes a table by calling `safe_fixtable(Tab,true)`. The table remains fixed until the process releases it by calling `safe_fixtable(Tab,false)`, or until the process terminates.

If several processes fix a table, the table will remain fixed until all processes have released it (or terminated). A reference counter is kept on a per process basis, and N consecutive fixes requires N releases to actually release the table.

When a table is fixed, a sequence of `first/1` and `next/2` calls are guaranteed to succeed and each object in the table will only be returned once, even if objects are removed or inserted during the traversal. The keys for new objects inserted during the traversal *may* be returned by *next/2* (it depends on the internal ordering of the keys). An example:

```
clean_all_with_value(Tab,X) ->
    safe_fixtable(Tab,true),
    clean_all_with_value(Tab,X,ets:first(Tab)),
    safe_fixtable(Tab,false).
```

```
clean_all_with_value(Tab,X,'$end_of_table') ->
    true;
clean_all_with_value(Tab,X,Key) ->
    case ets:lookup(Tab,Key) of
        [{Key,X}] ->
            ets:delete(Tab,Key);
        _ ->
            true
    end,
    clean_all_with_value(Tab,X,ets:next(Tab,Key)).
```

性能陷阱:

Note that no deleted objects are actually removed from a fixed table until it has been released. If a process fixes a table but never releases it, the memory used by the deleted objects will never be freed. The performance of operations on the table will also degrade significantly.

Use `info/2` to retrieve information about which processes have fixed which tables. A system with a lot of processes fixing tables may need a monitor which sends alarms when tables have been fixed for too long.

Note that for tables of the `ordered_set` type, `safe_fixtable/2` is not necessary as calls to `first/1` and `next/2` will always succeed.

`select(Tab, MatchSpec) -> [Match]`

Types:

Tab = tid() | atom()

Match = term()

MatchSpec = match_spec()

Matches the objects in the table `Tab` using a *match_spec*. This is a more general call than the `ets:match/2` and `ets:match_object/2` calls. In its simplest forms the match_specs look like this:

区别是什么??????

- MatchSpec = [MatchFunction]
- MatchFunction = {MatchHead, [Guard], [Result]}
- MatchHead = "Pattern as in ets:match"
- Guard = {"Guardtest name", ...}
- Result = "Term construct"

This means that the match_spec is always a list of one or more tuples (of arity 3). The tuples first element should be a pattern as described in the documentation of `ets:match/2`. The second element of the tuple should be a list of 0 or more guard tests (described below). The third element of the tuple should be a list containing a description of the value to actually return. In almost all normal cases the list contains exactly one term which fully describes the value to return for each object.

The return value is constructed using the "match variables" bound in the MatchHead or using the special match variables `'$_'` (the whole matching object) and `'$$'` (all match variables in a list), so that the following `ets:match/2` expression:

```
ets:match(Tab,{'$1','$2','$3'})
```

貌似match更简单

is exactly equivalent to:

```
ets:select(Tab,[{{'$1','$2','$3'},[],['$$']}])
```

- and the following `ets:match_object/2` call:

```
ets:match_object(Tab,{'$1','$2','$1'})
```

is exactly equivalent to

```
ets:select(Tab,[{{'$1','$2','$1'},[],['$_']}])
```

Composite terms can be constructed in the `Result` part either by simply writing a list, so that this code:

```
ets:select(Tab,[{{'$1','$2','$3'},[],['$$']}])
```

gives the same output as:

```
ets:select(Tab,[{{'$1','$2','$3'},[],[['$1','$2','$3']]}])
```

i.e. all the bound variables in the match head as a list. If tuples are to be constructed, one has to write a tuple of arity 1 with the single element in the tuple being the tuple one wants to construct (as an ordinary tuple could be mistaken for a `Guard`). Therefore the following call:

```
ets:select(Tab,[{{'$1','$2','$1'},[],['$_']}])
```

gives the same output as:

```
ets:select(Tab,[{{'$1','$2','$1'},[],[{{'$1','$2','$3'}}]}])
```

- this syntax is equivalent to the syntax used in the trace patterns (see *dbg(3)*).

The `Guard`s are constructed as tuples where the first element is the name of the test and the rest of the elements are the parameters of the test. To check for a specific type (say a list) of the element bound to the match variable `'$1'`, one would write the test as `{is_list, '$1'}`. If the test fails, the object in the table will not match and the next `MatchFunction` (if any) will be tried. Most guard tests present in Erlang can be used, but only the new versions prefixed `is_` are allowed (like `is_float`, `is_atom` etc).

The `Guard` section can also contain logic and arithmetic operations, which are written with the same syntax as the guard tests (prefix notation), so that a guard test written in Erlang looking like this:

```
is_integer(X), is_integer(Y), X + Y < 4711
```

is expressed like this (X replaced with '$1' and Y with '$2'):

```
[{is_integer, '$1'}, {is_integer, '$2'}, {'<', {'+', '$1', '$2'}, 4711}]
```

On tables of the `ordered_set` type, objects are visited in the same order as in a `first/next` traversal. This means that the match specification will be executed against objects with keys in the `first/next` order and the corresponding result list will be in the order of that execution.

`select(Tab, MatchSpec, Limit) -> {[Match],Continuation} | '$end_of_table'`

Types:

Tab = tid() | atom()

Match = term()

MatchSpec = match_spec()

Continuation = term()

Works like `ets:select/2` but only returns a limited (`Limit`) number of matching objects. The `Continuation` term can then be used in subsequent calls to `ets:select/1` to get the next chunk of matching objects. This is a space efficient way to work on objects in a table which is still faster than traversing the table object by object using `ets:first/1` and `ets:next/1`.

`'$end_of_table'` is returned if the table is empty.

`select(Continuation) -> {[Match],Continuation} | '$end_of_table'`

Types:

Match = term()

Continuation = term()

Continues a match started with `ets:select/3`. The next chunk of the size given in the initial `ets:select/3` call is returned together with a new `Continuation` that can be used in subsequent calls to this function.

`'$end_of_table'` is returned when there are no more objects in the table.

`select_count(Tab, MatchSpec) -> NumMatched`

这个恰恰是我需要的:只需要获得匹配的结果集数量 不需要获得详细数据

Types:

Tab = tid() | atom()

Object = tuple()

MatchSpec = match_spec()

NumMatched = integer()

Matches the objects in the table `Tab` using a *match_spec*. If the match_spec returns `true` for an object, that object considered a match and is counted. For any other result from the match_spec the object is not considered a match and is therefore not counted.

The function could be described as a `match_delete/2` that does not actually delete any elements, but only counts them.

The function returns the number of objects matched.

```
-record(test, {    id,    ip  }).

t() ->
   ets:new(test, [{keypos, #test.id}, named_table, public, set]),
   ets:insert(test, #test{ id = 1,ip = 1}),
   ets:insert(test, #test{id = 2, ip = 1}),

   MS_Filter = ets:fun2ms(fun(#test{id = ID, _='_'}) when ID == 1 -> true end),

   io:format("~p~n",[MS_Filter]),

    T0 = ets:select_count(test,MS_Filter),
    io:format("MS_Filter ~p~n",[T0]),

   T1 = ets:select_count(test,[{{test,1,'_'},[],[true]}]),
   io:format("~p~n",[T1]),
   T2 = ets:select(test, MS_Filter),
   io:format("~p~n",[T2]).
```

`select_delete(Tab, MatchSpec) -> NumDeleted`

Types:

Tab = tid() | atom()

Object = tuple()

MatchSpec = match_spec()

NumDeleted = integer()

Matches the objects in the table `Tab` using a *match_spec*. If the match_spec returns `true` for an object, that object is removed from the table. For any other result from the match_spec the object is retained. This is a more general call than the `ets:match_delete/2` call.

The function returns the number of objects actually deleted from the table.

> **Note:**
>
> The `match_spec` has to return the atom `true` if the object is to be deleted. No other return value will get the object deleted, why one can not use the same match specification for looking up elements as for deleting them.

`select_reverse(Tab, MatchSpec) -> [Match]`

Types:

设计规约:就不要在没有顺序概念的数据结构上去执行这种有顺序相关的方法了.没有意义.

Tab = tid() | atom()

Match = term()

MatchSpec = match_spec()

Works like `select/2`, but returns the list in reverse order for the `ordered_set` table type. For all other table types, the return value is identical to that of `select/2`.

`select_reverse(Tab, MatchSpec, Limit) -> {[Match],Continuation} | '$end_of_table'`

Types:

Tab = tid() | atom()

Match = term()

MatchSpec = match_spec()

Continuation = term()

Works like `select/3`, but for the `ordered_set` table type, traversing is done starting at the last object in Erlang term order and moves towards the first. For all other table types, the return value is identical to that of `select/3`.

Note that this is *not* equivalent to reversing the result list of a `select/3` call, as the result list is not only reversed, but also contains the last `Limit` matching objects in the table, not the first.

`select_reverse(Continuation) -> {[Match],Continuation} | '$end_of_table'`

Types:

Match = term()

Continuation = term()

Continues a match started with `ets:select_reverse/3`. If the table is an `ordered_set`, the traversal of the table will continue towards objects with keys earlier in the Erlang term order. The returned list will also contain objects with keys in reverse order.

For all other table types, the behaviour is exatly that of `select/1`.

Example:

```
1> T = ets:new(x,[ordered_set]).
2> [ ets:insert(T,{N}) || N <- lists:seq(1,10) ].
...
```

```
3> {R0,C0} = ets:select_reverse(T,[{'_',[],['$_']}],4).
...
4> R0.
[{10},{9},{8},{7}]
5> {R1,C1} = ets:select_reverse(C0).
...
6> R1.
[{6},{5},{4},{3}]
7> {R2,C2} = ets:select_reverse(C1).
...
8> R2.
[{2},{1}]
9> '$end_of_table' = ets:select_reverse(C2).
...
```

setopts(Tab, Opts) -> true 几乎用不到

Types:

Tab = tid() | atom()

Opts = Opt | [Opt]

Opt = {heir,pid(),HeirData} | {heir,none}

HeirData = term()

Set table options. The only option that currently is allowed to be set after the table has been created is *heir*. The calling process must be the table owner.

slot(Tab, I) -> [Object] | '$end_of_table'

Types: 调试使用的方法,不要在生产环境使用

Tab = tid() | atom()

I = int()

Object = tuple()

This function is mostly for debugging purposes. Normally one should use `first/next` or `last/prev` instead.

Returns all objects in the `I`:th slot of the table `Tab`. A table can be traversed by repeatedly calling the function, starting with the first slot `I=0` and ending when `'$end_of_table'` is returned. The function will fail with reason `badarg` if the `I` argument is out of range.

Unless a table of type `set`, `bag` or `duplicate_bag` is protected using `safe_fixtable/2`, see above, a traversal may fail if concurrent updates are made to the table. If the table is of type `ordered_set`, the function returns a list containing the `I`:th object in Erlang term order.

tab2file(Tab, Filename) -> ok | {error,Reason}

Types:

Tab = tid() | atom()

Filename = string() | atom()

Reason = term()

Dumps the table `Tab` to the file `Filename`.

Equivalent to `tab2file(Tab, Filename,[])`

tab2file(Tab, Filename, Options) -> ok | {error,Reason}

Types:

Tab = tid() | atom()

Filename = string() | atom()

Options = [Option]

Option = {extended_info, [ExtInfo]}

ExtInfo = object_count | md5sum

Reason = term()

Dumps the table `Tab` to the file `Filename`.

When dumping the table, certain information about the table is dumped to a header at the beginning of the dump. This information contains data about the table type, name, protection, size, version and if it's a named table. It also contains notes about what extended information is added to the file, which can be a count of the objects in the file or a MD5 sum of the header and records in the file.

The size field in the header might not correspond to the actual number of records in the file if the table is public and records are added or removed from the table during dumping. Public tables updated during dump, and that one wants to verify when reading, needs at least one field of extended information for the read verification process to be reliable later.

The `extended_info` option specifies what extra information is written to the table dump:

`object_count`

The number of objects actually written to the file is noted in the file footer, why verification of file truncation is possible even if the file was updated during dump.

`md5sum`

The header and objects in the file are checksummed using the built in MD5 functions. The MD5 sum of all objects is written in the file footer, so that verification while reading will detect the slightest bitflip in the file data. Using this costs a fair amount of CPU time.

Whenever the `extended_info` option is used, it results in a file not readable by versions of ets prior to that in stdlib-1.15.1

`tab2list(Tab) -> [Object]`

这个方法常用

Types:

Tab = tid() | atom()

Object = tuple()

Returns a list of all objects in the table `Tab`.

`tabfile_info(Filename) -> {ok, TableInfo} | {error, Reason}`

Types:

Filename = string() | atom()

TableInfo = [InfoItem]

InfoItem = {InfoTag, term()}

InfoTag = name | type | protection | named_table | keypos | size | extended_info | version

Reason = term()

Returns information about the table dumped to file by *tab2file/2* or *tab2file/3*

The following items are returned:

name

The name of the dumped table. If the table was a named table, a table with the same name cannot exist when the table is loaded from file with *file2tab/2*. If the table is not saved as a named table, this field has no significance at all when loading the table from file.

type

The ets type of the dumped table (i.e. `set`, `bag`, `duplicate_bag` or `ordered_set`). This type will be used when loading the table again.

protection

The protection of the dumped table (i.e. `private`, `protected` or `public`). A table loaded from the file will get the same protection.

named_table

`true` if the table was a named table when dumped to file, otherwise `false`. Note that when a named table is loaded from a file, there cannot exist a table in the system with the same name.

keypos

The `keypos` of the table dumped to file, which will be used when loading the table again.

size

The number of objects in the table when the table dump to file started, which in case of a `public` table need not correspond to the number of objects actually saved to the file, as objects might have been added or deleted by another process during table dump.

extended_info

The extended information written in the file footer to allow stronger verification during table loading from file, as specified to *tab2file/3*. Note that this function only tells *which* information is present, not the values in the file footer. The value is a list containing one or more of the atoms `object_count` and `md5sum`.

version

A tuple `{Major,Minor}` containing the major and minor version of the file format for ets table dumps. This version field was added beginning with stdlib-1.5.1, files dumped with older versions will return `{0,0}` in this field.

An error is returned if the file is inaccessible, badly damaged or not an file produced with *tab2file/2* or *tab2file/3*.

`table(Tab [, Options]) -> QueryHandle`

Types:

Tab = tid() | atom()

QueryHandle = - a query handle, see qlc(3) -

Options = [Option] | Option

Option = {n_objects, NObjects} | {traverse, TraverseMethod}

NObjects = default | integer() > 0

TraverseMethod = first_next | last_prev | select | {select, MatchSpec}

MatchSpec = match_spec()

Returns a QLC (Query List Comprehension) query handle. The module `qlc` implements a query language aimed mainly at Mnesia but ETS tables, Dets tables, and lists are also recognized by QLC as sources of data. Calling `ets:table/1,2` is the means to make the ETS table `Tab` usable to QLC.

When there are only simple restrictions on the key position QLC uses `ets:lookup/2` to look up the keys, but when that is not possible the whole table is traversed. The option `traverse` determines how this is done:

- `first_next`. The table is traversed one key at a time by calling `ets:first/1` and `ets:next/2`.
- `last_prev`. The table is traversed one key at a time by calling `ets:last/1` and `ets:prev/2`.
- `select`. The table is traversed by calling `ets:select/3` and `ets:select/1`. The option `n_objects` determines the number of objects returned (the third argument of `select/3`); the default is to return `100`

objects at a time. The *match_spec* (the second argument of `select/3`) is assembled by QLC: simple filters are translated into equivalent match_specs while more complicated filters have to be applied to all objects returned by `select/3` given a match_spec that matches all objects.

- `{select, MatchSpec}`. As for `select` the table is traversed by calling `ets:select/3` and `ets:select/1`. The difference is that the match_spec is explicitly given. This is how to state match_specs that cannot easily be expressed within the syntax provided by QLC.

The following example uses an explicit match_spec to traverse the table:

```
9> true = ets:insert(Tab = ets:new(t, []), [{1,a},{2,b},{3,c},{4,d}]),
MS = ets:fun2ms(fun({X,Y}) when (X > 1) or (X < 5) -> {Y} end),
QH1 = ets:table(Tab, [{traverse, {select, MS}}]).
```

An example with implicit match_spec:

```
10> QH2 = qlc:q([{Y} || {X,Y} <- ets:table(Tab), (X > 1) or (X < 5)]).
```

The latter example is in fact equivalent to the former which can be verified using the function `qlc:info/1`:

```
11> qlc:info(QH1) =:= qlc:info(QH2).
true
```

`qlc:info/1` returns information about a query handle, and in this case identical information is returned for the two query handles.

`test_ms(Tuple, MatchSpec) -> {ok, Result} | {error, Errors}`

Types:

Tuple = tuple()

MatchSpec = match_spec()

Result = term()

Errors = [{warning|error, string()}]

This function is a utility to test a *match_spec* used in calls to `ets:select/2`. The function both tests `MatchSpec` for "syntactic" correctness and runs the match_spec against the object `Tuple`. If the match_spec contains errors, the tuple `{error, Errors}` is returned where `Errors` is a list of natural language descriptions of what was wrong with the match_spec. If the match_spec is syntactically OK, the function returns `{ok,Term}` where `Term` is what would have been the result in a real `ets:select/2` call or `false` if the match_spec does not match the object `Tuple`.

This is a useful debugging and test tool, especially when writing complicated `ets:select/2` calls.

`to_dets(Tab, DetsTab) -> DetsTab`

Types:

Tab = tid() | atom()

DetsTab = atom()

Fills an already created/opened Dets table with the objects in the already opened ETS table named `Tab`. The Dets table is emptied before the objects are inserted.

update_counter(Tab, Key, UpdateOp) -> Result

update_counter(Tab, Key, [UpdateOp]) -> [Result]

update_counter(Tab, Key, Incr) -> Result

Types:

Tab = tid() | atom()

Key = term()

UpdateOp = {Pos,Incr} | {Pos,Incr,Threshold,SetValue}

Pos = Incr = Threshold = SetValue = Result = int()

This function provides an efficient way to update one or more counters, without the hassle of having to look up an object, update the object by incrementing an element and insert the resulting object into the table again. (The update is done atomically; i.e. no process can access the ets table in the middle of the operation.)

It will destructively update the object with key `Key` in the table `Tab` by adding `Incr` to the element at the `Pos`:th position. The new counter value is returned. If no position is specified, the element directly following the key (`<keypos>+1`) is updated.

If a `Threshold` is specified, the counter will be reset to the value `SetValue` if the following conditions occur:

- The `Incr` is not negative (`>= 0`) and the result would be greater than (`>`) `Threshold`
- The `Incr` is negative (`< 0`) and the result would be less than (`<`) `Threshold`

A list of `UpdateOp` can be supplied to do several update operations within the object. The operations are carried out in the order specified in the list. If the same counter position occurs more than one time in the list, the corresponding counter will thus be updated several times, each time based on the previous result. The return value is a list of the new counter values from each update operation in the same order as in the operation list. If an empty list is specified, nothing is updated and an empty list is returned. If the function should fail, no updates will be done at all.

The given Key is used to identify the object by either *matching* the key of an object in a `set` table, or *compare equal* to the key of an object in an `ordered_set` table (see *lookup/2* and *new/2* for details on the difference).

The function will fail with reason `badarg` if:

- the table is not of type `set` or `ordered_set`,
- no object with the right key exists,
- the object has the wrong arity,
- the element to update is not an integer,
- the element to update is also the key, or,
- any of `Pos`, `Incr`, `Threshold` or `SetValue` is not an integer

update_element(Tab, Key, {Pos,Value}) -> true | false

update_element(Tab, Key, [{Pos,Value}]) -> true | false

Types:

Tab = tid() | atom()

Key = Value = term()

Pos = int()

hassle :麻烦

This function provides an efficient way to update one or more elements within an object, without the hassle of having to look up, update and write back the entire object.

更新其中一个字段的值,省去了查询 更新 会写整个对象的麻烦

It will destructively update the object with key `Key` in the table `Tab`. The element at the `Pos`:th position will be given the value `Value`.

A list of `{Pos,Value}` can be supplied to update several elements within the same object. If the same position occurs more than one in the list, the last value in the list will be written. If the list is empty or the function fails, no updates will be done at all. The function is also atomic in the sense that other processes can never see any intermediate results.

The function returns `true` if an object with the key `Key` was found, `false` otherwise.

The given Key is used to identify the object by either *matching* the key of an object in a `set` table, or *compare equal* to the key of an object in an `ordered_set` table (see *lookup/2* and *new/2* for details on the difference).

The function will fail with reason `badarg` if:

- the table is not of type `set` or `ordered_set`,
- `Pos` is less than 1 or greater than the object arity, or,
- the element to update is also the key

The following functions used to be found in this module, but are now BIFs (i.e. implemented in C).

all/0
new/2
delete/1
delete/2
first/1
info/1
info/2
safe_fixtable/2
lookup/2
lookup_element/3
insert/2
is_compiled_ms/1
last/1
next/2
prev/2
rename/2
slot/2
match/1
match/2
match/3
match_object/1
match_object/2
match_object/3
match_spec_compile/1
match_spec_run_r/3
select/1
select/2
select/3
select_reverse/1
select_reverse/2
select_reverse/3
select_delete/2
update_counter/3

```
%'_' 通配符
A= ets:match(countries, {'$1','_','_' } ) ,
io:format(" ets:match(countries, {'$1','_','_' } ) Result : ~p ~n " ,[ A ]),
B= ets:match(countries , {'$1', '$0' ,'_' } ),
io:format(" ets:match(countries , {'$1', '$0' ,'_' } ), Result : ~p ~n " ,[ B ]),
C= ets:match(countries , {'$11', '$9' ,'_' } ),
io:format(" C= ets:match(countries , {'$11', '$9' ,'_' } ), Result : ~p ~n " ,[ C ]),
D= ets:match(countries , {'$11', '$99' ,'_' } ),
io:format(" ets:match(countries , {'$11', '$99' ,'_' } ), Result : ~p ~n " ,[ D ]),
E= ets:match(countries , {'$101', '$9' ,'_' } ),
io:format("ets:match(countries , {'$101', '$9' ,'_' } ), Result : ~p ~n " ,[ E ]),
F= ets:match(countries,{'$2',ireland,'_'}),
G= ets:match(countries,{'_',ireland,'_'}), % [[],[]] 如果没有数字占位符 是没有结果输出的 只是空列表
H= ets:match(countries,{'$2',cook,'_'}),
I= ets:match(countries,{'$0','$1',cook}),
J= ets:match(countries,{'$0','$0',cook}),
```

file_sorter

Erlang module

The functions of this module sort terms on files, merge already sorted files, and check files for sortedness. Chunks containing binary terms are read from a sequence of files, sorted internally in memory and written on temporary files, which are merged producing one sorted file as output. Merging is provided as an optimization; it is faster when the files are already sorted, but it always works to sort instead of merge.

On a file, a term is represented by a header and a binary. Two options define the format of terms on files:

- `{header, HeaderLength}`. HeaderLength determines the number of bytes preceding each binary and containing the length of the binary in bytes. Default is 4. The order of the header bytes is defined as follows: if `B` is a binary containing a header only, the size `Size` of the binary is calculated as `<<Size:HeaderLength/unit:8>> = B`.
- `{format, Format}`. The format determines the function that is applied to binaries in order to create the terms that will be sorted. The default value is `binary_term`, which is equivalent to `fun binary_to_term/1`. The value `binary` is equivalent to `fun(X) -> X end`, which means that the binaries will be sorted as they are. This is the fastest format. If `Format` is `term`, `io:read/2` is called to read terms. In that case only the default value of the `header` option is allowed. The `format` option also determines what is written to the sorted output file: if `Format` is `term` then `io:format/3` is called to write each term, otherwise the binary prefixed by a header is written. Note that the binary written is the same binary that was read; the results of applying the `Format` function are thrown away as soon as the terms have been sorted. Reading and writing terms using the `io` module is very much slower than reading and writing binaries.

Other options are:

- `{order, Order}`. The default is to sort terms in ascending order, but that can be changed by the value `descending` or by giving an ordering function `Fun`. An ordering function is antisymmetric, transitive and total. `Fun(A, B)` should return `true` if `A` comes before `B` in the ordering, `false` otherwise. An example of a typical ordering function is less than or equal to, `=</2`. Using an ordering function will slow down the sort considerably. The `keysort`, `keymerge` and `keycheck` functions do not accept ordering functions.
- `{unique, bool()}`. When sorting or merging files, only the first of a sequence of terms that compare equal (==) is output if this option is set to `true`. The default value is `false` which implies that all terms that compare equal are output. When checking files for sortedness, a check that no pair of consecutive terms compares equal is done if this option is set to `true`.
- `{tmpdir, TempDirectory}`. The directory where temporary files are put can be chosen explicitly. The default, implied by the value `""`, is to put temporary files on the same directory as the sorted output file. If output is a function (see below), the directory returned by `file:get_cwd()` is used instead. The names of temporary files are derived from the Erlang nodename (`node()`), the process identifier of the current Erlang emulator (`os:getpid()`), and a timestamp (`erlang:now()`); a typical name would be `fs_mynode@myhost_1763_1043_337000_266005.17`, where `17` is a sequence number. Existing files will be overwritten. Temporary files are deleted unless some uncaught EXIT signal occurs.
- `{compressed, bool()}`. Temporary files and the output file may be compressed. The default value `false` implies that written files are not compressed. Regardless of the value of the `compressed` option, compressed files can always be read. Note that reading and writing compressed files is significantly slower than reading and writing uncompressed files.
- `{size, Size}`. By default approximately 512*1024 bytes read from files are sorted internally. This option should rarely be needed.
- `{no_files, NoFiles}`. By default 16 files are merged at a time. This option should rarely be needed.

To summarize, here is the syntax of the options:

- `Options = [Option] | Option`
- `Option = {header, HeaderLength} | {format, Format} | {order, Order} | {unique, bool()} | {tmpdir, TempDirectory} | {compressed, bool()} | {size, Size} | {no_files, NoFiles}`
- `HeaderLength = int() > 0`
- `Format = binary_term | term | binary | FormatFun`
- `FormatFun = fun(Binary) -> Term`
- `Order = ascending | descending | OrderFun`
- `OrderFun = fun(Term, Term) -> bool()`
- `TempDirectory = "" | file_name()`
- `Size = int() >= 0`
- `NoFiles = int() > 1`

As an alternative to sorting files, a function of one argument can be given as input. When called with the argument `read` the function is assumed to return `end_of_input` or `{end_of_input, Value}` when there is no more input (`Value` is explained below), or `{Objects, Fun}`, where `Objects` is a list of binaries or terms depending on the format and `Fun` is a new input function. Any other value is immediately returned as value of the current call to `sort` or `keysort`. Each input function will be called exactly once, and should an error occur, the last function is called with the argument `close`, the reply of which is ignored.

A function of one argument can be given as output. The results of sorting or merging the input is collected in a non-empty sequence of variable length lists of binaries or terms depending on the format. The output function is called with one list at a time, and is assumed to return a new output function. Any other return value is immediately returned as value of the current call to the sort or merge function. Each output function is called exactly once. When some output function has been applied to all of the results or an error occurs, the last function is called with the argument `close`, and the reply is returned as value of the current call to the sort or merge function. If a function is given as input and the last input function returns `{end_of_input, Value}`, the function given as output will be called with the argument `{value, Value}`. This makes it easy to initiate the sequence of output functions with a value calculated by the input functions.

As an example, consider sorting the terms on a disk log file. A function that reads chunks from the disk log and returns a list of binaries is used as input. The results are collected in a list of terms.

```
sort(Log) ->
    {ok, _} = disk_log:open([{name,Log}, {mode,read_only}]),
    Input = input(Log, start),
    Output = output([]),
    Reply = file_sorter:sort(Input, Output, {format,term}),
    ok = disk_log:close(Log),
    Reply.

input(Log, Cont) ->
    fun(close) ->
            ok;
       (read) ->
            case disk_log:chunk(Log, Cont) of
                {error, Reason} ->
                    {error, Reason};
                {Cont2, Terms} ->
                    {Terms, input(Log, Cont2)};
                {Cont2, Terms, _Badbytes} ->
                    {Terms, input(Log, Cont2)};
                eof ->
                    end_of_input
            end
```

```
    end.

output(L) ->
    fun(close) ->
            lists:append(lists:reverse(L));
       (Terms) ->
            output([Terms | L])
    end.
```

Further examples of functions as input and output can be found at the end of the `file_sorter` module; the `term` format is implemented with functions.

The possible values of `Reason` returned when an error occurs are:

- `bad_object`, `{bad_object, FileName}`. Applying the format function failed for some binary, or the key(s) could not be extracted from some term.
- `{bad_term, FileName}`. `io:read/2` failed to read some term.
- `{file_error, FileName, Reason2}`. See `file(3)` for an explanation of `Reason2`.
- `{premature_eof, FileName}`. End-of-file was encountered inside some binary term.

Types

```
Binary = binary()
FileName = file_name()
FileNames = [FileName]
ICommand = read | close
IReply = end_of_input | {end_of_input, Value} | {[Object], Infun} | InputReply
Infun = fun(ICommand) -> IReply
Input = FileNames | Infun
InputReply = Term
KeyPos = int() > 0 | [int() > 0]
OCommand = {value, Value} | [Object] | close
OReply = Outfun | OutputReply
Object = Term | Binary
Outfun = fun(OCommand) -> OReply
Output = FileName | Outfun
OutputReply = Term
Term = term()
Value = Term
```

Exports

`sort(FileName) -> Reply`

`sort(Input, Output) -> Reply`

`sort(Input, Output, Options) -> Reply`

Types:

Reply = ok | {error, Reason} | InputReply | OutputReply

Sorts terms on files.

`sort(FileName)` is equivalent to `sort([FileName], FileName)`.

`sort(Input, Output)` is equivalent to `sort(Input, Output, [])`.

keysort(KeyPos, FileName) -> Reply

keysort(KeyPos, Input, Output) -> Reply

keysort(KeyPos, Input, Output, Options) -> Reply

Types:

Reply = ok | {error, Reason} | InputReply | OutputReply

Sorts tuples on files. The sort is performed on the element(s) mentioned in `KeyPos`. If two tuples compare equal (`==`) on one element, next element according to `KeyPos` is compared. The sort is stable.

`keysort(N, FileName)` is equivalent to `keysort(N, [FileName], FileName)`.

`keysort(N, Input, Output)` is equivalent to `keysort(N, Input, Output, [])`.

merge(FileNames, Output) -> Reply

merge(FileNames, Output, Options) -> Reply

Types:

Reply = ok | {error, Reason} | OutputReply

Merges terms on files. Each input file is assumed to be sorted.

`merge(FileNames, Output)` is equivalent to `merge(FileNames, Output, [])`.

keymerge(KeyPos, FileNames, Output) -> Reply

keymerge(KeyPos, FileNames, Output, Options) -> Reply

Types:

Reply = ok | {error, Reason} | OutputReply

Merges tuples on files. Each input file is assumed to be sorted on key(s).

`keymerge(KeyPos, FileNames, Output)` is equivalent to `keymerge(KeyPos, FileNames, Output, [])`.

check(FileName) -> Reply

check(FileNames, Options) -> Reply

Types:

Reply = {ok, [Result]} | {error, Reason}

Result = {FileName, TermPosition, Term}

TermPosition = int() > 1

Checks files for sortedness. If a file is not sorted, the first out-of-order element is returned. The first term on a file has position 1.

`check(FileName)` is equivalent to `check([FileName], [])`.

`keycheck(KeyPos, FileName) -> CheckReply`

`keycheck(KeyPos, FileNames, Options) -> Reply`

Types:

Reply = {ok, [Result]} | {error, Reason}

Result = {FileName, TermPosition, Term}

TermPosition = int() > 1

Checks files for sortedness. If a file is not sorted, the first out-of-order element is returned. The first term on a file has position 1.

`keycheck(KeyPos, FileName)` is equivalent to `keycheck(KeyPos, [FileName], [])`.

filelib

Erlang module

This module contains utilities on a higher level than the `file` module.

The module supports Unicode file names, so that it will match against regular expressions given in Unicode and that it will find and process raw file names (i.e. files named in a way that does not confirm to the expected encoding).

If the VM operates in Unicode file naming mode on a machine with transparent file naming, the `fun()` provided to `fold_files/5` needs to be prepared to handle binary file names.

For more information about raw file names, see the *file* module.

DATA TYPES

```
filename() =  = string() | atom() | DeepList | RawFilename
  DeepList = [char() | atom() | DeepList]
  RawFilename = binary()
  If VM is in unicode filename mode, string() and char() are allowed to be > 255.
  RawFilename is a filename not subject to Unicode translation, meaning that it
  can contain characters not conforming to the Unicode encoding expected from the
  filesystem (i.e. non-UTF-8 characters although the VM is started in Unicode
  filename mode).
dirname() = filename()
```

Exports

`ensure_dir(Name) -> ok | {error, Reason}`

Types:

Name = filename() | dirname()

Reason = posix() -- see file(3)

The `ensure_dir/1` function ensures that all parent directories for the given file or directory name `Name` exist, trying to create them if necessary.

Returns `ok` if all parent directories already exist or could be created, or `{error, Reason}` if some parent directory does not exist and could not be created for some reason.

`file_size(Filename) -> integer()`

The `file_size` function returns the size of the given file.

`fold_files(Dir, RegExp, Recursive, Fun, AccIn) -> AccOut`

Types:

Dir = dirname()

RegExp = regular_expression_string()

Recursive = true|false

Fun = fun(F, AccIn) -> AccOut

AccIn = AccOut = term()

The `fold_files/5` function folds the function `Fun` over all (regular) files `F` in the directory `Dir` that match the regular expression `RegExp` (see the *re* module for a description of the allowed regular expressions). If `Recursive` is true all sub-directories to `Dir` are processed. The regular expression matching is done on just the filename without the directory part.

If Unicode file name translation is in effect and the file system is completely transparent, file names that cannot be interpreted as Unicode may be encountered, in which case the `fun()` must be prepared to handle raw file names (i.e. binaries). If the regular expression contains codepoints beyond 255, it will not match file names that does not conform to the expected character encoding (i.e. are not encoded in valid UTF-8).

For more information about raw file names, see the *file* module.

`is_dir(Name) -> true | false`

Types:

Name = filename() | dirname()

The `is_dir/1` function returns `true` if `Name` refers to a directory, and `false` otherwise.

`is_file(Name) -> true | false`

Types:

Name = filename() | dirname()

The `is_file/1` function returns `true` if `Name` refers to a file or a directory, and `false` otherwise.

`is_regular(Name) -> true | false`

Types:

Name = filename()

The `is_regular/1` function returns `true` if `Name` refers to a file (regular file), and `false` otherwise.

`last_modified(Name) -> {{Year,Month,Day},{Hour,Min,Sec}} | 0`

Types:

Name = filename() | dirname()

The `last_modified/1` function returns the date and time the given file or directory was last modified, or 0 if the file does not exist.

`wildcard(Wildcard) -> list()`

Types:

Wildcard = filename() | dirname()

The `wildcard/1` function returns a list of all files that match Unix-style wildcard-string `Wildcard`.

The wildcard string looks like an ordinary filename, except that certain "wildcard characters" are interpreted in a special way. The following characters are special:

?

Matches one character.

*

Matches any number of characters up to the end of the filename, the next dot, or the next slash.

[Character1,Character2,...]

Matches any of the characters listed. Two characters separated by a hyphen will match a range of characters. Example: `[A-Z]` will match any uppercase letter.

{Item,...}

Alternation. Matches one of the alternatives.

Other characters represent themselves. Only filenames that have exactly the same character in the same position will match. (Matching is case-sensitive; i.e. "a" will not match "A").

Note that multiple "*" characters are allowed (as in Unix wildcards, but opposed to Windows/DOS wildcards).

Examples:

The following examples assume that the current directory is the top of an Erlang/OTP installation.

To find all `.beam` files in all applications, the following line can be used:

```
filelib:wildcard("lib/*/ebin/*.beam").
```

To find either `.erl` or `.hrl` in all applications `src` directories, the following

```
filelib:wildcard("lib/*/src/*.?rl")
```

or the following line

```
filelib:wildcard("lib/*/src/*.{erl,hrl}")
```

can be used.

To find all `.hrl` files in either `src` or `include` directories, use:

```
filelib:wildcard("lib/*/{src,include}/*.hrl").
```

To find all `.erl` or `.hrl` files in either `src` or `include` directories, use:

```
filelib:wildcard("lib/*/{src,include}/*.{erl,hrl}")
```

`wildcard(Wildcard, Cwd) -> list()`

Types:

Wildcard = filename() | dirname()

Cwd = dirname()

The `wildcard/2` function works like `wildcard/1`, except that instead of the actual working directory, `Cwd` will be used.

filename

Erlang module

The module `filename` provides a number of useful functions for analyzing and manipulating file names. These functions are designed so that the Erlang code can work on many different platforms with different formats for file names. With file name is meant all strings that can be used to denote a file. They can be short relative names like `foo.erl`, very long absolute name which include a drive designator and directory names like `D:\usr/local\bin\erl/lib\tools\foo.erl`, or any variations in between.

In Windows, all functions return file names with forward slashes only, even if the arguments contain back slashes. Use `join/1` to normalize a file name by removing redundant directory separators.

The module supports raw file names in the way that if a binary is present, or the file name cannot be interpreted according to the return value of *file:native_name_encoding/0*, a raw file name will also be returned. For example filename:join/1 provided with a path component being a binary (and also not being possible to interpret under the current native file name encoding) will result in a raw file name being returned (the join operation will have been performed of course). For more information about raw file names, see the *file* module.

DATA TYPES

```
name() = string() | atom() | DeepList | RawFilename
  DeepList = [char() | atom() | DeepList]
  RawFilename = binary()
  If VM is in unicode filename mode, string() and char() are allowed to be > 255.
  RawFilename is a filename not subject to Unicode translation, meaning that it
  can contain characters not conforming to the Unicode encoding expected from the
  filesystem (i.e. non-UTF-8 characters although the VM is started in Unicode
  filename mode).
```

Exports

`absname(Filename) -> string()`

Types:

Filename = name()

Converts a relative `Filename` and returns an absolute name. No attempt is made to create the shortest absolute name, because this can give incorrect results on file systems which allow links.

Unix examples:

```
1> pwd().
"/usr/local"
2> filename:absname("foo").
"/usr/local/foo"
3> filename:absname("../x").
"/usr/local/../x"
4> filename:absname("/").
"/"
```

Windows examples:

```
1> pwd().
"D:/usr/local"
2> filename:absname("foo").
"D:/usr/local/foo"
3> filename:absname("../x").
"D:/usr/local/../x"
4> filename:absname("/").
"D:/"
```

absname(Filename, Dir) -> string()

Types:

Filename = name()

Dir = string()

This function works like `absname/1`, except that the directory to which the file name should be made relative is given explicitly in the `Dir` argument.

absname_join(Dir, Filename) -> string()

Types:

Dir = string()

Filename = name()

Joins an absolute directory with a relative filename. Similar to `join/2`, but on platforms with tight restrictions on raw filename length and no support for symbolic links (read: VxWorks), leading parent directory components in `Filename` are matched against trailing directory components in `Dir` so they can be removed from the result - minimizing its length.

basename(Filename) -> string()

Types:

Filename = name()

Returns the last component of `Filename`, or `Filename` itself if it does not contain any directory separators.

```
5> filename:basename("foo").
"foo"
6> filename:basename("/usr/foo").
"foo"
7> filename:basename("/").
[]
```

basename(Filename, Ext) -> string()

Types:

Filename = Ext = name()

Returns the last component of `Filename` with the extension `Ext` stripped. This function should be used to remove a specific extension which might, or might not, be there. Use `rootname(basename(Filename))` to remove an extension that exists, but you are not sure which one it is.

```
8> filename:basename("~/src/kalle.erl", ".erl").
"kalle"
```

```
9> filename:basename("~/src/kalle.beam", ".erl").
"kalle.beam"
10> filename:basename("~/src/kalle.old.erl", ".erl").
"kalle.old"
11> filename:rootname(filename:basename("~/src/kalle.erl")).
"kalle"
12> filename:rootname(filename:basename("~/src/kalle.beam")).
"kalle"
```

dirname(Filename) -> string()

Types:

Filename = name()

Returns the directory part of `Filename`.

```
13> filename:dirname("/usr/src/kalle.erl").
"/usr/src"
14> filename:dirname("kalle.erl").
"."

5> filename:dirname("\\usr\\src/kalle.erl"). % Windows
"/usr/src"
```

extension(Filename) -> string()

Types: mochiweb_mime:from_extension(filename:extension(File))

Filename = name()

Returns the file extension of `Filename`, including the period. Returns an empty string if there is no extension.

```
15> filename:extension("foo.erl").
".erl"
16> filename:extension("beam.src/kalle").
[]
```

flatten(Filename) -> string()

Types:

Filename = name()

Converts a possibly deep list filename consisting of characters and atoms into the corresponding flat string filename.

join(Components) -> string()

Types:

Components = [string()]

Joins a list of file name `Components` with directory separators. If one of the elements of `Components` includes an absolute path, for example `"/xxx"`, the preceding elements, if any, are removed from the result.

The result is "normalized":

- Redundant directory separators are removed.
- In Windows, all directory separators are forward slashes and the drive letter is in lower case.

```
17> filename:join(["/usr", "local", "bin"]).
"/usr/local/bin"
18> filename:join(["a/b///c/"]).
"a/b/c"

6> filename:join(["B:a\\b///c/"]). % Windows
"b:a/b/c"
```

join(Name1, Name2) -> string()

Types:

Name1 = Name2 = string()

Joins two file name components with directory separators. Equivalent to `join([Name1, Name2])`.

nativename(Path) -> string()

Types:

Path = string()

Converts `Path` to a form accepted by the command shell and native applications on the current platform. On Windows, forward slashes is converted to backward slashes. On all platforms, the name is normalized as done by `join/1`.

```
19> filename:nativename("/usr/local/bin/"). % Unix
"/usr/local/bin"

7> filename:nativename("/usr/local/bin/"). % Windows
"\\usr\\local\\bin"
```

pathtype(Path) -> absolute | relative | volumerelative

Returns the type of path, one of `absolute`, `relative`, or `volumerelative`.

`absolute`

The path name refers to a specific file on a specific volume.

Unix example: `/usr/local/bin`

Windows example: `D:/usr/local/bin`

`relative`

The path name is relative to the current working directory on the current volume.

Example: `foo/bar, ../src`

`volumerelative`

The path name is relative to the current working directory on a specified volume, or it is a specific file on the current working volume.

Windows example: `D:bar.erl, /bar/foo.erl`

rootname(Filename) -> string()

rootname(Filename, Ext) -> string()

Types:

Filename = Ext = name()

Remove a filename extension. `rootname/2` works as `rootname/1`, except that the extension is removed only if it is `Ext`.

```
20> filename:rootname("/beam.src/kalle").
/beam.src/kalle"
21> filename:rootname("/beam.src/foo.erl").
"/beam.src/foo"
22> filename:rootname("/beam.src/foo.erl", ".erl").
"/beam.src/foo"
23> filename:rootname("/beam.src/foo.beam", ".erl").
"/beam.src/foo.beam"
```

split(Filename) -> Components

Types:

Filename = name()

Components = [string()]

Returns a list whose elements are the path components of `Filename`.

```
24> filename:split("/usr/local/bin").
["/","usr","local","bin"]
25> filename:split("foo/bar").
["foo","bar"]
26> filename:split("a:\\msdev\\include").
["a:/","msdev","include"]
```

find_src(Beam) -> {SourceFile, Options} | {error,{ErrorReason,Module}}

find_src(Beam, Rules) -> {SourceFile, Options} | {error,{ErrorReason,Module}}

Types:

Beam = Module | Filename

Module = atom()

Filename = string() | atom()

SourceFile = string()

Options = [Opt]

Opt = {i, string()} | {outdir, string()} | {d, atom()}

ErrorReason = non_existing | preloaded | interpreted

Finds the source filename and compiler options for a module. The result can be fed to `compile:file/2` in order to compile the file again.

The `Beam` argument, which can be a string or an atom, specifies either the module name or the path to the source code, with or without the `".erl"` extension. In either case, the module must be known by the code server, i.e. `code:which(Module)` must succeed.

`Rules` describes how the source directory can be found, when the object code directory is known. It is a list of tuples `{BinSuffix, SourceSuffix}` and is interpreted as follows: If the end of the directory name where the object is located matches `BinSuffix`, then the source code directory has the same name, but with `BinSuffix` replaced by `SourceSuffix`. `Rules` defaults to:

```
[{"", ""}, {"ebin", "src"}, {"ebin", "esrc"}]
```

If the source file is found in the resulting directory, then the function returns that location together with `Options`. Otherwise, the next rule is tried, and so on.

The function returns `{SourceFile, Options}` if it succeeds. `SourceFile` is the absolute path to the source file without the `".erl"` extension. `Options` include the options which are necessary to recompile the file with `compile:file/2`, but excludes options such as `report` or `verbose` which do not change the way code is generated. The paths in the `{outdir, Path}` and `{i, Path}` options are guaranteed to be absolute.

gb_sets

Erlang module

An implementation of ordered sets using Prof. Arne Andersson's General Balanced Trees. This can be much more efficient than using ordered lists, for larger sets, but depends on the application.

This module considers two elements as different if and only if they do not compare equal (==).

Complexity note

The complexity on set operations is bounded by either O(|S|) or O(|T| * log(|S|)), where S is the largest given set, depending on which is fastest for any particular function call. For operating on sets of almost equal size, this implementation is about 3 times slower than using ordered-list sets directly. For sets of very different sizes, however, this solution can be arbitrarily much faster; in practical cases, often between 10 and 100 times. This implementation is particularly suited for accumulating elements a few at a time, building up a large set (more than 100-200 elements), and repeatedly testing for membership in the current set.

As with normal tree structures, lookup (membership testing), insertion and deletion have logarithmic complexity.

Compatibility

All of the following functions in this module also exist and do the same thing in the `sets` and `ordsets` modules. That is, by only changing the module name for each call, you can try out different set representations.

- `add_element/2`
- `del_element/2`
- `filter/2`
- `fold/3`
- `from_list/1`
- `intersection/1`
- `intersection/2`
- `is_element/2`
- `is_set/1`
- `is_subset/2`
- `new/0`
- `size/1`
- `subtract/2`
- `to_list/1`
- `union/1`
- `union/2`

DATA TYPES

```
gb_set() = a GB set
```

Exports

`add(Element, Set1) -> Set2`

`add_element(Element, Set1) -> Set2`

Types:

Element = **term()**

Set1 = **Set2** = **gb_set()**

Returns a new gb_set formed from `Set1` with `Element` inserted. If `Element` is already an element in `Set1`, nothing is changed.

`balance(Set1) -> Set2`

Types:

Set1 = **Set2** = **gb_set()**

Rebalances the tree representation of `Set1`. Note that this is rarely necessary, but may be motivated when a large number of elements have been deleted from the tree without further insertions. Rebalancing could then be forced in order to minimise lookup times, since deletion only does not rebalance the tree.

`delete(Element, Set1) -> Set2`

Types:

Element = **term()**

Set1 = **Set2** = **gb_set()**

Returns a new gb_set formed from `Set1` with `Element` removed. Assumes that `Element` is present in `Set1`.

`delete_any(Element, Set1) -> Set2`

`del_element(Element, Set1) -> Set2`

Types:

Element = **term()**

Set1 = **Set2** = **gb_set()**

Returns a new gb_set formed from `Set1` with `Element` removed. If `Element` is not an element in `Set1`, nothing is changed.

`difference(Set1, Set2) -> Set3`

`subtract(Set1, Set2) -> Set3`

Types:

Set1 = **Set2** = **Set3** = **gb_set()**

Returns only the elements of `Set1` which are not also elements of `Set2`.

`empty() -> Set`

`new() -> Set`

Types:

Set = gb_set()

Returns a new empty gb_set.

filter(Pred, Set1) -> Set2

Types:

Pred = fun (E) -> bool()
E = term()
Set1 = Set2 = gb_set()

Filters elements in `Set1` using predicate function `Pred`.

fold(Function, Acc0, Set) -> Acc1

Types:

Function = fun (E, AccIn) -> AccOut
Acc0 = Acc1 = AccIn = AccOut = term()
E = term()
Set = gb_set()

Folds `Function` over every element in `Set` returning the final value of the accumulator.

from_list(List) -> Set

Types:

List = [term()]
Set = gb_set()

Returns a gb_set of the elements in `List`, where `List` may be unordered and contain duplicates.

from_ordset(List) -> Set

Types:

List = [term()]
Set = gb_set()

Turns an ordered-set list `List` into a gb_set. The list must not contain duplicates.

insert(Element, Set1) -> Set2

Types:

Element = term()
Set1 = Set2 = gb_set()

Returns a new gb_set formed from `Set1` with `Element` inserted. Assumes that `Element` is not present in `Set1`.

intersection(Set1, Set2) -> Set3

Types:

Set1 = Set2 = Set3 = gb_set()

Returns the intersection of `Set1` and `Set2`.

intersection(SetList) -> Set

Types:

SetList = [gb_set()]

Set = gb_set()

Returns the intersection of the non-empty list of gb_sets.

`is_disjoint(Set1, Set2) -> bool()`

Types:

Set1 = Set2 = gb_set()

Returns `true` if `Set1` and `Set2` are disjoint (have no elements in common), and `false` otherwise.

`is_empty(Set) -> bool()`

Types:

Set = gb_set()

Returns `true` if `Set` is an empty set, and `false` otherwise.

`is_member(Element, Set) -> bool()`

`is_element(Element, Set) -> bool()`

Types:

Element = term()

Set = gb_set()

Returns `true` if `Element` is an element of `Set`, otherwise `false`.

`is_set(Term) -> bool()`

Types:

Term = term()

Returns `true` if `Set` appears to be a gb_set, otherwise `false`.

`is_subset(Set1, Set2) -> bool()`

Types:

Set1 = Set2 = gb_set()

Returns `true` when every element of `Set1` is also a member of `Set2`, otherwise `false`.

`iterator(Set) -> Iter`

Types:

Set = gb_set()

Iter = term()

Returns an iterator that can be used for traversing the entries of `Set`; see `next/1`. The implementation of this is very efficient; traversing the whole set using `next/1` is only slightly slower than getting the list of all elements using `to_list/1` and traversing that. The main advantage of the iterator approach is that it does not require the complete list of all elements to be built in memory at one time.

`largest(Set) -> term()`

Types:

Set = gb_set()

Returns the largest element in `Set`. Assumes that `Set` is nonempty.

`next(Iter1) -> {Element, Iter2} | none`

Types:

Iter1 = Iter2 = Element = term()

Returns `{Element, Iter2}` where `Element` is the smallest element referred to by the iterator `Iter1`, and `Iter2` is the new iterator to be used for traversing the remaining elements, or the atom `none` if no elements remain.

`singleton(Element) -> gb_set()`

Types:

Element = term()

Returns a gb_set containing only the element `Element`.

`size(Set) -> int()`

Types:

Set = gb_set()

Returns the number of elements in `Set`.

`smallest(Set) -> term()`

Types:

Set = gb_set()

Returns the smallest element in `Set`. Assumes that `Set` is nonempty.

`take_largest(Set1) -> {Element, Set2}`

Types:

Set1 = Set2 = gb_set()

Element = term()

Returns `{Element, Set2}`, where `Element` is the largest element in `Set1`, and `Set2` is this set with `Element` deleted. Assumes that `Set1` is nonempty.

`take_smallest(Set1) -> {Element, Set2}`

Types:

Set1 = Set2 = gb_set()

Element = term()

Returns `{Element, Set2}`, where `Element` is the smallest element in `Set1`, and `Set2` is this set with `Element` deleted. Assumes that `Set1` is nonempty.

`to_list(Set) -> List`

Types:

Set = gb_set()

List = [term()]

Returns the elements of `Set` as a list.

union(Set1, Set2) -> Set3

Types:

Set1 = Set2 = Set3 = gb_set()

Returns the merged (union) gb_set of `Set1` and `Set2`.

union(SetList) -> Set

Types:

SetList = [gb_set()]

Set = gb_set()

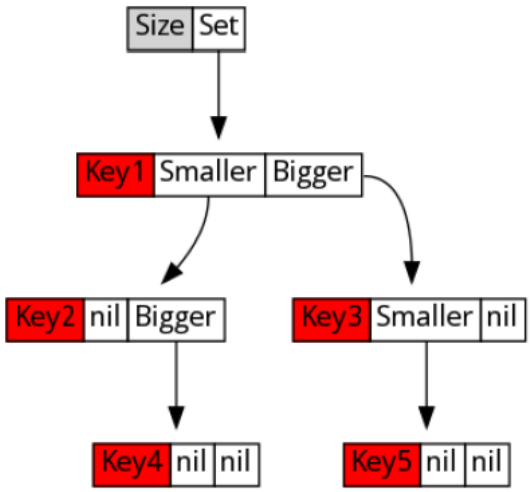
Returns the merged (union) gb_set of the list of gb_sets.

SEE ALSO

gb_trees(3), ordsets(3), sets(3)

gb_sets

- gb_sets 同 gb_trees 结构基本一致，只是结点上只存储键，不再存储值。
- gb_sets 内部结构如下图所示：



- gb_sets 进行集合间操作时，总是会将元素数量较少的集合 X 转换为有序列表，并根据另一个集合 Y 的元素数量使用不同的方法进行遍历操作：
  - 若 **|Y| < 10**，则集合 Y 也被转换为有序列表，X 和 Y 之间的集合操作通过列表归并进行，结果列表再转换为 gb_sets 结构；
  - 若 **|Y|/|X| < c*log[2](|Y|)**，则仍然将 Y 转换为有序列表处理，这里 c 设为 1，且因 math 模块中没有以 2 为底的对数函数，条件被改为 **|Y| < |X|*c1*ln(|Y|)**，其中 **c1=c/ln(2)=1.46**；
  - 以上条件都不满足时，保持 Y 为树形结构不变进行遍历；

General Balanced Trees, on the other hand, have a bunch more functions leaving you more direct control over how the structure is to be used. There are basically two modes for gb_trees: the mode where you know your structure in and out (I call this the 'smart mode', and the mode where you can't assume much about it (I call this one the 'naive mode'). In naive mode, the functions are gb_trees:enter/2, gb_trees:lookup/2 and gb_trees:delete_any/2. The related smart functions are gb_trees:insert/3, gb_trees:get/2, gb_trees:update/3 and gb_trees:delete/2. There is also gb_trees:map/2, which is always a nice thing when you need it.

The disadvantage of 'naive' functions over 'smart' ones is that because gb_trees are balanced trees, whenever you insert a new element (or delete a bunch), it might be possible that the tree will need to balance itself. This can take time and memory (even in useless checks just to make sure). The 'smart' function all assume that the key is present in the tree: this lets you skip all the safety checks and results in faster times.

gb_trees

Erlang module

An efficient implementation of Prof. Arne Andersson's General Balanced Trees. These have no storage overhead compared to unbalanced binary trees, and their performance is in general better than AVL trees.

This module considers two keys as different if and only if they do not compare equal (==).

```
5>  gb_trees:insert(b,c,gb_trees:empty()).
{1,{b,c,nil,nil}}
6>  gb_trees:insert(q,w,v(5)).
{2,{b,c,nil,{q,w,nil,nil}}}
7>  gb_trees:insert(k,r,v(6)).
{3,{b,c,nil,{q,w,{k,r,nil,nil},nil}}}
8>  gb_trees:insert(g,r,v(7)).
{4,{b,c,nil,{q,w,{k,r,{g,r,nil,nil},nil},nil}}}
9>  gb_trees:insert(g,t,v(7)).
{4,{b,c,nil,{q,w,{k,r,{g,t,nil,nil},nil},nil}}}
```

Data structure

Data structure:

通用平衡树,这里没有非平衡二叉树的存储问题.其性能整体上要优于AVL树.

```
- {Size, Tree}, where `Tree' is composed of nodes of the form:
  - {Key, Value, Smaller, Bigger}, and the "empty tree" node:
  - nil.
```

There is no attempt to balance trees after deletions. Since deletions do not increase the height of a tree, this should be OK. 删除之后没有尝试重新新平衡树 因为删除操作并增加树高

Original balance condition *h(T) <= ceil(c * log(|T|))* has been changed to the similar (but not quite equivalent) condition *2 ^ h(T) <= |T| ^ c*. This should also be OK.

Performance is comparable to the AVL trees in the Erlang book (and faster in general due to less overhead); the difference is that deletion works for these trees, but not for the book's trees. Behaviour is logarithmic (as it should be).

DATA TYPES

```
gb_tree() = a GB tree
```

Exports

```
14> mochiweb_headers:make([{"set-cookie", foo}, {"set-cookie", bar}, {"set-cookie", baz}]).
{1, {"set-cookie",  {"set-cookie",{array,["baz","bar","foo"]}},  nil,nil}}
15> H =mochiweb_headers:make([{"set-cookie", foo}, {"set-cookie", bar}, {"set-cookie", baz}]).
{1, {"set-cookie",  {"set-cookie",{array,["baz","bar","foo"]}},  nil,nil}}
16> mochiweb_headers:get_value("set-cookie", H).
"foo, bar, baz"
```

`balance(Tree1) -> Tree2`

Types:

Tree1 = Tree2 = gb_tree()

Rebalances `Tree1`. Note that this is rarely necessary, but may be motivated when a large number of nodes have been deleted from the tree without further insertions. Rebalancing could then be forced in order to minimise lookup times, since deletion only does not rebalance the tree.

`delete(Key, Tree1) -> Tree2`

Types:

Key = term()

Tree1 = Tree2 = gb_tree()

Removes the node with key `Key` from `Tree1`; returns new tree. Assumes that the key is present in the tree, crashes otherwise.

delete_any(Key, Tree1) -> Tree2

Types:

Key = term()

Tree1 = Tree2 = gb_tree()

Removes the node with key `Key` from `Tree1` if the key is present in the tree, otherwise does nothing; returns new tree.

empty() -> Tree

Types:

Tree = gb_tree()

Returns a new empty tree

enter(Key, Val, Tree1) -> Tree2

```
10>
10>  gb_trees:enter(g,kill,v(7)).
{4,{b,c,nil,{q,w,{k,r,{g,kill,nil,nil},nil},nil}}}
11>  gb_trees:enter(tk,too,v(7)).
{4,{b,c,nil,{q,w,{k,r,nil,nil},{tk,too,nil,nil}}}}
12>
```

Types:

Key = Val = term()

Tree1 = Tree2 = gb_tree()

Inserts `Key` with value `Val` into `Tree1` if the key is not present in the tree, otherwise updates `Key` to value `Val` in `Tree1`. Returns the new tree.

from_orddict(List) -> Tree

Types:

List = [{Key, Val}]

Key = Val = term()

Tree = gb_tree()

Turns an ordered list `List` of key-value tuples into a tree. The list must not contain duplicate keys.

get(Key, Tree) -> Val

Types:

Key = Val = term()

Tree = gb_tree()

Retrieves the value stored with `Key` in `Tree`. Assumes that the key is present in the tree, crashes otherwise.

lookup(Key, Tree) -> {value, Val} | none

Types:

Key = Val = term()

Tree = gb_tree()

Looks up `Key` in `Tree`; returns `{value, Val}`, or `none` if `Key` is not present.

insert(Key, Val, Tree1) -> Tree2

Types:

Key = Val = term()

Tree1 = Tree2 = gb_tree()

Inserts `Key` with value `Val` into `Tree1`; returns the new tree. Assumes that the key is not present in the tree, crashes otherwise.

is_defined(Key, Tree) -> bool()

Types:

Tree = gb_tree()

Returns `true` if `Key` is present in `Tree`, otherwise `false`.

is_empty(Tree) -> bool()

Types:

Tree = gb_tree()

Returns `true` if `Tree` is an empty tree, and `false` otherwise.

iterator(Tree) -> Iter

Types:

Tree = gb_tree()

Iter = term()

Returns an iterator that can be used for traversing the entries of `Tree`; see `next/1`. The implementation of this is very efficient; traversing the whole tree using `next/1` is only slightly slower than getting the list of all elements using `to_list/1` and traversing that. The main advantage of the iterator approach is that it does not require the complete list of all elements to be built in memory at one time.

keys(Tree) -> [Key]

Types:

Tree = gb_tree()

Key = term()

Returns the keys in `Tree` as an ordered list.

largest(Tree) -> {Key, Val}

Types:

Tree = gb_tree()

Key = Val = term()

Returns `{Key, Val}`, where `Key` is the largest key in `Tree`, and `Val` is the value associated with this key. Assumes that the tree is nonempty.

map(Function, Tree1) -> Tree2

Types:

Function = fun(K, V1) -> V2

Tree1 = Tree2 = gb_tree()

maps the function F(K, V1) -> V2 to all key-value pairs of the tree Tree1 and returns a new tree Tree2 with the same set of keys as Tree1 and the new set of values V2.

next(Iter1) -> {Key, Val, Iter2} | none

Types:

Iter1 = Iter2 = Key = Val = term()

Returns `{Key, Val, Iter2}` where `Key` is the smallest key referred to by the iterator `Iter1`, and `Iter2` is the new iterator to be used for traversing the remaining nodes, or the atom `none` if no nodes remain.

`size(Tree) -> int()`

Types:

Tree = **gb_tree()**

Returns the number of nodes in `Tree`.

`smallest(Tree) -> {Key, Val}`

Types:

Tree = **gb_tree()**

Key = **Val** = **term()**

Returns `{Key, Val}`, where `Key` is the smallest key in `Tree`, and `Val` is the value associated with this key. Assumes that the tree is nonempty.

`take_largest(Tree1) -> {Key, Val, Tree2}`

Types:

Tree1 = **Tree2** = **gb_tree()**

Key = **Val** = **term()**

Returns `{Key, Val, Tree2}`, where `Key` is the largest key in `Tree1`, `Val` is the value associated with this key, and `Tree2` is this tree with the corresponding node deleted. Assumes that the tree is nonempty.

`take_smallest(Tree1) -> {Key, Val, Tree2}`

Types:

Tree1 = **Tree2** = **gb_tree()**

Key = **Val** = **term()**

Returns `{Key, Val, Tree2}`, where `Key` is the smallest key in `Tree1`, `Val` is the value associated with this key, and `Tree2` is this tree with the corresponding node deleted. Assumes that the tree is nonempty.

`to_list(Tree) -> [{Key, Val}]`

Types:

Tree = **gb_tree()**

Key = **Val** = **term()**

Converts a tree into an ordered list of key-value tuples.

`update(Key, Val, Tree1) -> Tree2`

Types:

Key = **Val** = **term()**

Tree1 = **Tree2** = **gb_tree()**

Updates `Key` to value `Val` in `Tree1`; returns the new tree. Assumes that the key is present in the tree.

`values(Tree) -> [Val]`

Types:

Tree = gb_tree()

Val = term()

Returns the values in `Tree` as an ordered list, sorted by their corresponding keys. Duplicates are not removed.

SEE ALSO

gb_sets(3), dict(3)

gb_trees 是一种平衡二叉查找树结构，基于论文 General Balanced Trees 中描述的算法构建，它可以在不额外在结点上记录数据的情况下，通过简单的平衡策略达到平均状况下接近 AVL 树或红黑树的访问效率；

原文中每个 GB 树需要记录 2 个全局数据：树中结点个数 |T| 和自上次全局重平衡后的结点删除次数 d(T)。

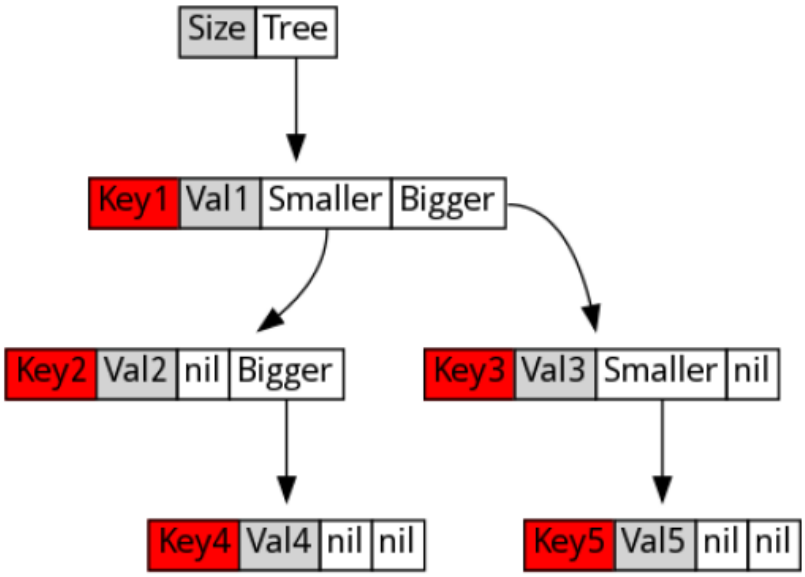
基于这 2 个数据可定义重平衡策略如下：
[1] 发生结点插入时，若新增结点 v 的高度 h(v) 满足条件 h(v) > ceil(c*log[2](|T|+d(T)))，则沿插入路径回溯寻找高度最小的结点 u 使得 h(u) > ceil(c*log[2](|u|))，然后对以 u 为根的子树进行局部重平衡操作，最后增加全局计数 |T|；

[2] 发生结点删除时，增加全局计数 d(T)，若条件 d(T) >= (2^(b/c)-1)*|T| 满足，则对整个树进行全局重平衡，并重置 d(T) 为 0。
这里 c、b 为常数且满足条件 c > 1、b > 0，使用这种重平衡策略后树的最大高度为 max(h(T)) = ceil(c*log[2](|T|)+b)

Erlang 的 gb_trees 并没有完全按照论文的方式实现，而是进行了若干折中处理：
因为结点删除不会增加树的高度，故删除操作不进行自动重平衡，以降低删除开销。在大规模删除结点后，为了让查找效率最优化，开发人员可以显式调用 balance/1 方法对整棵树进行重平衡；

由于不进行自动重平衡，gb_trees 不需要记录删除次数 d(T)，只有一个全局数据 |T| 需要记录；
局部重平衡子树的根结点查找策略从原来的 h(u) > ceil(c*log[2](|u|)) 改为 2^h(u) > |u|^c，以提高计算效率。
虽然变换前后的条件不完全相同，但实测差异不大；

- gb_trees 内部结构如下图所示：



- 其中字段含义为：
  - Size - 树中结点总数 |T|；
  - Tree - 保存整棵树结构；
  - Key - 结点键；
  - Val - 结点值；
  - Smaller - 左子树结构，其中所有结点键都小于当前结点键；
  - Bigger - 右子树结构，其中所有结点键都大于当前结点键；

smallest/1、largest/1、is_defined/2、lookup/2、get/2 都是尾递归操作，附加空间开销很低；
所有对树进行变更的操作都是普通递归操作，数据集较大时附加空间开销很大；
iterator/1 创建的迭代器是一个列表，记录了树中序遍历调用栈，next/2 只是基于迭代器列表进行简单的弹栈、压栈操作而已；

gen_event

Erlang module

A behaviour module for implementing event handling functionality. The OTP event handling model consists of a generic event manager process with an arbitrary number of event handlers which are added and deleted dynamically.

An event manager implemented using this module will have a standard set of interface functions and include functionality for tracing and error reporting. It will also fit into an OTP supervision tree. Refer to *OTP Design Principles* for more information.

Each event handler is implemented as a callback module exporting a pre-defined set of functions. The relationship between the behaviour functions and the callback functions can be illustrated as follows:

```
gen_event module                    Callback module
----------------                    ---------------
gen_event:start_link       ----->   -

gen_event:add_handler
gen_event:add_sup_handler  ----->   Module:init/1

gen_event:notify
gen_event:sync_notify      ----->   Module:handle_event/2

gen_event:call             ----->   Module:handle_call/2

-                          ----->   Module:handle_info/2

gen_event:delete_handler   ----->   Module:terminate/2

gen_event:swap_handler
gen_event:swap_sup_handler ----->   Module1:terminate/2
                                    Module2:init/1

gen_event:which_handlers   ----->   -

gen_event:stop             ----->   Module:terminate/2

-                          ----->   Module:code_change/3
```

Since each event handler is one callback module, an event manager will have several callback modules which are added and deleted dynamically. Therefore `gen_event` is more tolerant of callback module errors than the other behaviours. If a callback function for an installed event handler fails with `Reason`, or returns a bad value `Term`, the event manager will not fail. It will delete the event handler by calling the callback function `Module:terminate/2` (see below), giving as argument `{error,{'EXIT',Reason}}` or `{error,Term}`, respectively. No other event handler will be affected.

A gen_event process handles system messages as documented in *sys(3)*. The `sys` module can be used for debugging an event manager.

Note that an event manager *does* trap exit signals automatically.

The gen_event process can go into hibernation (see *erlang(3)*) if a callback function in a handler module specifies `'hibernate'` in its return value. This might be useful if the server is expected to be idle for a long time. However this feature should be used with care as hibernation implies at least two garbage collections (when hibernating and shortly after waking up) and is not something you'd want to do between each event handled by a busy event manager.

gen_event的容错性更强,已经加载的event_handler出错,Event Manager也不会失败.
采取的措施是:Event Manager会调用该handler的terminate方法完成对这个event handler的移除
其它的event_handler并不受影响.

It's also worth noting that when multiple event handlers are invoked, it's sufficient that one single event handler returns a `'hibernate'` request for the whole event manager to go into hibernation.

Unless otherwise stated, all functions in this module fail if the specified event manager does not exist or if bad arguments are given.

Exports

`start_link() -> Result`

`start_link(EventMgrName) -> Result`

Types:

EventMgrName = {local,Name} | {global,Name}

Name = atom()

Result = {ok,Pid} | {error,{already_started,Pid}}

Pid = pid()

Creates an event manager process as part of a supervision tree. The function should be called, directly or indirectly, by the supervisor. It will, among other things, ensure that the event manager is linked to the supervisor.

If `EventMgrName={local,Name}`, the event manager is registered locally as `Name` using `register/2`. If `EventMgrName={global,Name}`, the event manager is registered globally as `Name` using `global:register_name/2`. If no name is provided, the event manager is not registered.

If the event manager is successfully created the function returns `{ok,Pid}`, where `Pid` is the pid of the event manager. If there already exists a process with the specified `EventMgrName` the function returns `{error,{already_started,Pid}}`, where `Pid` is the pid of that process.

`start() -> Result`

`start(EventMgrName) -> Result`

Types:

EventMgrName = {local,Name} | {global,Name}

Name = atom()

Result = {ok,Pid} | {error,{already_started,Pid}}

Pid = pid()

Creates a stand-alone event manager process, i.e. an event manager which is not part of a supervision tree and thus has no supervisor.

See `start_link/0,1` for a description of arguments and return values.

`add_handler(EventMgrRef, Handler, Args) -> Result`

Types:

EventMgr = Name | {Name,Node} | {global,Name} | pid()

Name = Node = atom()

Handler = Module | {Module,Id}

Module = atom()

Id = term()

Args = term()

Result = ok | {'EXIT',Reason} | term()

Reason = term()

Adds a new event handler to the event manager `EventMgrRef`. The event manager will call `Module:init/1` to initiate the event handler and its internal state.

`EventMgrRef` can be:

- the pid,
- `Name`, if the event manager is locally registered,
- `{Name,Node}`, if the event manager is locally registered at another node, or
- `{global,Name}`, if the event manager is globally registered.

`Handler` is the name of the callback module `Module` or a tuple `{Module,Id}`, where `Id` is any term. The `{Module,Id}` representation makes it possible to identify a specific event handler when there are several event handlers using the same callback module.

`Args` is an arbitrary term which is passed as the argument to `Module:init/1`.

If `Module:init/1` returns a correct value, the event manager adds the event handler and this function returns `ok`. If `Module:init/1` fails with `Reason` or returns an unexpected value `Term`, the event handler is ignored and this function returns `{'EXIT',Reason}` or `Term`, respectively.

`add_sup_handler(EventMgrRef, Handler, Args) -> Result`

Types:

EventMgr = Name | {Name,Node} | {global,Name} | pid()

Name = Node = atom()

Handler = Module | {Module,Id}

Module = atom()

Id = term()

Args = term()

Result = ok | {'EXIT',Reason} | term()

Reason = term()

Adds a new event handler in the same way as `add_handler/3` but will also supervise the connection between the event handler and the calling process.

- If the calling process later terminates with `Reason`, the event manager will delete the event handler by calling `Module:terminate/2` with `{stop,Reason}` as argument.
- If the event handler later is deleted, the event manager sends a message`{gen_event_EXIT,Handler,Reason}` to the calling process. `Reason` is one of the following:
  - `normal`, if the event handler has been removed due to a call to `delete_handler/3`, or `remove_handler` has been returned by a callback function (see below).
  - `shutdown`, if the event handler has been removed because the event manager is terminating.
  - `{swapped,NewHandler,Pid}`, if the process `Pid` has replaced the event handler with another event handler `NewHandler` using a call to `swap_handler/3` or `swap_sup_handler/3`.
  - a term, if the event handler is removed due to an error. Which term depends on the error.

See `add_handler/3` for a description of the arguments and return values.

notify(EventMgrRef, Event) -> ok

sync_notify(EventMgrRef, Event) -> ok

Types:

EventMgrRef = Name | {Name,Node} | {global,Name} | pid()
Name = Node = atom()
Event = term()

Sends an event notification to the event manager `EventMgrRef`. The event manager will call `Module:handle_event/2` for each installed event handler to handle the event.

`notify` is asynchronous and will return immediately after the event notification has been sent. `sync_notify` is synchronous in the sense that it will return `ok` after the event has been handled by all event handlers.

See `add_handler/3` for a description of `EventMgrRef`.

`Event` is an arbitrary term which is passed as one of the arguments to `Module:handle_event/2`.

`notify` will not fail even if the specified event manager does not exist, unless it is specified as `Name`.

call(EventMgrRef, Handler, Request) -> Result

call(EventMgrRef, Handler, Request, Timeout) -> Result

Types:

EventMgrRef = Name | {Name,Node} | {global,Name} | pid()
Name = Node = atom()
Handler = Module | {Module,Id}
Module = atom()
Id = term()
Request = term()
Timeout = int()>0 | infinity
Result = Reply | {error,Error}
Reply = term()
Error = bad_module | {'EXIT',Reason} | term()
Reason = term()

Makes a synchronous call to the event handler `Handler` installed in the event manager `EventMgrRef` by sending a request and waiting until a reply arrives or a timeout occurs. The event manager will call `Module:handle_call/2` to handle the request.

See `add_handler/3` for a description of `EventMgrRef` and `Handler`.

`Request` is an arbitrary term which is passed as one of the arguments to `Module:handle_call/2`.

`Timeout` is an integer greater than zero which specifies how many milliseconds to wait for a reply, or the atom `infinity` to wait indefinitely. Default value is 5000. If no reply is received within the specified time, the function call fails.

The return value `Reply` is defined in the return value of `Module:handle_call/2`. If the specified event handler is not installed, the function returns `{error,bad_module}`. If the callback function fails with `Reason` or returns an unexpected value `Term`, this function returns `{error,{'EXIT',Reason}}` or `{error,Term}`, respectively.

delete_handler(EventMgrRef, Handler, Args) -> Result

Types:

EventMgrRef = Name | {Name,Node} | {global,Name} | pid()
Name = Node = atom()
Handler = Module | {Module,Id}
Module = atom()
Id = term()
Args = term()
Result = term() | {error,module_not_found} | {'EXIT',Reason}
Reason = term()

Deletes an event handler from the event manager `EventMgrRef`. The event manager will call `Module:terminate/2` to terminate the event handler.

See `add_handler/3` for a description of `EventMgrRef` and `Handler`.

`Args` is an arbitrary term which is passed as one of the arguments to `Module:terminate/2`.

The return value is the return value of `Module:terminate/2`. If the specified event handler is not installed, the function returns `{error,module_not_found}`. If the callback function fails with `Reason`, the function returns `{'EXIT',Reason}`.

swap_handler(EventMgrRef, {Handler1,Args1}, {Handler2,Args2}) -> Result

Types:

EventMgrRef = Name | {Name,Node} | {global,Name} | pid()
Name = Node = atom()
Handler1 = Handler2 = Module | {Module,Id}
Module = atom()
Id = term()
Args1 = Args2 = term()
Result = ok | {error,Error}
Error = {'EXIT',Reason} | term()
Reason = term()

Replaces an old event handler with a new event handler in the event manager `EventMgrRef`.

See `add_handler/3` for a description of the arguments.

First the old event handler `Handler1` is deleted. The event manager calls `Module1:terminate(Args1, ...)`, where `Module1` is the callback module of `Handler1`, and collects the return value.

Then the new event handler `Handler2` is added and initiated by calling `Module2:init({Args2,Term})`, where `Module2` is the callback module of `Handler2` and `Term` the return value of `Module1:terminate/2`. This makes it possible to transfer information from `Handler1` to `Handler2`.

The new handler will be added even if the the specified old event handler is not installed in which case `Term=error`, or if `Module1:terminate/2` fails with `Reason` in which case `Term={'EXIT',Reason}`. The old handler will be deleted even if `Module2:init/1` fails.

If there was a supervised connection between `Handler1` and a process `Pid`, there will be a supervised connection between `Handler2` and `Pid` instead.

If `Module2:init/1` returns a correct value, this function returns `ok`. If `Module2:init/1` fails with `Reason` or returns an unexpected value `Term`, this this function returns `{error,{'EXIT',Reason}}` or `{error,Term}`, respectively.

`swap_sup_handler(EventMgrRef, {Handler1,Args1}, {Handler2,Args2}) -> Result`

Types:

EventMgrRef = Name | {Name,Node} | {global,Name} | pid()
Name = Node = atom()
Handler1 = Handler 2 = Module | {Module,Id}
Module = atom()
Id = term()
Args1 = Args2 = term()
Result = ok | {error,Error}
Error = {'EXIT',Reason} | term()
Reason = term()

Replaces an event handler in the event manager `EventMgrRef` in the same way as `swap_handler/3` but will also supervise the connection between `Handler2` and the calling process.

See `swap_handler/3` for a description of the arguments and return values.

`which_handlers(EventMgrRef) -> [Handler]`

Types:

EventMgrRef = Name | {Name,Node} | {global,Name} | pid()
Name = Node = atom()
Handler = Module | {Module,Id}
Module = atom()
Id = term()

Returns a list of all event handlers installed in the event manager `EventMgrRef`.

See `add_handler/3` for a description of `EventMgrRef` and `Handler`.

`stop(EventMgrRef) -> ok`

Types:

EventMgrRef = Name | {Name,Node} | {global,Name} | pid()
Name = Node = atom()

Terminates the event manager `EventMgrRef`. Before terminating, the event manager will call `Module:terminate(stop,...)` for each installed event handler.

See `add_handler/3` for a description of the argument.

CALLBACK FUNCTIONS

The following functions should be exported from a `gen_event` callback module.

Exports

`Module:init(InitArgs) -> {ok,State} | {ok,State,hibernate}`

Types:

InitArgs = Args | {Args,Term}

Args = Term = term()

State = term()

Whenever a new event handler is added to an event manager, this function is called to initialize the event handler.

If the event handler is added due to a call to `gen_event:add_handler/3` or `gen_event:add_sup_handler/3`, `InitArgs` is the `Args` argument of these functions.

If the event handler is replacing another event handler due to a call to `gen_event:swap_handler/3` or `gen_event:swap_sup_handler/3`, or due to a `swap` return tuple from one of the other callback functions, `InitArgs` is a tuple `{Args,Term}` where `Args` is the argument provided in the function call/return tuple and `Term` is the result of terminating the old event handler, see `gen_event:swap_handler/3`.

The function should return `{ok,State}` or `{ok,State, hibernate}` where `State` is the initial internal state of the event handler.

If `{ok,State,hibernate}` is returned, the event manager will go into hibernation (by calling *proc_lib:hibernate/3*), waiting for the next event to occur.

`Module:handle_event(Event, State) -> Result`

Types:

Event = term()

State = term()

Result = {ok,NewState} | {ok,NewState,hibernate}

| {swap_handler,Args1,NewState,Handler2,Args2} | remove_handler

NewState = term()

Args1 = Args2 = term()

Handler2 = Module2 | {Module2,Id}

Module2 = atom()

Id = term()

Whenever an event manager receives an event sent using `gen_event:notify/2` or `gen_event:sync_notify/2`, this function is called for each installed event handler to handle the event.

`Event` is the `Event` argument of `notify`/`sync_notify`.

`State` is the internal state of the event handler.

If the function returns `{ok,NewState}` or `{ok,NewState,hibernate}` the event handler will remain in the event manager with the possible updated internal state `NewState`.

If `{ok,NewState,hibernate}` is returned, the event manager will also go into hibernation (by calling *proc_lib:hibernate/3*), waiting for the next event to occur. It is sufficient that one of the event handlers return `{ok,NewState,hibernate}` for the whole event manager process to hibernate.

If the function returns `{swap_handler,Args1,NewState,Handler2,Args2}` the event handler will be replaced by `Handler2` by first calling `Module:terminate(Args1,NewState)` and then `Module2:init({Args2,Term})` where `Term` is the return value of `Module:terminate/2`. See `gen_event:swap_handler/3` for more information.

If the function returns `remove_handler` the event handler will be deleted by calling `Module:terminate(remove_handler,State)`.

`Module:handle_call(Request, State) -> Result`

Types:

Request = term()
State = term()
Result = {ok,Reply,NewState} | {ok,Reply,NewState,hibernate}
| {swap_handler,Reply,Args1,NewState,Handler2,Args2}
| {remove_handler, Reply}
Reply = term()
NewState = term()
Args1 = Args2 = term()
Handler2 = Module2 | {Module2,Id}
Module2 = atom()
Id = term()

Whenever an event manager receives a request sent using `gen_event:call/3,4`, this function is called for the specified event handler to handle the request.

`Request` is the `Request` argument of `call`.

`State` is the internal state of the event handler.

The return values are the same as for `handle_event/2` except they also contain a term `Reply` which is the reply given back to the client as the return value of `call`.

`Module:handle_info(Info, State) -> Result`

Types:

Info = term()
State = term()
Result = {ok,NewState} | {ok,NewState,hibernate}
| {swap_handler,Args1,NewState,Handler2,Args2} | remove_handler
NewState = term()
Args1 = Args2 = term()
Handler2 = Module2 | {Module2,Id}
Module2 = atom()
Id = term()

This function is called for each installed event handler when an event manager receives any other message than an event or a synchronous request (or a system message).

`Info` is the received message.

See `Module:handle_event/2` for a description of State and possible return values.

`Module:terminate(Arg, State) -> term()`

Types:

Arg = Args | {stop,Reason} | stop | remove_handler
| {error,{'EXIT',Reason}} | {error,Term}

Args = Reason = Term = term()

Whenever an event handler is deleted from an event manager, this function is called. It should be the opposite of `Module:init/1` and do any necessary cleaning up.

If the event handler is deleted due to a call to `gen_event:delete_handler`, `gen_event:swap_handler/3` or `gen_event:swap_sup_handler/3`, `Arg` is the `Args` argument of this function call.

`Arg={stop,Reason}` if the event handler has a supervised connection to a process which has terminated with reason `Reason`.

`Arg=stop` if the event handler is deleted because the event manager is terminating.

The event manager will terminate if it is part of a supervision tree and it is ordered by its supervisor to terminate. Even if it is *not* part of a supervision tree, it will terminate if it receives an `'EXIT'` message from its parent.

`Arg=remove_handler` if the event handler is deleted because another callback function has returned `remove_handler` or `{remove_handler,Reply}`.

`Arg={error,Term}` if the event handler is deleted because a callback function returned an unexpected value `Term`, or `Arg={error,{'EXIT',Reason}}` if a callback function failed.

`State` is the internal state of the event handler.

The function may return any term. If the event handler is deleted due to a call to `gen_event:delete_handler`, the return value of that function will be the return value of this function. If the event handler is to be replaced with another event handler due to a swap, the return value will be passed to the `init` function of the new event handler. Otherwise the return value is ignored.

`Module:code_change(OldVsn, State, Extra) -> {ok, NewState}`

Types:

OldVsn = Vsn | {down, Vsn}

Vsn = term()

State = NewState = term()

Extra = term()

This function is called for an installed event handler which should update its internal state during a release upgrade/downgrade, i.e. when the instruction `{update,Module,Change,...}` where `Change={advanced,Extra}` is given in the `.appup` file. See *OTP Design Principles* for more information.

In the case of an upgrade, `OldVsn` is `Vsn`, and in the case of a downgrade, `OldVsn` is `{down,Vsn}`. `Vsn` is defined by the `vsn` attribute(s) of the old version of the callback module `Module`. If no such attribute is defined, the version is the checksum of the BEAM file.

`State` is the internal state of the event handler.

`Extra` is passed as-is from the `{advanced,Extra}` part of the update instruction.

The function should return the updated internal state.

`Module:format_status(Opt, [PDict, State]) -> Status`

Types:

Opt = normal | terminate

PDict = [{Key, Value}]

State = term()

Status = term()

> **Note:**
>
> This callback is optional, so event handler modules need not export it. If a handler does not export this function, the gen_event module uses the handler state directly for the purposes described below.

This function is called by a gen_event process when:

- One of *sys:get_status/1,2* is invoked to get the gen_event status. `Opt` is set to the atom `normal` for this case.
- The event handler terminates abnormally and gen_event logs an error. `Opt` is set to the atom `terminate` for this case.

This function is useful for customising the form and appearance of the event handler state for these cases. An event handler callback module wishing to customise the `sys:get_status/1,2` return value as well as how its state appears in termination error logs exports an instance of `format_status/2` that returns a term describing the current state of the event handler.

`PDict` is the current value of the gen_event's process dictionary.

`State` is the internal state of the event handler.

The function should return `Status`, a term that customises the details of the current state of the event handler. Any term is allowed for `Status`. The gen_event module uses `Status` as follows:

- When `sys:get_status/1,2` is called, gen_event ensures that its return value contains `Status` in place of the event handler's actual state term.
- When an event handler terminates abnormally, gen_event logs `Status` in place of the event handler's actual state term.

One use for this function is to return compact alternative state representations to avoid having large state terms printed in logfiles.

SEE ALSO

supervisor(3), sys(3)

gen_fsm

Erlang module

A behaviour module for implementing a finite state machine. A generic finite state machine process (gen_fsm) implemented using this module will have a standard set of interface functions and include functionality for tracing and error reporting. It will also fit into an OTP supervision tree. Refer to *OTP Design Principles* for more information.

A gen_fsm assumes all specific parts to be located in a callback module exporting a pre-defined set of functions. The relationship between the behaviour functions and the callback functions can be illustrated as follows:

```
gen_fsm module                  Callback module
--------------                  ---------------
gen_fsm:start_link              -----> Module:init/1

gen_fsm:send_event              -----> Module:StateName/2

gen_fsm:send_all_state_event    -----> Module:handle_event/3

gen_fsm:sync_send_event         -----> Module:StateName/3

gen_fsm:sync_send_all_state_event -----> Module:handle_sync_event/4

-                               -----> Module:handle_info/3

-                               -----> Module:terminate/3

-                               -----> Module:code_change/4
```

这里的方法可以分成两组,普通的事件派发和同步调用等待返回结果的事件触发,由于后者需要把结果进行返回所以多了一个From参数

If a callback function fails or returns a bad value, the gen_fsm will terminate.

A gen_fsm handles system messages as documented in *sys(3)*. The `sys` module can be used for debugging a gen_fsm.

Note that a gen_fsm does not trap exit signals automatically, this must be explicitly initiated in the callback module.

Unless otherwise stated, all functions in this module fail if the specified gen_fsm does not exist or if bad arguments are given.

The gen_fsm process can go into hibernation (see *erlang(3)*) if a callback function specifies `'hibernate'` instead of a timeout value. This might be useful if the server is expected to be idle for a long time. However this feature should be used with care as hibernation implies at least two garbage collections (when hibernating and shortly after waking up) and is not something you'd want to do between each call to a busy state machine.

在被频繁调用的情况下不要使用hibernate选项 除了服务性质的进程 我觉得可以忽略掉这个属性

Exports

如果回调函数调用失败或者返回了不恰当的值,gen_fsm就会terminate
注意:gen_fsm并不是gen_server的实现

`start_link(Module, Args, Options) -> Result`

`start_link(FsmName, Module, Args, Options) -> Result`

Types:

FsmName = {local,Name} | {global,GlobalName}

Name = atom()

GlobalName = term()

Module = atom()

Args = term()

Options = [Option]
Option = {debug,Dbgs} | {timeout,Time} | {spawn_opt,SOpts}
Dbgs = [Dbg]
Dbg = trace | log | statistics
| {log_to_file,FileName} | {install,{Func,FuncState}}
SOpts = [SOpt]
SOpt - see erlang:spawn_opt/2,3,4,5
Result = {ok,Pid} | ignore | {error,Error}
Pid = pid()
Error = {already_started,Pid} | term()

Creates a gen_fsm process as part of a supervision tree. The function should be called, directly or indirectly, by the supervisor. It will, among other things, ensure that the gen_fsm is linked to the supervisor.

The gen_fsm process calls `Module:init/1` to initialize. To ensure a synchronized start-up procedure, `start_link/3,4` does not return until `Module:init/1` has returned.

If `FsmName={local,Name}`, the gen_fsm is registered locally as `Name` using `register/2`. If `FsmName={global,GlobalName}`, the gen_fsm is registered globally as `GlobalName` using `global:register_name/2`. If no name is provided, the gen_fsm is not registered.

`Module` is the name of the callback module.

`Args` is an arbitrary term which is passed as the argument to `Module:init/1`.

If the option `{timeout,Time}` is present, the gen_fsm is allowed to spend `Time` milliseconds initializing or it will be terminated and the start function will return `{error,timeout}`.

If the option `{debug,Dbgs}` is present, the corresponding `sys` function will be called for each item in `Dbgs`. See *sys(3)*.

If the option `{spawn_opt,SOpts}` is present, `SOpts` will be passed as option list to the `spawn_opt` BIF which is used to spawn the gen_fsm process. See *erlang(3)*.

> **Note:**
>
> Using the spawn option `monitor` is currently not allowed, but will cause the function to fail with reason `badarg`.

If the gen_fsm is successfully created and initialized the function returns `{ok,Pid}`, where `Pid` is the pid of the gen_fsm. If there already exists a process with the specified `FsmName`, the function returns `{error,{already_started,Pid}}` where `Pid` is the pid of that process.

If `Module:init/1` fails with `Reason`, the function returns `{error,Reason}`. If `Module:init/1` returns `{stop,Reason}` or `ignore`, the process is terminated and the function returns `{error,Reason}` or `ignore`, respectively.

gen_server提供了一个start方法可以在shell中启动,为什么fsm没有?
这个设计是出于什么考虑?

靠,继续往下看,还真有

`start(Module, Args, Options) -> Result`

`start(FsmName, Module, Args, Options) -> Result`

Types:

FsmName = {local,Name} | {global,GlobalName}

Name = atom()
GlobalName = term()
Module = atom()
Args = term()
Options = [Option]
Option = {debug,Dbgs} | {timeout,Time} | {spawn_opt,SOpts}
Dbgs = [Dbg]
Dbg = trace | log | statistics
| {log_to_file,FileName} | {install,{Func,FuncState}}
SOpts = [term()]
Result = {ok,Pid} | ignore | {error,Error}
Pid = pid()
Error = {already_started,Pid} | term()

Creates a stand-alone gen_fsm process, i.e. a gen_fsm which is not part of a supervision tree and thus has no supervisor.

See *start_link/3,4* for a description of arguments and return values.

同步

`send_event(FsmRef, Event) -> ok`

gen_fsm是事件触发状态改变,所以使用的是send_event,
这里使用的第一个参数是Fsm的ref

Types:

FsmRef = Name | {Name,Node} | {global,GlobalName} | pid()
Name = Node = atom()
GlobalName = term()
Event = term()

当前状态会捕捉这个事件,进行处理
Event实际上是EventArg 是具体的消息数据了

Sends an event asynchronously to the gen_fsm `FsmRef` and returns `ok` immediately. The gen_fsm will call `Module:StateName/2` to handle the event, where `StateName` is the name of the current state of the gen_fsm.

`FsmRef` can be:

- the pid,
- `Name`, if the gen_fsm is locally registered,
- `{Name,Node}`, if the gen_fsm is locally registered at another node, or
- `{global,GlobalName}`, if the gen_fsm is globally registered.

`Event` is an arbitrary term which is passed as one of the arguments to `Module:StateName/2`.

Event是作为可以传送给Module:StateName/2的任意Eralng项

异步

`send_all_state_event(FsmRef, Event) -> ok`

Types:

由于调用的时候外部事件根本不知道当前的状态,所以这里只有一个Event参数

FsmRef = Name | {Name,Node} | {global,GlobalName} | pid()
Name = Node = atom()
GlobalName = term()
Event = term()

Sends an event asynchronously to the gen_fsm `FsmRef` and returns `ok` immediately. The gen_fsm will call `Module:handle_event/3` to handle the event.

See *send_event/2* for a description of the arguments.

这个方法的应用场景是什么? 无论现在状态机在什么状态,接收到这个事件之后都会执行同一个逻辑,比如清理结束的逻辑

send_event vs. send_all_state_event
区别在于用来处理事件的回调函数不同.这个方法的使用场景是:所有状态接受到这个事件都会执行相同的逻辑 ,所以只需要一个handle_event回调函数就可以了.而不是在每一个state都写一个相同的处理方式.

The difference between `send_event` and `send_all_state_event` is which callback function is used to handle the event. This function is useful when sending events that are handled the same way in every state, as only one `handle_event` clause is needed to handle the event instead of one clause in each state name function.

`sync_send_event(FsmRef, Event) -> Reply`

`sync_send_event(FsmRef, Event, Timeout) -> Reply`

Types: 同步发送事件并获取事件发生的结果

FsmRef = Name | {Name,Node} | {global,GlobalName} | pid()

Name = Node = atom()

GlobalName = term()

Event = term()

Timeout = int()>0 | infinity

Reply = term()

Sends an event to the gen_fsm `FsmRef` and waits until a reply arrives or a timeout occurs. The gen_fsm will call `Module:StateName/3` to handle the event, where `StateName` is the name of the current state of the gen_fsm.

See *send_event/2* for a description of `FsmRef` and `Event`.

`Timeout` is an integer greater than zero which specifies how many milliseconds to wait for a reply, or the atom `infinity` to wait indefinitely. Default value is 5000. If no reply is received within the specified time, the function call fails. 默认值是5000,还可以指定为infinity

The return value `Reply` is defined in the return value of `Module:StateName/3`.

The ancient behaviour of sometimes consuming the server exit message if the server died during the call while linked to the client has been removed in OTP R12B/Erlang 5.6.

`sync_send_all_state_event(FsmRef, Event) -> Reply`

`sync_send_all_state_event(FsmRef, Event, Timeout) -> Reply`

Types:

FsmRef = Name | {Name,Node} | {global,GlobalName} | pid()

Name = Node = atom()

GlobalName = term()

Event = term()

Timeout = int()>0 | infinity

Reply = term()

Sends an event to the gen_fsm `FsmRef` and waits until a reply arrives or a timeout occurs. The gen_fsm will call `Module:handle_sync_event/4` to handle the event.

See *send_event/2* for a description of `FsmRef` and `Event`. See *sync_send_event/3* for a description of `Timeout` and `Reply`.

See *send_all_state_event/2* for a discussion about the difference between `sync_send_event` and `sync_send_all_state_event`.

`reply(Caller, Reply) -> true`

Types:

Caller - see below

Reply = term()

This function can be used by a gen_fsm to explicitly send a reply to a client process that called *sync_send_event/2,3* or *sync_send_all_state_event/2,3*, when the reply cannot be defined in the return value of `Module:State/3` or `Module:handle_sync_event/4`.

`Caller` must be the `From` argument provided to the callback function. `Reply` is an arbitrary term, which will be given back to the client as the return value of `sync_send_event/2,3` or `sync_send_all_state_event/2,3`.

`send_event_after(Time, Event) -> Ref`

Types:

Time = integer() 延迟发送消息,返回值是TimerReference

Event = term()

Ref = reference()

Sends a delayed event internally in the gen_fsm that calls this function after `Time` ms. Returns immediately a reference that can be used to cancel the delayed send using *cancel_timer/1*.

The gen_fsm will call `Module:StateName/2` to handle the event, where `StateName` is the name of the current state of the gen_fsm at the time the delayed event is delivered.

`Event` is an arbitrary term which is passed as one of the arguments to `Module:StateName/2`.

`start_timer(Time, Msg) -> Ref`

Types:

Time = integer()

Msg = term()

Ref = reference()

Sends a timeout event internally in the gen_fsm that calls this function after `Time` ms. Returns immediately a reference that can be used to cancel the timer using *cancel_timer/1*.

The gen_fsm will call `Module:StateName/2` to handle the event, where `StateName` is the name of the current state of the gen_fsm at the time the timeout message is delivered.

`Msg` is an arbitrary term which is passed in the timeout message, `{timeout, Ref, Msg}`, as one of the arguments to `Module:StateName/2`.

`cancel_timer(Ref) -> RemainingTime | false`

Types:

Ref = reference()

RemainingTime = integer()

Cancels an internal timer referred by `Ref` in the gen_fsm that calls this function.

`Ref` is a reference returned from *send_event_after/2* or *start_timer/2*.

If the timer has already timed out, but the event not yet been delivered, it is cancelled as if it had *not* timed out, so there will be no false timer event after returning from this function.

Returns the remaining time in ms until the timer would have expired if `Ref` referred to an active timer, `false` otherwise.

```
enter_loop(Module, Options, StateName, StateData)

enter_loop(Module, Options, StateName, StateData, FsmName)

enter_loop(Module, Options, StateName, StateData, Timeout)

enter_loop(Module, Options, StateName, StateData, FsmName, Timeout)
```

Types:

Module = atom()
Options = [Option]
Option = {debug,Dbgs}
Dbgs = [Dbg]
Dbg = trace | log | statistics
| {log_to_file,FileName} | {install,{Func,FuncState}}
StateName = atom()
StateData = term()
FsmName = {local,Name} | {global,GlobalName}
Name = atom()
GlobalName = term()
Timeout = int() | infinity

Makes an existing process into a gen_fsm. Does not return, instead the calling process will enter the gen_fsm receive loop and become a gen_fsm process. The process *must* have been started using one of the start functions in `proc_lib`, see *proc_lib(3)*. The user is responsible for any initialization of the process, including registering a name for it.

This function is useful when a more complex initialization procedure is needed than the gen_fsm behaviour provides.

`Module`, `Options` and `FsmName` have the same meanings as when calling *start[_link]/3,4*. However, if `FsmName` is specified, the process must have been registered accordingly *before* this function is called.

`StateName`, `StateData` and `Timeout` have the same meanings as in the return value of *Module:init/1*. Also, the callback module `Module` does not need to export an `init/1` function.

Failure: If the calling process was not started by a `proc_lib` start function, or if it is not registered according to `FsmName`.

CALLBACK FUNCTIONS

The following functions should be exported from a `gen_fsm` callback module.

In the description, the expression *state name* is used to denote a state of the state machine. *state data* is used to denote the internal state of the Erlang process which implements the state machine.

Exports

有限状态机的loopdata分为两部分:状态+数据

```
Module:init(Args) -> Result
```

Types:

Args = term()
Return = {ok,StateName,StateData} | {ok,StateName,StateData,Timeout}
| {ok,StateName,StateData,hibernate}

这里的超时值是这样的作用:在Timeout毫秒时间内,如果没有事件或者消息到来,timeout事件就会被触发.
就是说初始化状态之后,首先会等待一个超时
(demo@192.168.1.199)1> rgy:start().
{ok,<0.38.0>}
(demo@192.168.1.199)2> red timeout.

| {stop,Reason} | ignore
StateName = atom() 状态机的初始态
StateData = term()
Timeout = int()>0 | infinity
Reason = term()

Whenever a gen_fsm is started using *gen_fsm:start/3,4* or *gen_fsm:start_link/3,4*, this function is called by the new process to initialize.

`Args` is the `Args` argument provided to the start function.

If initialization is successful, the function should return `{ok,StateName,StateData}`, `{ok,StateName,StateData,Timeout}` or `{ok,StateName,StateData,hibernate}`, where `StateName` is the initial state name and `StateData` the initial state data of the gen_fsm.

If an integer timeout value is provided, a timeout will occur unless an event or a message is received within `Timeout` milliseconds. A timeout is represented by the atom `timeout` and should be handled by the `Module:StateName/2` callback functions. The atom `infinity` can be used to wait indefinitely, this is the default value.

If `hibernate` is specified instead of a timeout value, the process will go into hibernation when waiting for the next message to arrive (by calling *proc_lib:hibernate/3*).

If something goes wrong during the initialization the function should return `{stop,Reason}`, where `Reason` is any term, or `ignore`.

`Module:StateName(Event, StateData) -> Result`

Types: 这个方法实际上是自解释的,包含了当前状态 当前状态数据 和发送过来的事件 所以发送send_event方法也只需要发送ent参数信息就够了

Event = timeout | term()
StateData = term()
Result = {next_state,NextStateName,NewStateData}
| {next_state,NextStateName,NewStateData,Timeout}
| {next_state,NextStateName,NewStateData,hibernate}
| {stop,Reason,NewStateData}
NextStateName = atom()
NewStateData = term()
Timeout = int()>0 | infinity
Reason = term()

There should be one instance of this function for each possible state name. Whenever a gen_fsm receives an event sent using *gen_fsm:send_event/2*, the instance of this function with the same name as the current state name `StateName` is called to handle the event. It is also called if a timeout occurs.

`Event` is either the atom `timeout`, if a timeout has occurred, or the `Event` argument provided to `send_event/2`.

`StateData` is the state data of the gen_fsm.

If the function returns `{next_state,NextStateName,NewStateData}`, `{next_state,NextStateName,NewStateData,Timeout}` or `{next_state,NextStateName,NewStateData,hibernate}`, the gen_fsm will continue executing with the current state name set to `NextStateName` and with the possibly updated state data `NewStateData`. See `Module:init/1` for a description of `Timeout` and `hibernate`.

If the function returns `{stop,Reason,NewStateData}`, the gen_fsm will call `Module:terminate(Reason,NewStateData)` and terminate.

`Module:handle_event(Event, StateName, StateData) -> Result`

Types:

Event = term()
StateName = atom()
StateData = term()
Result = {next_state,NextStateName,NewStateData}
| {next_state,NextStateName,NewStateData,Timeout}
| {next_state,NextStateName,NewStateData,hibernate}
| {stop,Reason,NewStateData}
NextStateName = atom()
NewStateData = term()
Timeout = int()>0 | infinity
Reason = term()

Whenever a gen_fsm receives an event sent using *gen_fsm:send_all_state_event/2*, this function is called to handle the event.

`StateName` is the current state name of the gen_fsm.

See `Module:StateName/2` for a description of the other arguments and possible return values.

`Module:StateName(Event, From, StateData) -> Result`

Types:

Event = term()
From = {pid(),Tag}
StateData = term()
Result = {reply,Reply,NextStateName,NewStateData}
| {reply,Reply,NextStateName,NewStateData,Timeout}
| {reply,Reply,NextStateName,NewStateData,hibernate}
| {next_state,NextStateName,NewStateData}
| {next_state,NextStateName,NewStateData,Timeout}
| {next_state,NextStateName,NewStateData,hibernate}
| {stop,Reason,Reply,NewStateData} | {stop,Reason,NewStateData}
Reply = term()
NextStateName = atom()
NewStateData = term()
Timeout = int()>0 | infinity
Reason = normal | term()

There should be one instance of this function for each possible state name. Whenever a gen_fsm receives an event sent using *gen_fsm:sync_send_event/2,3*, the instance of this function with the same name as the current state name `StateName` is called to handle the event.

`Event` is the `Event` argument provided to `sync_send_event`.

`From` is a tuple `{Pid,Tag}` where `Pid` is the pid of the process which called `sync_send_event/2,3` and `Tag` is a unique tag.

`StateData` is the state data of the gen_fsm.

If the function returns `{reply,Reply,NextStateName,NewStateData}`, `{reply,Reply,NextStateName,NewStateData,Timeout}` or `{reply,Reply,NextStateName,NewStateData,hibernate}`, `Reply` will be given back to `From` as the return value of `sync_send_event/2,3`. The gen_fsm then continues executing with the current state name set to `NextStateName` and with the possibly updated state data `NewStateData`. See `Module:init/1` for a description of `Timeout` and `hibernate`.

If the function returns `{next_state,NextStateName,NewStateData}`, `{next_state,NextStateName,NewStateData,Timeout}` or `{next_state,NextStateName,NewStateData,hibernate}`, the gen_fsm will continue executing in `NextStateName` with `NewStateData`. Any reply to `From` must be given explicitly using *gen_fsm:reply/2*.

If the function returns `{stop,Reason,Reply,NewStateData}`, `Reply` will be given back to `From`. If the function returns `{stop,Reason,NewStateData}`, any reply to `From` must be given explicitly using `gen_fsm:reply/2`. The gen_fsm will then call `Module:terminate(Reason,NewStateData)` and terminate.

`Module:handle_sync_event(Event, From, StateName, StateData) -> Result`

Types:

Event = term()
From = {pid(),Tag}
StateName = atom()
StateData = term()
Result = {reply,Reply,NextStateName,NewStateData}
| {reply,Reply,NextStateName,NewStateData,Timeout}
| {reply,Reply,NextStateName,NewStateData,hibernate}
| {next_state,NextStateName,NewStateData}
| {next_state,NextStateName,NewStateData,Timeout}
| {next_state,NextStateName,NewStateData,hibernate}
| {stop,Reason,Reply,NewStateData} | {stop,Reason,NewStateData}
Reply = term()
NextStateName = atom()
NewStateData = term()
Timeout = int()>0 | infinity
Reason = term()

Whenever a gen_fsm receives an event sent using *gen_fsm:sync_send_all_state_event/2,3*, this function is called to handle the event.

`StateName` is the current state name of the gen_fsm.

See `Module:StateName/3` for a description of the other arguments and possible return values.

`Module:handle_info(Info, StateName, StateData) -> Result`

Types:

Info = term()
StateName = atom()
StateData = term()
Result = {next_state,NextStateName,NewStateData}

> | **{next_state,NextStateName,NewStateData,Timeout}**

> | **{next_state,NextStateName,NewStateData,hibernate}**

> | **{stop,Reason,NewStateData}**

NextStateName = atom()

NewStateData = term()

Timeout = int()>0 | infinity

Reason = normal | term()

This function is called by a gen_fsm when it receives any other message than a synchronous or asynchronous event (or a system message).

`Info` is the received message.

See `Module:StateName/2` for a description of the other arguments and possible return values.

`Module:terminate(Reason, StateName, StateData)`

Types:

Reason = normal | shutdown | {shutdown,term()} | term()

StateName = atom()

StateData = term()

This function is called by a gen_fsm when it is about to terminate. It should be the opposite of `Module:init/1` and do any necessary cleaning up. When it returns, the gen_fsm terminates with `Reason`. The return value is ignored.

`Reason` is a term denoting the stop reason, `StateName` is the current state name, and `StateData` is the state data of the gen_fsm.

`Reason` depends on why the gen_fsm is terminating. If it is because another callback function has returned a stop tuple `{stop,..}`, `Reason` will have the value specified in that tuple. If it is due to a failure, `Reason` is the error reason.

If the gen_fsm is part of a supervision tree and is ordered by its supervisor to terminate, this function will be called with `Reason=shutdown` if the following conditions apply:

- the gen_fsm has been set to trap exit signals, and
- the shutdown strategy as defined in the supervisor's child specification is an integer timeout value, not `brutal_kill`.

Even if the gen_fsm is *not* part of a supervision tree, this function will be called if it receives an `'EXIT'` message from its parent. `Reason` will be the same as in the `'EXIT'` message.

Otherwise, the gen_fsm will be immediately terminated.

Note that for any other reason than `normal`, `shutdown`, or `{shutdown,Term}` the gen_fsm is assumed to terminate due to an error and an error report is issued using *error_logger:format/2*.

`Module:code_change(OldVsn, StateName, StateData, Extra) -> {ok, NextStateName, NewStateData}`

Types:

OldVsn = Vsn | {down, Vsn}

Vsn = term()

StateName = NextStateName = atom()

StateData = NewStateData = term()

Extra = term()

This function is called by a gen_fsm when it should update its internal state data during a release upgrade/downgrade, i.e. when the instruction `{update,Module,Change,...}` where `Change={advanced,Extra}` is given in the `appup` file. See *OTP Design Principles*.

In the case of an upgrade, `OldVsn` is `Vsn`, and in the case of a downgrade, `OldVsn` is `{down,Vsn}`. `Vsn` is defined by the `vsn` attribute(s) of the old version of the callback module `Module`. If no such attribute is defined, the version is the checksum of the BEAM file.

`StateName` is the current state name and `StateData` the internal state data of the gen_fsm.

`Extra` is passed as-is from the `{advanced,Extra}` part of the update instruction.

The function should return the new current state name and updated internal data.

`Module:format_status(Opt, [PDict, StateData]) -> Status`

Types:

Opt = normal | terminate

PDict = [{Key, Value}]

StateData = term()

Status = term()

> Note:
>
> This callback is optional, so callback modules need not export it. The gen_fsm module provides a default implementation of this function that returns the callback module state data.

This function is called by a gen_fsm process when:

- One of *sys:get_status/1,2* is invoked to get the gen_fsm status. `Opt` is set to the atom `normal` for this case.
- The gen_fsm terminates abnormally and logs an error. `Opt` is set to the atom `terminate` for this case.

This function is useful for customising the form and appearance of the gen_fsm status for these cases. A callback module wishing to customise the `sys:get_status/1,2` return value as well as how its status appears in termination error logs exports an instance of `format_status/2` that returns a term describing the current status of the gen_fsm.

`PDict` is the current value of the gen_fsm's process dictionary.

`StateData` is the internal state data of the gen_fsm.

The function should return `Status`, a term that customises the details of the current state and status of the gen_fsm. There are no restrictions on the form `Status` can take, but for the `sys:get_status/1,2` case (when `Opt` is `normal`), the recommended form for the `Status` value is `[{data, [{"StateData", Term}]}]` where `Term` provides relevant details of the gen_fsm state data. Following this recommendation isn't required, but doing so will make the callback module status consistent with the rest of the `sys:get_status/1,2` return value.

One use for this function is to return compact alternative state data representations to avoid having large state terms printed in logfiles.

SEE ALSO

gen_event(3), *gen_server(3)*, *supervisor(3)*, *proc_lib(3)*, *sys(3)*

gen_server

Erlang module

A behaviour module for implementing the server of a client-server relation. A generic server process (gen_server) implemented using this module will have a standard set of interface functions and include functionality for tracing and error reporting. It will also fit into an OTP supervision tree. Refer to *OTP Design Principles* for more information.

A gen_server assumes all specific parts to be located in a callback module exporting a pre-defined set of functions. The relationship between the behaviour functions and the callback functions can be illustrated as follows:

```
gen_server module            Callback module
-----------------            ---------------
gen_server:start_link ---> Module:init/1

gen_server:call
gen_server:multi_call -----> Module:handle_call/3

gen_server:cast
gen_server:abcast     -----> Module:handle_cast/2

-                     -----> Module:handle_info/2

-                     -----> Module:terminate/2

-                     -----> Module:code_change/3
```

If a callback function fails or returns a bad value, the gen_server will terminate.

A gen_server handles system messages as documented in *sys(3)*. The `sys` module can be used for debugging a gen_server.

Note that a gen_server does not trap exit signals automatically, this must be explicitly initiated in the callback module.

Unless otherwise stated, all functions in this module fail if the specified gen_server does not exist or if bad arguments are given.

The gen_server process can go into hibernation (see *erlang(3)*) if a callback function specifies `'hibernate'` instead of a timeout value. This might be useful if the server is expected to be idle for a long time. However this feature should be used with care as hibernation implies at least two garbage collections (when hibernating and shortly after waking up) and is not something you'd want to do between each call to a busy server.

Exports

`start_link(Module, Args, Options) -> Result`

`start_link(ServerName, Module, Args, Options) -> Result`

Types:

ServerName = {local,Name} | {global,GlobalName}

Name = atom()

GlobalName = term()

Module = atom()

Args = term()

Options = [Option]
Option = {debug,Dbgs} | {timeout,Time} | {spawn_opt,SOpts}
Dbgs = [Dbg]
Dbg = trace | log | statistics | {log_to_file,FileName} | {install,{Func,FuncState}}
SOpts = [term()]
Result = {ok,Pid} | ignore | {error,Error}
Pid = pid()
Error = {already_started,Pid} | term()

Creates a gen_server process as part of a supervision tree. The function should be called, directly or indirectly, by the supervisor. It will, among other things, ensure that the gen_server is linked to the supervisor.

The gen_server process calls `Module:init/1` to initialize. To ensure a synchronized start-up procedure, `start_link/3,4` does not return until `Module:init/1` has returned.

If `ServerName={local,Name}` the gen_server is registered locally as `Name` using `register/2`. If `ServerName={global,GlobalName}` the gen_server is registered globally as `GlobalName` using `global:register_name/2`. If no name is provided, the gen_server is not registered.

`Module` is the name of the callback module.

`Args` is an arbitrary term which is passed as the argument to `Module:init/1`.

If the option `{timeout,Time}` is present, the gen_server is allowed to spend `Time` milliseconds initializing or it will be terminated and the start function will return `{error,timeout}`.

If the option `{debug,Dbgs}` is present, the corresponding `sys` function will be called for each item in `Dbgs`. See *sys(3)*.

If the option `{spawn_opt,SOpts}` is present, `SOpts` will be passed as option list to the `spawn_opt` BIF which is used to spawn the gen_server. See *erlang(3)*.

> Note:
>
> Using the spawn option `monitor` is currently not allowed, but will cause the function to fail with reason `badarg`.

If the gen_server is successfully created and initialized the function returns `{ok,Pid}`, where `Pid` is the pid of the gen_server. If there already exists a process with the specified `ServerName` the function returns `{error,{already_started,Pid}}`, where `Pid` is the pid of that process.

If `Module:init/1` fails with `Reason`, the function returns `{error,Reason}`. If `Module:init/1` returns `{stop,Reason}` or `ignore`, the process is terminated and the function returns `{error,Reason}` or `ignore`, respectively.

`start(Module, Args, Options) -> Result`

`start(ServerName, Module, Args, Options) -> Result`

Types:

ServerName = {local,Name} | {global,GlobalName}
Name = atom()
GlobalName = term()

Module = atom()
Args = term()
Options = [Option]
Option = {debug,Dbgs} | {timeout,Time} | {spawn_opt,SOpts}
Dbgs = [Dbg]
Dbg = trace | log | statistics | {log_to_file,FileName} | {install,{Func,FuncState}}
SOpts = [term()]
Result = {ok,Pid} | ignore | {error,Error}
Pid = pid()
Error = {already_started,Pid} | term()

Creates a stand-alone gen_server process, i.e. a gen_server which is not part of a supervision tree and thus has no supervisor.

See *start_link/3,4* for a description of arguments and return values.

`call(ServerRef, Request) -> Reply`

`call(ServerRef, Request, Timeout) -> Reply`

Types:

ServerRef = Name | {Name,Node} | {global,GlobalName} | pid()
Node = atom()
GlobalName = term()
Request = term()
Timeout = int()>0 | infinity
Reply = term()

对于单个gen_sever的操作可以设置超时的时间,注意如果是全局节点,那么也要标识出来时global

Makes a synchronous call to the gen_server `ServerRef` by sending a request and waiting until a reply arrives or a timeout occurs. The gen_server will call `Module:handle_call/3` to handle the request.

`ServerRef` can be:

- the pid,
- `Name`, if the gen_server is locally registered,
- `{Name,Node}`, if the gen_server is locally registered at another node, or
- `{global,GlobalName}`, if the gen_server is globally registered.

`Request` is an arbitrary term which is passed as one of the arguments to `Module:handle_call/3`.

`Timeout` is an integer greater than zero which specifies how many milliseconds to wait for a reply, or the atom `infinity` to wait indefinitely. Default value is 5000. If no reply is received within the specified time, the function call fails. If the caller catches the failure and continues running, and the server is just late with the reply, it may arrive at any time later into the caller's message queue. The caller must in this case be prepared for this and discard any such garbage messages that are two element tuples with a reference as the first element.

The return value `Reply` is defined in the return value of `Module:handle_call/3`.

The call may fail for several reasons, including timeout and the called gen_server dying before or during the call.

The ancient behaviour of sometimes consuming the server exit message if the server died during the call while linked to the client has been removed in OTP R12B/Erlang 5.6.

multi_call(Name, Request) -> Result

multi_call(Nodes, Name, Request) -> Result

multi_call(Nodes, Name, Request, Timeout) -> Result

Types:

Nodes = [Node]
Node = atom()
Name = atom()
Request = term()
Timeout = int()>=0 | infinity
Result = {Replies,BadNodes}
Replies = [{Node,Reply}]
Reply = term()
BadNodes = [Node]

```
%% Call apply(M,F,A) on all nodes in parallel

multicall(Nodes, M, F, A, infinity)
  when is_list(Nodes), is_atom(M), is_atom(F), is_list(A) ->
    do_multicall(Nodes, M, F, A, infinity);
multicall(Nodes, M, F, A, Timeout)
  when is_list(Nodes), is_atom(M), is_atom(F), is_list(A), is_integer(Timeout),
       Timeout >= 0 ->
    do_multicall(Nodes, M, F, A, Timeout).

do_multicall(Nodes, M, F, A, Timeout) ->
    {Rep,Bad} = gen_server:multi_call(Nodes, ?NAME,
          {call, M,F,A, group_leader()},
          Timeout),
    {lists:map(fun({_,R}) -> R end, Rep), Bad}.
```

Makes a synchronous call to all gen_servers locally registered as `Name` at the specified nodes by first sending a request to every node and then waiting for the replies. The gen_servers will call `Module:handle_call/3` to handle the request.

The function returns a tuple `{Replies,BadNodes}` where `Replies` is a list of `{Node,Reply}` and `BadNodes` is a list of node that either did not exist, or where the gen_server `Name` did not exist or did not reply.

`Nodes` is a list of node names to which the request should be sent. Default value is the list of all known nodes `[node()|nodes()]`.

`Name` is the locally registered name of each gen_server.

`Request` is an arbitrary term which is passed as one of the arguments to `Module:handle_call/3`.

`Timeout` is an integer greater than zero which specifies how many milliseconds to wait for each reply, or the atom `infinity` to wait indefinitely. Default value is `infinity`. If no reply is received from a node within the specified time, the node is added to `BadNodes`.

When a reply `Reply` is received from the gen_server at a node `Node`, `{Node,Reply}` is added to `Replies`. `Reply` is defined in the return value of `Module:handle_call/3`.

> ### Warning:
>
> If one of the nodes is not capable of process monitors, for example C or Java nodes, and the gen_server is not started when the requests are sent, but starts within 2 seconds, this function waits the whole `Timeout`, which may be infinity.
>
> This problem does not exist if all nodes are Erlang nodes.

To avoid that late answers (after the timeout) pollutes the caller's message queue, a middleman process is used to do the actual calls. Late answers will then be discarded when they arrive to a terminated process.

cast(ServerRef, Request) -> ok

Types:

ServerRef = Name | {Name,Node} | {global,GlobalName} | pid()

Node = atom()
GlobalName = term()
Request = term()

Sends an asynchronous request to the gen_server `ServerRef` and returns `ok` immediately, ignoring if the destination node or gen_server does not exist. The gen_server will call `Module:handle_cast/2` to handle the request.

See *call/2,3* for a description of `ServerRef`.

`Request` is an arbitrary term which is passed as one of the arguments to `Module:handle_cast/2`.

`abcast(Name, Request) -> abcast`

`abcast(Nodes, Name, Request) -> abcast`

Types:

Nodes = [Node]
Node = atom()
Name = atom()
Request = term()

Sends an asynchronous request to the gen_servers locally registered as `Name` at the specified nodes. The function returns immediately and ignores nodes that do not exist, or where the gen_server `Name` does not exist. The gen_servers will call `Module:handle_cast/2` to handle the request.

See *multi_call/2,3,4* for a description of the arguments.

`reply(Client, Reply) -> Result`

Types:

Client - see below
Reply = term()
Result = term()

This function can be used by a gen_server to explicitly send a reply to a client that called `call/2,3` or `multi_call/2,3,4`, when the reply cannot be defined in the return value of `Module:handle_call/3`.

`Client` must be the `From` argument provided to the callback function. `Reply` is an arbitrary term, which will be given back to the client as the return value of `call/2,3` or `multi_call/2,3,4`.

The return value `Result` is not further defined, and should always be ignored.

`enter_loop(Module, Options, State)`

`enter_loop(Module, Options, State, ServerName)`

`enter_loop(Module, Options, State, Timeout)`

`enter_loop(Module, Options, State, ServerName, Timeout)`

Types:

Module = atom()
Options = [Option]
Option = {debug,Dbgs}

Dbgs = [Dbg]
Dbg = trace | log | statistics
| {log_to_file,FileName} | {install,{Func,FuncState}}
State = term()
ServerName = {local,Name} | {global,GlobalName}
Name = atom()
GlobalName = term()
Timeout = int() | infinity

Makes an existing process into a gen_server. Does not return, instead the calling process will enter the gen_server receive loop and become a gen_server process. The process *must* have been started using one of the start functions in `proc_lib`, see *proc_lib(3)*. The user is responsible for any initialization of the process, including registering a name for it.

This function is useful when a more complex initialization procedure is needed than the gen_server behaviour provides.

`Module`, `Options` and `ServerName` have the same meanings as when calling *gen_server:start[_link]/3,4*. However, if `ServerName` is specified, the process must have been registered accordingly *before* this function is called.

`State` and `Timeout` have the same meanings as in the return value of *Module:init/1*. Also, the callback module `Module` does not need to export an `init/1` function.

Failure: If the calling process was not started by a `proc_lib` start function, or if it is not registered according to `ServerName`.

CALLBACK FUNCTIONS

The following functions should be exported from a `gen_server` callback module.

Exports

`Module:init(Args) -> Result`

Types:

Args = term()
Result = {ok,State} | {ok,State,Timeout} | {ok,State,hibernate}
| {stop,Reason} | ignore
State = term()
Timeout = int()>=0 | infinity
Reason = term()

Whenever a gen_server is started using *gen_server:start/3,4* or *gen_server:start_link/3,4*, this function is called by the new process to initialize.

`Args` is the `Args` argument provided to the start function.

If the initialization is successful, the function should return `{ok,State}`, `{ok,State,Timeout}` or `{ok,State,hibernate}`, where `State` is the internal state of the gen_server.

If an integer timeout value is provided, a timeout will occur unless a request or a message is received within `Timeout` milliseconds. A timeout is represented by the atom `timeout` which should be handled by the `handle_info/2` callback function. The atom `infinity` can be used to wait indefinitely, this is the default value.

If `hibernate` is specified instead of a timeout value, the process will go into hibernation when waiting for the next message to arrive (by calling *proc_lib:hibernate/3*).

这里提供了一个超时时间,除非是接收到请求或者消息,就会发生timeout超时

If something goes wrong during the initialization the function should return `{stop,Reason}` where `Reason` is any term, or `ignore`.

`Module:handle_call(Request, From, State) -> Result`

Types:

Request = term()
From = {pid(),Tag}
State = term()
Result = {reply,Reply,NewState} | {reply,Reply,NewState,Timeout}
| {reply,Reply,NewState,hibernate}
| {noreply,NewState} | {noreply,NewState,Timeout}
| {noreply,NewState,hibernate}
| {stop,Reason,Reply,NewState} | {stop,Reason,NewState}
Reply = term()
NewState = term()
Timeout = int()>=0 | infinity
Reason = term()

Whenever a gen_server receives a request sent using *gen_server:call/2,3* or *gen_server:multi_call/2,3,4*, this function is called to handle the request.

`Request` is the `Request` argument provided to `call` or `multi_call`.

`From` is a tuple `{Pid,Tag}` where `Pid` is the pid of the client and `Tag` is a unique tag.

`State` is the internal state of the gen_server.

If the function returns `{reply,Reply,NewState}`, `{reply,Reply,NewState,Timeout}` or `{reply,Reply,NewState,hibernate}`, `Reply` will be given back to `From` as the return value of `call/2,3` or included in the return value of `multi_call/2,3,4`. The gen_server then continues executing with the possibly updated internal state `NewState`. See `Module:init/1` for a description of `Timeout` and `hibernate`.

If the functions returns `{noreply,NewState}`, `{noreply,NewState,Timeout}` or `{noreply,NewState,hibernate}`, the gen_server will continue executing with `NewState`. Any reply to `From` must be given explicitly using *gen_server:reply/2*.

If the function returns `{stop,Reason,Reply,NewState}`, `Reply` will be given back to `From`. If the function returns `{stop,Reason,NewState}`, any reply to `From` must be given explicitly using `gen_server:reply/2`. The gen_server will then call `Module:terminate(Reason,NewState)` and terminate.

`Module:handle_cast(Request, State) -> Result`

Types:

Request = term()
State = term()
Result = {noreply,NewState} | {noreply,NewState,Timeout}
| {noreply,NewState,hibernate}
| {stop,Reason,NewState}
NewState = term()
Timeout = int()>=0 | infinity
Reason = term()

Whenever a gen_server receives a request sent using *gen_server:cast/2* or *gen_server:abcast/2,3*, this function is called to handle the request.

See `Module:handle_call/3` for a description of the arguments and possible return values.

`Module:handle_info(Info, State) -> Result`

Types:

Info = timeout | term()

State = term()

Result = {noreply,NewState} | {noreply,NewState,Timeout}

| {noreply,NewState,hibernate}

| {stop,Reason,NewState}

NewState = term()

Timeout = int()>=0 | infinity

Reason = normal | term()

This function is called by a gen_server when a timeout occurs or when it receives any other message than a synchronous or asynchronous request (or a system message).

`Info` is either the atom `timeout`, if a timeout has occurred, or the received message.

See `Module:handle_call/3` for a description of the other arguments and possible return values.

`Module:terminate(Reason, State)`

Types:

Reason = normal | shutdown | {shutdown,term()} | term()

State = term()

This function is called by a gen_server when it is about to terminate. It should be the opposite of `Module:init/1` and do any necessary cleaning up. When it returns, the gen_server terminates with `Reason`. The return value is ignored.

`Reason` is a term denoting the stop reason and `State` is the internal state of the gen_server.

`Reason` depends on why the gen_server is terminating. If it is because another callback function has returned a stop tuple `{stop,..}`, `Reason` will have the value specified in that tuple. If it is due to a failure, `Reason` is the error reason.

If the gen_server is part of a supervision tree and is ordered by its supervisor to terminate, this function will be called with `Reason=shutdown` if the following conditions apply:

- the gen_server has been set to trap exit signals, and
- the shutdown strategy as defined in the supervisor's child specification is an integer timeout value, not `brutal_kill`.

Even if the gen_server is *not* part of a supervision tree, this function will be called if it receives an `'EXIT'` message from its parent. `Reason` will be the same as in the `'EXIT'` message.

Otherwise, the gen_server will be immediately terminated.

Note that for any other reason than `normal`, `shutdown`, or `{shutdown,Term}` the gen_server is assumed to terminate due to an error and an error report is issued using *error_logger:format/2*.

`Module:code_change(OldVsn, State, Extra) -> {ok, NewState}`

Types:

OldVsn = Vsn | {down, Vsn}

Vsn = term()

State = NewState = term()

Extra = term()

This function is called by a gen_server when it should update its internal state during a release upgrade/downgrade, i.e. when the instruction `{update,Module,Change,...}` where `Change={advanced,Extra}` is given in the `appup` file. See *OTP Design Principles* for more information.

In the case of an upgrade, `OldVsn` is `Vsn`, and in the case of a downgrade, `OldVsn` is `{down,Vsn}`. `Vsn` is defined by the `vsn` attribute(s) of the old version of the callback module `Module`. If no such attribute is defined, the version is the checksum of the BEAM file.

`State` is the internal state of the gen_server.

`Extra` is passed as-is from the `{advanced,Extra}` part of the update instruction.

The function should return the updated internal state.

`Module:format_status(Opt, [PDict, State]) -> Status`

Types:

Opt = normal | terminate

PDict = [{Key, Value}]

State = term()

Status = term()

> Note:
>
> This callback is optional, so callback modules need not export it. The gen_server module provides a default implementation of this function that returns the callback module state.

This function is called by a gen_server process when:

- One of *sys:get_status/1,2* is invoked to get the gen_server status. `Opt` is set to the atom `normal` for this case.
- The gen_server terminates abnormally and logs an error. `Opt` is set to the atom `terminate` for this case.

This function is useful for customising the form and appearance of the gen_server status for these cases. A callback module wishing to customise the `sys:get_status/1,2` return value as well as how its status appears in termination error logs exports an instance of `format_status/2` that returns a term describing the current status of the gen_server.

`PDict` is the current value of the gen_server's process dictionary.

`State` is the internal state of the gen_server.

The function should return `Status`, a term that customises the details of the current state and status of the gen_server. There are no restrictions on the form `Status` can take, but for the `sys:get_status/1,2` case (when `Opt` is `normal`), the recommended form for the `Status` value is `[{data, [{"State", Term}]}]` where `Term` provides relevant details of the gen_server state. Following this recommendation isn't required, but doing so will make the callback module status consistent with the rest of the `sys:get_status/1,2` return value.

One use for this function is to return compact alternative state representations to avoid having large state terms printed in logfiles.

SEE ALSO

gen_event(3), gen_fsm(3), supervisor(3), proc_lib(3), sys(3)

io

Erlang module

This module provides an interface to standard Erlang IO servers. The output functions all return `ok` if they are successful, or exit if they are not.

In the following description, all functions have an optional parameter `IoDevice`. If included, it must be the pid of a process which handles the IO protocols. Normally, it is the `IoDevice` returned by *file:open/2*.

For a description of the IO protocols refer to the STDLIB Users Guide.

> ### Warning:
>
> As of R13A, data supplied to the *put_chars* function should be in the `chardata()` format described below. This means that programs supplying binaries to this function need to convert them to UTF-8 before trying to output the data on an `io_device()`.
>
> If an io_device() is set in binary mode, the functions *get_chars* and *get_line* may return binaries instead of lists. The binaries will, as of R13A, be encoded in UTF-8.
>
> To work with binaries in ISO-latin-1 encoding, use the *file* module instead.
>
> For conversion functions between character encodings, see the *unicode* module.

DATA TYPES

```
io_device()
  as returned by file:open/2, a process handling IO protocols
```

```
unicode_binary() = binary() with characters encoded in UTF-8 coding standard
unicode_char() = integer() representing valid unicode codepoint

chardata() = charlist() | unicode_binary()

charlist() = [unicode_char() | unicode_binary() | charlist()]
  a unicode_binary is allowed as the tail of the list
```

Exports

`columns([IoDevice]) -> {ok,int()} | {error, enotsup}`

Types:

IoDevice = io_device()

Retrieves the number of columns of the `IoDevice` (i.e. the width of a terminal). The function only succeeds for terminal devices, for all other devices the function returns `{error, enotsup}`

`put_chars([IoDevice,] IoData) -> ok`

Types:

IoDevice = io_device()

IoData = chardata()

Writes the characters of `IoData` to the io_server() (`IoDevice`).

`nl([IoDevice]) -> ok`

Types:

IoDevice = io_device()

Writes new line to the standard output (`IoDevice`).

`get_chars([IoDevice,] Prompt, Count) -> Data | eof`

Types:

IoDevice = io_device()

Prompt = atom() | string()

Count = int()

Data = [unicode_char()] | unicode_binary()

Reads `Count` characters from standard input (`IoDevice`), prompting it with `Prompt`. It returns:

`Data`

The input characters. If the device supports Unicode, the data may represent codepoints larger than 255 (the latin1 range). If the io_server() is set to deliver binaries, they will be encoded in UTF-8 (regardless of if the device actually supports Unicode or not).

`eof`

End of file was encountered.

`{error,Reason}`

Other (rare) error condition, for instance `{error,estale}` if reading from an NFS file system.

`get_line([IoDevice,] Prompt) -> Data | eof | {error,Reason}`

Types:

IoDevice = io_device()

Prompt = atom() | string()

Data = [unicode_char()] | unicode_binary()

Reads a line from the standard input (`IoDevice`), prompting it with `Prompt`. It returns:

`Data`

The characters in the line terminated by a LF (or end of file). If the device supports Unicode, the data may represent codepoints larger than 255 (the latin1 range). If the io_server() is set to deliver binaries, they will be encoded in UTF-8 (regardless of if the device actually supports Unicode or not).

`eof`

End of file was encountered.

`{error,Reason}`

Other (rare) error condition, for instance `{error,estale}` if reading from an NFS file system.

`getopts([IoDevice]) -> Opts`

Types:

IoDevice = io_device()

Opts = [Opt]

Opt = {atom(),Value}

Value = term()

This function requests all available options and their current values for a specific io_device(). Example:

```
1> {ok,F} = file:open("/dev/null",[read]).
{ok,<0.42.0>}
2> io:getopts(F).
[{binary,false},{encoding,latin1}]
```

Here the file I/O-server returns all available options for a file, which are the expected ones, `encoding` and `binary`. The standard shell however has some more options:

```
3> io:getopts().
[{expand_fun,#Fun<group.0.120017273>},
 {echo,true},
 {binary,false},
 {encoding,unicode}]
```

This example is, as can be seen, run in an environment where the terminal supports Unicode input and output.

`setopts([IoDevice,] Opts) -> ok | {error, Reason}`

Types:

IoDevice = io_device()

Opts = [Opt]

Opt = atom() | {atom(),Value}

Value = term()

Reason = term()

Set options for the io_device() (`IoDevice`).

Possible options and values vary depending on the actual io_device(). For a list of supported options and their current values on a specific device, use the *getopts/1* function.

The options and values supported by the current OTP io_devices are:

`binary, list or {binary, bool()}`

If set in binary mode (binary or {binary,true}), the io_server() sends binary data (encoded in UTF-8) as answers to the get_line, get_chars and, if possible, get_until requests (see the I/O protocol description in STDLIB User's Guide for details). The immediate effect is that `get_chars/2,3` and `get_line/1,2` return UTF-8 binaries instead of lists of chars for the affected device.

By default, all io_devices in OTP are set in list mode, but the io functions can handle any of these modes and so should other, user written, modules behaving as clients to I/O-servers.

This option is supported by the standard shell (group.erl), the 'oldshell' (user.erl) and the file I/O servers.

`{echo, bool()}`

Denotes if the terminal should echo input. Only supported for the standard shell I/O-server (group.erl)

`{expand_fun, fun()}`

Provide a function for tab-completion (expansion) like the erlang shell. This function is called when the user presses the Tab key. The expansion is active when calling line-reading functions such as `get_line/1,2`.

The function is called with the current line, upto the cursor, as a reversed string. It should return a three-tuple: `{yes|no, string(), [string(), ...]}`. The first element gives a beep if `no`, otherwise the expansion is silent, the second is a string that will be entered at the cursor position, and the third is a list of possible expansions. If this list is non-empty, the list will be printed and the current input line will be written once again.

Trivial example (beep on anything except empty line, which is expanded to "quit"):

```
fun("") -> {yes, "quit", []};
   (_) -> {no, "", ["quit"]} end
```

This option is supported by the standard shell only (group.erl).

`{encoding, latin1 | unicode}`

Specifies how characters are input or output from or to the actual device, implying that i.e. a terminal is set to handle Unicode input and output or a file is set to handle UTF-8 data encoding.

The option *does not* affect how data is returned from the io-functions or how it is sent in the I/O-protocol, it only affects how the io_device() is to handle Unicode characters towards the "physical" device.

The standard shell will be set for either unicode or latin1 encoding when the system is started. The actual encoding is set with the help of the "LANG" or "LC_CTYPE" environment variables on Unix-like system or by other means on other systems. The bottom line is that the user can input Unicode characters and the device will be in {encoding, unicode} mode if the device supports it. The mode can be changed, if the assumption of the runtime system is wrong, by setting this option.

The io_device() used when Erlang is started with the "-oldshell" or "-noshell" flags is by default set to latin1 encoding, meaning that any characters beyond codepoint 255 will be escaped and that input is expected to be plain 8-bit ISO-latin-1. If the encoding is changed to Unicode, input and output from the standard file descriptors will be in UTF-8 (regardless of operating system).

Files can also be set in {encoding, unicode}, meaning that data is written and read as UTF-8. More encodings are possible for files, see below.

{encoding, unicode | latin1} is supported by both the standard shell (group.erl including werl on windows), the 'oldshell' (user.erl) and the file I/O servers.

`{encoding, utf8 | utf16 | utf32 | {utf16,big} | {utf16,little} | {utf32,big} | {utf32,little}}`

For disk files, the encoding can be set to various UTF variants. This will have the effect that data is expected to be read as the specified encoding from the file and the data will be written in the specified encoding to the disk file.

{encoding, utf8} will have the same effect as {encoding,unicode} on files.

The extended encodings are only supported on disk files (opened by the *file:open/2* function)

`write([IoDevice,] Term) -> ok`

Types:

IoDevice = io_device()

Term = term()

Writes the term `Term` to the standard output (`IoDevice`).

`read([IoDevice,] Prompt) -> Result`

Types:

IoDevice = io_device()
Prompt = atom() | string()
Result = {ok, Term} | eof | {error, ErrorInfo}
Term = term()
ErrorInfo -- see section Error Information below

Reads a term `Term` from the standard input (`IoDevice`), prompting it with `Prompt`. It returns:

`{ok, Term}`

The parsing was successful.

`eof`

End of file was encountered.

`{error, ErrorInfo}`

The parsing failed.

`read(IoDevice, Prompt, StartLine) -> Result`

Types:

IoDevice = io_device()
Prompt = atom() | string()
StartLine = int()
Result = {ok, Term, EndLine} | {eof, EndLine} | {error, ErrorInfo, EndLine}
Term = term()
EndLine = int()
ErrorInfo -- see section Error Information below

Reads a term `Term` from `IoDevice`, prompting it with `Prompt`. Reading starts at line number `StartLine`. It returns:

`{ok, Term, EndLine}`

The parsing was successful.

`{eof, EndLine}`

End of file was encountered.

`{error, ErrorInfo, EndLine}`

The parsing failed.

fwrite(Format) ->

fwrite([IoDevice,] Format, Data) -> ok

format(Format) ->

format([IoDevice,] Format, Data) -> ok

Types:

IoDevice = io_device()

Format = atom() | string() | binary()

Data = [term()]

Writes the items in `Data` (`[]`) on the standard output (`IoDevice`) in accordance with `Format`. `Format` contains plain characters which are copied to the output device, and control sequences for formatting, see below. If `Format` is an atom or a binary, it is first converted to a list with the aid of `atom_to_list/1` or `binary_to_list/1`.

```
1> io:fwrite("Hello world!~n", []).
Hello world!
ok
```

The general format of a control sequence is `~F.P.PadModC`. The character `C` determines the type of control sequence to be used, `F` and `P` are optional numeric arguments. If `F`, `P`, or `Pad` is `*`, the next argument in `Data` is used as the numeric value of `F` or `P`.

`F` is the `field width` of the printed argument. A negative value means that the argument will be left justified within the field, otherwise it will be right justified. If no field width is specified, the required print width will be used. If the field width specified is too small, then the whole field will be filled with `*` characters.

`P` is the `precision` of the printed argument. A default value is used if no precision is specified. The interpretation of precision depends on the control sequences. Unless otherwise specified, the argument `within` is used to determine print width.

`Pad` is the padding character. This is the character used to pad the printed representation of the argument so that it conforms to the specified field width and precision. Only one padding character can be specified and, whenever applicable, it is used for both the field width and precision. The default padding character is `' '` (space).

`Mod` is the control sequence modifier. It is either a single character (currently only 't', for unicode translation, is supported) that changes the interpretation of Data.

The following control sequences are available:

`~`

The character ~ is written.

`c`

The argument is a number that will be interpreted as an ASCII code. The precision is the number of times the character is printed and it defaults to the field width, which in turn defaults to 1. The following example illustrates:

```
2> io:fwrite("|~10.5c|~-10.5c|~5c|~n", [$a, $b, $c]).
|     aaaaa|bbbbb     |ccccc|
ok
```

If the Unicode translation modifier ('t') is in effect, the integer argument can be any number representing a valid unicode codepoint, otherwise it should be an integer less than or equal to 255, otherwise it is masked with 16#FF:

```
1> io:fwrite("~tc~n",[1024]).
\x{400}
ok
2> io:fwrite("~c~n",[1024]).
^@
ok
```

`f`

The argument is a float which is written as `[-]ddd.ddd`, where the precision is the number of digits after the decimal point. The default precision is 6 and it cannot be less than 1.

`e`

The argument is a float which is written as `[-]d.ddde+-ddd`, where the precision is the number of digits written. The default precision is 6 and it cannot be less than 2.

`g`

The argument is a float which is written as `f`, if it is >= 0.1 and < 10000.0. Otherwise, it is written in the `e` format. The precision is the number of significant digits. It defaults to 6 and should not be less than 2. If the absolute value of the float does not allow it to be written in the `f` format with the desired number of significant digits, it is also written in the `e` format.

`s`

Prints the argument with the `string` syntax. The argument is, if no Unicode translation modifier is present, an *I/O list*, a binary, or an atom. If the Unicode translation modifier ('t') is in effect, the argument is chardata(), meaning that binaries are in UTF-8. The characters are printed without quotes. In this format, the printed argument is truncated to the given precision and field width.

This format can be used for printing any object and truncating the output so it fits a specified field:

```
3> io:fwrite("|~10w|~n", [{hey, hey, hey}]).
|**********|
ok
4> io:fwrite("|~10s|~n", [io_lib:write({hey, hey, hey})]).
|{hey,hey,h|
ok
```

A list with integers larger than 255 is considered an error if the Unicode translation modifier is not given:

```
1> io:fwrite("~ts~n",[[1024]]).
\x{400}
ok
2> io:fwrite("~s~n",[[1024]]).
** exception exit: {badarg,[{io,format,[<0.26.0>,"~s~n",[[1024]]]},
   ...
```

w

Writes data with the standard syntax. This is used to output Erlang terms. Atoms are printed within quotes if they contain embedded non-printable characters, and floats are printed accurately as the shortest, correctly rounded string.

p

Writes the data with standard syntax in the same way as ~w, but breaks terms whose printed representation is longer than one line into many lines and indents each line sensibly. It also tries to detect lists of printable characters and to output these as strings. For example:

```
5> T = [{attributes,[[{id,age,1.50000},{mode,explicit},
{typename,"INTEGER"}], [{id,cho},{mode,explicit},{typename,'Cho'}]]},
{typename,'Person'},{tag,{'PRIVATE',3}},{mode,implicit}].
...
6> io:fwrite("~w~n", [T]).
[{attributes,[[{id,age,1.5},{mode,explicit},{typename,
[73,78,84,69,71,69,82]}],[{id,cho},{mode,explicit},{typena
me,'Cho'}]]},{typename,'Person'},{tag,{'PRIVATE',3}},{mode
,implicit}]
ok
7> io:fwrite("~62p~n", [T]).
[{attributes,[[{id,age,1.5},
               {mode,explicit},
               {typename,"INTEGER"}],
              [{id,cho},{mode,explicit},{typename,'Cho'}]]},
 {typename,'Person'},
 {tag,{'PRIVATE',3}},
 {mode,implicit}]
ok
```

The field width specifies the maximum line length. It defaults to 80. The precision specifies the initial indentation of the term. It defaults to the number of characters printed on this line in the `same` call to `io:fwrite` or `io:format`. For example, using `T` above:

```
8> io:fwrite("Here T = ~62p~n", [T]).
Here T = [{attributes,[[{id,age,1.5},
                       {mode,explicit},
                       {typename,"INTEGER"}],
                      [{id,cho},
                       {mode,explicit},
                       {typename,'Cho'}]]},
          {typename,'Person'},
          {tag,{'PRIVATE',3}},
          {mode,implicit}]
ok
```

W

Writes data in the same way as ~w, but takes an extra argument which is the maximum depth to which terms are printed. Anything below this depth is replaced with `...`. For example, using `T` above:

```
9> io:fwrite("~W~n", [T,9]).
[{attributes,[[{id,age,1.5},{mode,explicit},{typename,...}],
[{id,cho},{mode,...},{...}]]},{typename,'Person'},
{tag,{'PRIVATE',3}},{mode,implicit}]
```

```
ok
```

If the maximum depth has been reached, then it is impossible to read in the resultant output. Also, the `,...` form in a tuple denotes that there are more elements in the tuple but these are below the print depth.

P

Writes data in the same way as ~p, but takes an extra argument which is the maximum depth to which terms are printed. Anything below this depth is replaced with `...`. For example:

```
10> io:fwrite("~62P~n", [T,9]).
[{attributes,[[{id,age,1.5},{mode,explicit},{typename,...}],
              [{id,cho},{mode,...},{...}]]},
 {typename,'Person'},
 {tag,{'PRIVATE',3}},
 {mode,implicit}]
ok
```

B

Writes an integer in base 2..36, the default base is 10. A leading dash is printed for negative integers.

The precision field selects base. For example:

```
11> io:fwrite("~.16B~n", [31]).
1F
ok
12> io:fwrite("~.2B~n", [-19]).
-10011
ok
13> io:fwrite("~.36B~n", [5*36+35]).
5Z
ok
```

X

Like B, but takes an extra argument that is a prefix to insert before the number, but after the leading dash, if any.

The prefix can be a possibly deep list of characters or an atom.

```
14> io:fwrite("~X~n", [31,"10#"]).
10#31
ok
15> io:fwrite("~.16X~n", [-31,"0x"]).
-0x1F
ok
```

#

Like B, but prints the number with an Erlang style '#'-separated base prefix.

```
16> io:fwrite("~.10#~n", [31]).
10#31
ok
17> io:fwrite("~.16#~n", [-31]).
-16#1F
```

```
ok
```

`b`

Like `B`, but prints lowercase letters.

`x`

Like `X`, but prints lowercase letters.

`+`

Like #, but prints lowercase letters.

`n`

Writes a new line.

`i`

Ignores the next term.

Returns:

`ok`

The formatting succeeded.

If an error occurs, there is no output. For example:

```
18> io:fwrite("~s ~w ~i ~w ~c ~n",['abc def', 'abc def', {foo, 1},{foo, 1}, 65]).
abc def 'abc def'  {foo,1} A
ok
19> io:fwrite("~s", [65]).
** exception exit: {badarg,[{io,format,[<0.22.0>,"~s","A"]},
                            {erl_eval,do_apply,5},
                            {shell,exprs,6},
                            {shell,eval_exprs,6},
                            {shell,eval_loop,3}]}
     in function  io:o_request/2
```

In this example, an attempt was made to output the single character '65' with the aid of the string formatting directive "~s".

`fread([IoDevice,] Prompt, Format) -> Result`

Types:

IoDevice = io_device()

Prompt = atom() | string()

Format = string()

Result = {ok, Terms} | eof | {error, What}

Terms = [term()]

What = term()

Reads characters from the standard input (`IoDevice`), prompting it with `Prompt`. Interprets the characters in accordance with `Format`. `Format` contains control sequences which directs the interpretation of the input.

`Format` may contain:

- White space characters (SPACE, TAB and NEWLINE) which cause input to be read to the next non-white space character.

- Ordinary characters which must match the next input character.
- Control sequences, which have the general format `~*FMC`. The character `*` is an optional return suppression character. It provides a method to specify a field which is to be omitted. `F` is the `field width` of the input field, `M` is an optional translation modifier (of which 't' is the only currently supported, meaning Unicode translation) and `C` determines the type of control sequence.

  Unless otherwise specified, leading white-space is ignored for all control sequences. An input field cannot be more than one line wide. The following control sequences are available:

  `~`

  A single ~ is expected in the input.

  `d`

  A decimal integer is expected.

  `u`

  An unsigned integer in base 2..36 is expected. The field width parameter is used to specify base. Leading white-space characters are not skipped.

  `-`

  An optional sign character is expected. A sign character '-' gives the return value `-1`. Sign character '+' or none gives `1`. The field width parameter is ignored. Leading white-space characters are not skipped.

  `#`

  An integer in base 2..36 with Erlang-style base prefix (for example `"16#ffff"`) is expected.

  `f`

  A floating point number is expected. It must follow the Erlang floating point number syntax.

  `s`

  A string of non-white-space characters is read. If a field width has been specified, this number of characters are read and all trailing white-space characters are stripped. An Erlang string (list of characters) is returned.

  If Unicode translation is in effect (~ts), characters larger than 255 are accepted, otherwise not. With the translation modifier, the list returned may as a consequence also contain integers larger than 255:

```
1> io:fread("Prompt> ","~s").
Prompt> <Characters beyond latin1 range not printable in this medium>
{error,{fread,string}}
2> io:fread("Prompt> ","~ts").
Prompt> <Characters beyond latin1 range not printable in this medium>
{ok,[[1091,1085,1080,1094,1086,1076,1077]]}
```

  `a`

  Similar to `s`, but the resulting string is converted into an atom.

  The Unicode translation modifier is not allowed (atoms can not contain characters beyond the latin1 range).

  `c`

  The number of characters equal to the field width are read (default is 1) and returned as an Erlang string. However, leading and trailing white-space characters are not omitted as they are with `s`. All characters are returned.

  The Unicode translation modifier works as with `s`:

```
1> io:fread("Prompt> ","~c").
Prompt> <Character beyond latin1 range not printable in this medium>
{error,{fread,string}}
2> io:fread("Prompt> ","~tc").
Prompt> <Character beyond latin1 range not printable in this medium>
{ok,[[1091]]}
```

`l`

Returns the number of characters which have been scanned up to that point, including white-space characters.

It returns:

`{ok, Terms}`

The read was successful and `Terms` is the list of successfully matched and read items.

`eof`

End of file was encountered.

`{error, What}`

The read operation failed and the parameter `What` gives a hint about the error.

Examples:

```
20> io:fread('enter>', "~f~f~f").
enter>1.9 35.5e3 15.0
{ok,[1.9,3.55e4,15.0]}
21> io:fread('enter>', "~10f~d").
enter>     5.67899
{ok,[5.678,99]}
22> io:fread('enter>', ":~10s:~10c:").
enter>:   alan   :   joe    :
{ok, ["alan", "   joe    "]}
```

`rows([IoDevice]) -> {ok,int()} | {error, enotsup}`

Types:

IoDevice = io_device()

Retrieves the number of rows of the `IoDevice` (i.e. the height of a terminal). The function only succeeds for terminal devices, for all other devices the function returns `{error, enotsup}`

`scan_erl_exprs(Prompt) ->`

`scan_erl_exprs([IoDevice,] Prompt, StartLine) -> Result`

Types:

IoDevice = io_device()

Prompt = atom() | string()

StartLine = int()

Result = {ok, Tokens, EndLine} | {eof, EndLine} | {error, ErrorInfo, EndLine}

Tokens -- see erl_scan(3)

EndLine = int()

ErrorInfo -- see section Error Information below

Reads data from the standard input (`IoDevice`), prompting it with `Prompt`. Reading starts at line number `StartLine` (1). The data is tokenized as if it were a sequence of Erlang expressions until a final `'.'` is reached. This token is also returned. It returns:

`{ok, Tokens, EndLine}`

The tokenization succeeded.

`{eof, EndLine}`

End of file was encountered.

`{error, ErrorInfo, EndLine}`

An error occurred.

Example:

```
23> io:scan_erl_exprs('enter>').
enter>abc(), "hey".
{ok,[{atom,1,abc},{'(',1},{')',1},{',',1},{string,1,"hey"},{dot,1}],2}
24> io:scan_erl_exprs('enter>').
enter>1.0er.
{error,{1,erl_scan,{illegal,float}},2}
```

`scan_erl_form(Prompt) ->`

`scan_erl_form([IoDevice,] Prompt, StartLine) -> Result`

Types:

IoDevice = io_device()

Prompt = atom() | string()

StartLine = int()

Result = {ok, Tokens, EndLine} | {eof, EndLine} | {error, ErrorInfo, EndLine}

Tokens -- see erl_scan(3)

EndLine = int()

ErrorInfo -- see section Error Information below

Reads data from the standard input (`IoDevice`), prompting it with `Prompt`. Starts reading at line number `StartLine` (1). The data is tokenized as if it were an Erlang form - one of the valid Erlang expressions in an Erlang source file - until a final `'.'` is reached. This last token is also returned. The return values are the same as for `scan_erl_exprs/1,2,3` above.

`parse_erl_exprs(Prompt) ->`

`parse_erl_exprs([IoDevice,] Prompt, StartLine) -> Result`

Types:

IoDevice = io_device()

Prompt = atom() | string()

StartLine = int()

Result = {ok, Expr_list, EndLine} | {eof, EndLine} | {error, ErrorInfo, EndLine}

Expr_list -- see erl_parse(3)

EndLine = int()

ErrorInfo -- see section Error Information below

Reads data from the standard input (`IoDevice`), prompting it with `Prompt`. Starts reading at line number `StartLine` (1). The data is tokenized and parsed as if it were a sequence of Erlang expressions until a final '.' is reached. It returns:

`{ok, Expr_list, EndLine}`

The parsing was successful.

`{eof, EndLine}`

End of file was encountered.

`{error, ErrorInfo, EndLine}`

An error occurred.

Example:

```
25> io:parse_erl_exprs('enter>').
enter>abc(), "hey".
{ok, [{call,1,{atom,1,abc},[]},{string,1,"hey"}],2}
26> io:parse_erl_exprs ('enter>').
enter>abc("hey".
{error,{1,erl_parse,["syntax error before: ",["'.'"]]},2}
```

`parse_erl_form(Prompt) ->`

`parse_erl_form([IoDevice,] Prompt, StartLine) -> Result`

Types:

IoDevice = io_device()

Prompt = atom() | string()

StartLine = int()

Result = {ok, AbsForm, EndLine} | {eof, EndLine} | {error, ErrorInfo, EndLine}

AbsForm -- see erl_parse(3)

EndLine = int()

ErrorInfo -- see section Error Information below

Reads data from the standard input (`IoDevice`), prompting it with `Prompt`. Starts reading at line number `StartLine` (1). The data is tokenized and parsed as if it were an Erlang form - one of the valid Erlang expressions in an Erlang source file - until a final '.' is reached. It returns:

`{ok, AbsForm, EndLine}`

The parsing was successful.

`{eof, EndLine}`

End of file was encountered.

`{error, ErrorInfo, EndLine}`

An error occurred.

Standard Input/Output

All Erlang processes have a default standard IO device. This device is used when no `IoDevice` argument is specified in the above function calls. However, it is sometimes desirable to use an explicit `IoDevice` argument which refers to the default IO device. This is the case with functions that can access either a file or the default IO device. The atom `standard_io` has this special meaning. The following example illustrates this:

```
27> io:read('enter>').
enter>foo.
{ok,foo}
28> io:read(standard_io, 'enter>').
enter>bar.
{ok,bar}
```

There is always a process registered under the name of `user`. This can be used for sending output to the user.

Standard Error

In certain situations, especially when the standard output is redirected, access to an io_server() specific for error messages might be convenient. The io_device 'standard_error' can be used to direct output to whatever the current operating system considers a suitable device for error output. Example on a Unix-like operating system:

```
$ erl -noshell -noinput -eval 'io:format(standard_error,"Error: ~s~n",["error 11"]),'\
'init:stop().' > /dev/null
Error: error 11
```

Error Information

The `ErrorInfo` mentioned above is the standard `ErrorInfo` structure which is returned from all IO modules. It has the format:

```
{ErrorLine, Module, ErrorDescriptor}
```

A string which describes the error is obtained with the following call:

```
Module:format_error(ErrorDescriptor)
```

io_lib

Erlang module

This module contains functions for converting to and from strings (lists of characters). They are used for implementing the functions in the `io` module. There is no guarantee that the character lists returned from some of the functions are flat, they can be deep lists. `lists:flatten/1` can be used for flattening deep lists.

DATA TYPES

```
chars() = [char() | chars()]
```

Exports

`nl() -> chars()`

Returns a character list which represents a new line character.

`write(Term) ->`

`write(Term, Depth) -> chars()`

```
1>  lists:flatten(io_lib:write([{1,2},{2,3},{3,4}])).
"[{1,2},{2,3},{3,4}]"
2> S= lists:flatten(io_lib:write([{1,2},{2,3},{3,4}])).
"[{1,2},{2,3},{3,4}]"
3> E=S++".".
"[{1,2},{2,3},{3,4}]."
```

Types:

Term = term()

Depth = int()

Returns a character list which represents `Term`. The `Depth` (-1) argument controls the depth of the structures written. When the specified depth is reached, everything below this level is replaced by "...". For example:

```
1> lists:flatten(io_lib:write({1,[2],[3],[4,5],6,7,8,9})).
"{1,[2],[3],[4,5],6,7,8,9}"
2> lists:flatten(io_lib:write({1,[2],[3],[4,5],6,7,8,9}, 5)).
"{1,[2],[3],[...],...}"
```

`print(Term) ->`

`print(Term, Column, LineLength, Depth) -> chars()`

Types:

Term = term()

Column = LineLenght = Depth = int()

Also returns a list of characters which represents `Term`, but breaks representations which are longer than one line into many lines and indents each line sensibly. It also tries to detect and output lists of printable characters as strings. `Column` is the starting column (1), `LineLength` the maximum line length (80), and `Depth` (-1) the maximum print depth.

fwrite(Format, Data) ->

format(Format, Data) -> chars() | UnicodeList

Types:

Format = atom() | string() | binary()

Data = [term()]

UnicodeList = [Unicode]

Unicode = int() representing valid unicode codepoint

Returns a character list which represents `Data` formatted in accordance with `Format`. See *io:fwrite/1,2,3* for a detailed description of the available formatting options. A fault is generated if there is an error in the format string or argument list.

If (and only if) the Unicode translation modifier is used in the format string (i.e. ~ts or ~tc), the resulting list may contain characters beyond the ISO-latin-1 character range (in other words, numbers larger than 255). If so, the result is not an ordinary Erlang string(), but can well be used in any context where Unicode data is allowed.

fread(Format, String) -> Result

Types:

Format = String = string()

Result = {ok, InputList, LeftOverChars} | {more, RestFormat, Nchars, InputStack} | {error, What}

InputList = chars()

LeftOverChars = string()

RestFormat = string()

Nchars = int()

InputStack = chars()

What = term()

Tries to read `String` in accordance with the control sequences in `Format`. See *io:fread/3* for a detailed description of the available formatting options. It is assumed that `String` contains whole lines. It returns:

`{ok, InputList, LeftOverChars}`

The string was read. `InputList` is the list of successfully matched and read items, and `LeftOverChars` are the input characters not used.

`{more, RestFormat, Nchars, InputStack}`

The string was read, but more input is needed in order to complete the original format string. `RestFormat` is the remaining format string, `NChars` the number of characters scanned, and `InputStack` is the reversed list of inputs matched up to that point.

`{error, What}`

The read operation failed and the parameter `What` gives a hint about the error.

Example:

```
3> io_lib:fread("~f~f~f", "15.6 17.3e-6 24.5").
{ok,[15.6,1.73e-5,24.5],[]}
```

fread(Continuation, String, Format) -> Return

Types:

Continuation = see below

String = Format = string()

Return = {done, Result, LeftOverChars} | {more, Continuation}

Result = {ok, InputList} | eof | {error, What}

InputList = chars()

What = term()()

LeftOverChars = string()

This is the re-entrant formatted reader. The continuation of the first call to the functions must be `[]`. Refer to Armstrong, Virding, Williams, 'Concurrent Programming in Erlang', Chapter 13 for a complete description of how the re-entrant input scheme works.

The function returns:

`{done, Result, LeftOverChars}`

The input is complete. The result is one of the following:

`{ok, InputList}`

The string was read. `InputList` is the list of successfully matched and read items, and `LeftOverChars` are the remaining characters.

`eof`

End of file has been encountered. `LeftOverChars` are the input characters not used.

`{error, What}`

An error occurred and the parameter `What` gives a hint about the error.

`{more, Continuation}`

More data is required to build a term. `Continuation` must be passed to `fread/3`, when more data becomes available.

write_atom(Atom) -> chars()

Types:

Atom = atom()

Returns the list of characters needed to print the atom `Atom`.

write_string(String) -> chars()

Types:

String = string()

Returns the list of characters needed to print `String` as a string.

write_char(Integer) -> chars()

Types:

Integer = int()

Returns the list of characters needed to print a character constant in the ISO-latin-1 character set.

`indentation(String, StartIndent) -> int()`

Types:

String = string()

StartIndent = int()

Returns the indentation if `String` has been printed, starting at `StartIndent`.

`char_list(Term) -> bool()`

Types:

Term = term()

Returns `true` if `Term` is a flat list of characters in the ISO-latin-1 range, otherwise it returns `false`.

`deep_char_list(Term) -> bool()`

Types:

Term = term()

Returns `true` if `Term` is a, possibly deep, list of characters in the ISO-latin-1 range, otherwise it returns `false`.

`printable_list(Term) -> bool()`

Types:

Term = term()

Returns `true` if `Term` is a flat list of printable ISO-latin-1 characters, otherwise it returns `false`.

lib

Erlang module

> ### Warning:
>
> This module is retained for compatibility. It may disappear without warning in a future release.

Exports

```
flush_receive() ->
  receive
    _Any ->
      flush_receive()
  after
      0 ->
    ok
  end.
```

`flush_receive() -> void()`

Flushes the message buffer of the current process.

`error_message(Format, Args) -> ok`

Types:

Format = atom() | string() | binary()

Args = [term()]

Prints error message `Args` in accordance with `Format`. Similar to `io:format/2`, see *io(3)*.

`progname() -> atom()`

Returns the name of the script that started the current Erlang session.

`nonl(String1) -> String2`

Types:

String1 = String2 = string()

Removes the last newline character, if any, in `String1`.

`send(To, Msg)`

Types:

To = pid() | Name | {Name,Node}

Name = Node = atom()

Msg = term()

This function to makes it possible to send a message using the `apply/3` BIF.

`sendw(To, Msg)`

Types:

To = pid() | Name | {Name,Node}

Name = Node = atom()

Msg = term()

As `send/2`, but waits for an answer. It is implemented as follows:

```
sendw(To, Msg) ->
    To ! {self(),Msg},
    receive
        Reply -> Reply
    end.
```

The message returned is not necessarily a reply to the message sent.

lists

Erlang module

erlang的list 和 string 库实现 有我们自己实现函数的时候可以模仿的很多手法， 看熟悉了 容易在头脑里面自然反射代码的实现，俺学c的时候就是学会了字符串操作的时候感觉功力大增。 -余锋

This module contains functions for list processing.

Unless otherwise stated, all functions assume that position numbering starts at 1. That is, the first element of a list is at position 1.

Two terms `T1` and `T2` compare equal if `T1 == T2` evaluates to `true`. They match if `T1 =:= T2` evaluates to `true`.

Whenever an *ordering function* `F` is expected as argument, it is assumed that the following properties hold of `F` for all x, y and z:

- if x `F` y and y `F` x then x = y (`F` is antisymmetric);
- if x `F` y and y `F` z then x `F` z (`F` is transitive);
- x `F` y or y `F` x (`F` is total).

An example of a typical ordering function is less than or equal to, `=</2`.

Exports

(zen_latest@192.168.1.199)3> lists:all(fun(X) -> X >10 end ,[1,2,3]).
false

注意:
latest@192.168.1.199)4> lists:all(fun(X) -> X >10 end ,[]).
true

原因: all(Pred, []) when is_function(Pred, 1) -> true.

`all(Pred, List) -> bool()`

Types:

Pred = **fun(Elem) -> bool()**

Elem = **term()**

List = **[term()]**

Returns `true` if `Pred(Elem)` returns `true` for all elements `Elem` in `List`, otherwise `false`.

1> Y=lists:any(fun(Item)->io:format("~p",[Item]), Item>5 end,[1,5,17,56,3,12,4,2]).
2>1517true

同样要注意这里:
(zen_latest@192.168.1.199)4> lists:any(fun(X)-> X>10 end,[]).
false

(zen_latest@192.168.1.199)5> lists:any(fun(X)-> X>10 end,[11,2]).
true
(zen_latest@192.168.1.199)6> lists:any(fun(X)-> X>10 end,[1,2]).
false
(zen_latest@192.168.1.199)7>

`any(Pred, List) -> bool()`

Types:

Pred = **fun(Elem) -> bool()**

Elem = **term()**

List = **[term()]**

Returns `true` if `Pred(Elem)` returns `true` for at least one element `Elem` in `List`.

append([E]) -> E;
append([H|T]) -> H ++ append(T);
append([]) -> [].

里面只有一个元素的时候把结果取出来

`append(ListOfLists) -> List1`

Types:

ListOfLists = **[List]**

List = **List1** = **[term()]**

Returns a list in which all the sub-lists of `ListOfLists` have been appended. For example:

```
> lists:append([[1, 2, 3], [a, b], [4, 5, 6]]).
[1,2,3,a,b,4,5,6]
```

`append(List1, List2) -> List3`

Types:

List1 = List2 = List3 = [term()]

Returns a new list `List3` which is made from the elements of `List1` followed by the elements of `List2`. For example:

```
> lists:append("abc", "def").
"abcdef"
```

`lists:append(A, B)` is equivalent to `A ++ B`.

把各种数据连接在一起成为一个字符串

```
7> lists:concat([abc,123,[1232],"abc",<<"34">>]).
** exception error: no function clause matching
                    lists:thing_to_list(<<"34">>)
     in function  lists:flatmap/2
     in call from lists:flatmap/2
8> lists:concat([abc,123,[1232],"abc"]).
[97,98,99,49,50,51,1232,97,98,99]
```

`concat(Things) -> string()`

Types:

Things = [Thing]

Thing = atom() | integer() | float() | string()

Concatenates the text representation of the elements of `Things`. The elements of `Things` can be atoms, integers, floats or strings.

```
> lists:concat([doc, '/', file, '.', 3]).
"doc/file.3"
```

`delete(Elem, List1) -> List2`

```
9> lists:delete(1,[1,2,3,1,2,3,4]).
[2,3,1,2,3,4]
10> lists:delete(2,[1,2,3,1,2,3,4]).
[1,3,1,2,3,4]
11> lists:delete(7,[1,2,3,1,2,3,4]).
[1,2,3,1,2,3,4]
```

Types:

Elem = term()

List1 = List2 = [term()]

Returns a copy of `List1` where the first element matching `Elem` is deleted, if there is such an element.

`dropwhile(Pred, List1) -> List2`

Types:

Pred = fun(Elem) -> bool()

Elem = term()

List1 = List2 = [term()]

Drops elements `Elem` from `List1` while `Pred(Elem)` returns `true` and returns the remaining list.

`duplicate(N, Elem) -> List`

Types:

N = int()

Elem = term()

List = [term()]

Returns a list which contains N copies of the term `Elem`. For example:

```
> lists:duplicate(5, xx).
[xx,xx,xx,xx,xx]
```

filter(Pred, List1) -> List2

Types:

Pred = fun(Elem) -> bool()

Elem = term()

List1 = List2 = [term()]

`List2` is a list of all elements `Elem` in `List1` for which `Pred(Elem)` returns `true`.

flatlength(DeepList) -> int()

Types:

DeepList = [term() | DeepList]

Equivalent to `length(flatten(DeepList))`, but more efficient.

flatmap(Fun, List1) -> List2

Types:

Fun = fun(A) -> [B]

List1 = [A]

List2 = [B]

A = B = term()

Takes a function from As to lists of Bs, and a list of As (`List1`) and produces a list of Bs by applying the function to every element in `List1` and appending the resulting lists.

That is, `flatmap` behaves as if it had been defined as follows:

```
flatmap(Fun, List1) ->
    append(map(Fun, List1)).
```

Example:

```
> lists:flatmap(fun(X)->[X,X] end, [a,b,c]).
[a,a,b,b,c,c]
```

flatten(DeepList) -> List

Types:

DeepList = [term() | DeepList]

List = [term()]

Returns a flattened version of `DeepList`.

flatten(DeepList, Tail) -> List

Types:

DeepList = [term() | DeepList]

Tail = List = [term()]

Returns a flattened version of `DeepList` with the tail `Tail` appended.

foldl(Fun, Acc0, List) -> Acc1

Types:

Fun = fun(Elem, AccIn) -> AccOut

Elem = term()

Acc0 = Acc1 = AccIn = AccOut = term()

List = [term()]

Calls `Fun(Elem, AccIn)` on successive elements `A` of `List`, starting with `AccIn == Acc0`. `Fun/2` must return a new accumulator which is passed to the next call. The function returns the final value of the accumulator. `Acc0` is returned if the list is empty. For example:

```
> lists:foldl(fun(X, Sum) -> X + Sum end, 0, [1,2,3,4,5]).
15
> lists:foldl(fun(X, Prod) -> X * Prod end, 1, [1,2,3,4,5]).
120
```

foldr(Fun, Acc0, List) -> Acc1

Types:

Fun = fun(Elem, AccIn) -> AccOut

Elem = term()

Acc0 = Acc1 = AccIn = AccOut = term()

List = [term()]

Like `foldl/3`, but the list is traversed from right to left. For example:

```
> P = fun(A, AccIn) -> io:format("~p ", [A]), AccIn end.
#Fun<erl_eval.12.2225172>
> lists:foldl(P, void, [1,2,3]).
1 2 3 void
> lists:foldr(P, void, [1,2,3]).
3 2 1 void
```

`foldl/3` is tail recursive and would usually be preferred to `foldr/3`.

foreach(Fun, List) -> void()

Types:

Fun = fun(Elem) -> void()

Elem = term()

List = [term()]

Calls `Fun(Elem)` for each element `Elem` in `List`. This function is used for its side effects and the evaluation order is defined to be the same as the order of the elements in the list.

keydelete(Key, N, TupleList1) -> TupleList2

lists:keydelete(1,1,[{1,2},{1,3},{2,4}]).

[{1,3},{2,4}]

Types:

Key = term()

N = 1..tuple_size(Tuple)

TupleList1 = TupleList2 = [Tuple]

Tuple = tuple()

Returns a copy of `TupleList1` where the first occurrence of a tuple whose `N`th element compares equal to `Key` is deleted, if there is such a tuple.

[BIF] **`keyfind(Key, N, TupleList) -> Tuple | false`**

Types:

Key = term()

N = 1..tuple_size(Tuple)

TupleList = [Tuple]

Tuple = tuple()

```
1> lists:keyfind(zen,2,[{ok,lk},error,{this,zen},2012,{12,zen}]).
{this,zen}
2> lists:keyfind(zen,2,[{ok,lk},{23,zen},error,{this,zen},2012,{12,zen}]).
{23,zen}
3> lists:keyfind(zen,2,[1,2,3,4,5]).
false
4>
```

Searches the list of tuples `TupleList` for a tuple whose `N`th element compares equal to `Key`. Returns `Tuple` if such a tuple is found, otherwise `false`.

`keymap(Fun, N, TupleList1) -> TupleList2`

Types:

Fun = fun(Term1) -> Term2

Term1 = Term2 = term()

N = 1..tuple_size(Tuple)

TupleList1 = TupleList2 = [tuple()]

Returns a list of tuples where, for each tuple in `TupleList1`, the `N`th element `Term1` of the tuple has been replaced with the result of calling `Fun(Term1)`.

Examples:

```
> Fun = fun(Atom) -> atom_to_list(Atom) end.
#Fun<erl_eval.6.10732646>
2> lists:keymap(Fun, 2, [{name,jane,22},{name,lizzie,20},{name,lydia,15}]).
[{name,"jane",22},{name,"lizzie",20},{name,"lydia",15}]
```

`keymember(Key, N, TupleList) -> bool()`

Types:

Key = term()

N = 1..tuple_size(Tuple)

TupleList = [Tuple]

Tuple = tuple()

Returns `true` if there is a tuple in `TupleList` whose `N`th element compares equal to `Key`, otherwise `false`.

`keymerge(N, TupleList1, TupleList2) -> TupleList3`

Types:

N = 1..tuple_size(Tuple)

TupleList1 = TupleList2 = TupleList3 = [Tuple]

Tuple = tuple()

Returns the sorted list formed by merging `TupleList1` and `TupleList2`. The merge is performed on the `N`th element of each tuple. Both `TupleList1` and `TupleList2` must be key-sorted prior to evaluating this function. When two tuples compare equal, the tuple from `TupleList1` is picked before the tuple from `TupleList2`.

`keyreplace(Key, N, TupleList1, NewTuple) -> TupleList2`

Types:

Key = term()

N = 1..tuple_size(Tuple)

TupleList1 = TupleList2 = [Tuple]

NewTuple = Tuple = tuple()

Returns a copy of `TupleList1` where the first occurrence of a `T` tuple whose `N`th element compares equal to `Key` is replaced with `NewTuple`, if there is such a tuple `T`.

`keysearch(Key, N, TupleList) -> {value, Tuple} | false`

Types:

Key = term()

N = 1..tuple_size(Tuple)

TupleList = [Tuple]

Tuple = tuple()

Searches the list of tuples `TupleList` for a tuple whose `N`th element compares equal to `Key`. Returns `{value, Tuple}` if such a tuple is found, otherwise `false`.

> **Note:**
>
> This function is retained for backward compatibility. The function `lists:keyfind/3` (introduced in R13A) is in most cases more convenient.

`keysort(N, TupleList1) -> TupleList2`

Types: 实在没法喜欢这种函数,参数不看说明是猜不出来是做什么用的

N = 1..tuple_size(Tuple)

TupleList1 = TupleList2 = [Tuple]

Tuple = tuple()

```
randomize(List) ->
  D = lists:map(fun(A) ->
                {random:uniform(), A}
        end, List),
  {_, D1} = lists:unzip(lists:keysort(1, D)),
  D1.
```

Returns a list containing the sorted elements of the list `TupleList1`. Sorting is performed on the `N`th element of the tuples. The sort is stable.

`keystore(Key, N, TupleList1, NewTuple) -> TupleList2`

Types:

```
T2 =[{{1,2},1},{{2,3},23},{{3,1},12},{{3,2},24},{{1,1},34},{{2,4},12},{{1,1},99}].
lists:keysort(1,T2).
```

Key = term()

N = 1..tuple_size(Tuple)

TupleList1 = TupleList2 = [Tuple]

NewTuple = Tuple = tuple()

Returns a copy of `TupleList1` where the first occurrence of a tuple `T` whose `Nth` element compares equal to `Key` is replaced with `NewTuple`, if there is such a tuple `T`. If there is no such tuple `T` a copy of `TupleList1` where [`NewTuple`] has been appended to the end is returned.

`keytake(Key, N, TupleList1) -> {value, Tuple, TupleList2} | false`

Types:

Key = term()

N = 1..tuple_size(Tuple)

TupleList1 = TupleList2 = [Tuple]

Tuple = tuple()

Searches the list of tuples `TupleList1` for a tuple whose `Nth` element compares equal to `Key`. Returns `{value, Tuple, TupleList2}` if such a tuple is found, otherwise `false`. `TupleList2` is a copy of `TupleList1` where the first occurrence of `Tuple` has been removed.

`last(List) -> Last`

Types:

List = [term()], length(List) > 0

Last = term()

Returns the last element in `List`.

`map(Fun, List1) -> List2`

Types:

Fun = fun(A) -> B

List1 = [A]

List2 = [B]

A = B = term()

Takes a function from `As` to `Bs`, and a list of `As` and produces a list of `Bs` by applying the function to every element in the list. This function is used to obtain the return values. The evaluation order is implementation dependent.

`mapfoldl(Fun, Acc0, List1) -> {List2, Acc1}`

Types:

Fun = fun(A, AccIn) -> {B, AccOut}

Acc0 = Acc1 = AccIn = AccOut = term()

List1 = [A]

List2 = [B]

A = B = term()

`mapfoldl` combines the operations of `map/2` and `foldl/3` into one pass. An example, summing the elements in a list and double them at the same time:

```
> lists:mapfoldl(fun(X, Sum) -> {2*X, X+Sum} end,
0, [1,2,3,4,5]).
{[2,4,6,8,10],15}
```

mapfoldr(Fun, Acc0, List1) -> {List2, Acc1}

Types:

Fun = fun(A, AccIn) -> {B, AccOut}

Acc0 = Acc1 = AccIn = AccOut = term()

List1 = [A]

List2 = [B]

A = B = term()

`mapfoldr` combines the operations of `map/2` and `foldr/3` into one pass.

max(List) -> Max

Types:

List = [term()], length(List) > 0

Max = term()

Returns the first element of `List` that compares greater than or equal to all other elements of `List`.

member(Elem, List) -> bool()

Types:

Elem = term()

List = [term()]

Returns `true` if `Elem` matches some element of `List`, otherwise `false`.

merge(ListOfLists) -> List1

Types:

ListOfLists = [List]

List = List1 = [term()]

Returns the sorted list formed by merging all the sub-lists of `ListOfLists`. All sub-lists must be sorted prior to evaluating this function. When two elements compare equal, the element from the sub-list with the lowest position in `ListOfLists` is picked before the other element.

merge(List1, List2) -> List3

Types:

List1 = List2 = List3 = [term()]

Returns the sorted list formed by merging `List1` and `List2`. Both `List1` and `List2` must be sorted prior to evaluating this function. When two elements compare equal, the element from `List1` is picked before the element from `List2`.

merge(Fun, List1, List2) -> List3

Types:

Fun = fun(A, B) -> bool()

List1 = [A]

List2 = [B]

List3 = [A | B]

A = B = term()

Returns the sorted list formed by merging `List1` and `List2`. Both `List1` and `List2` must be sorted according to the *ordering function* `Fun` prior to evaluating this function. `Fun(A, B)` should return `true` if `A` compares less than or equal to `B` in the ordering, `false` otherwise. When two elements compare equal, the element from `List1` is picked before the element from `List2`.

`merge3(List1, List2, List3) -> List4`

Types:

List1 = List2 = List3 = List4 = [term()]

Returns the sorted list formed by merging `List1`, `List2` and `List3`. All of `List1`, `List2` and `List3` must be sorted prior to evaluating this function. When two elements compare equal, the element from `List1`, if there is such an element, is picked before the other element, otherwise the element from `List2` is picked before the element from `List3`.

`min(List) -> Min`

Types:

List = [term()], length(List) > 0

Min = term()

Returns the first element of `List` that compares less than or equal to all other elements of `List`.

`nth(N, List) -> Elem`

Types:

N = 1..length(List)

List = [term()]

Elem = term()

Returns the `N`th element of `List`. For example:

```
> lists:nth(3, [a, b, c, d, e]).
c
```

`nthtail(N, List1) -> Tail`

Types:

N = 0..length(List1)

List1 = Tail = [term()]

Returns the `N`th tail of `List`, that is, the sublist of `List` starting at `N+1` and continuing up to the end of the list. For example:

```
> lists:nthtail(3, [a, b, c, d, e]).
[d,e]
> tl(tl(tl([a, b, c, d, e]))).
[d,e]
> lists:nthtail(0, [a, b, c, d, e]).
[a,b,c,d,e]
> lists:nthtail(5, [a, b, c, d, e]).
[]
```

partition(Pred, List) -> {Satisfying, NonSatisfying}

Types:

Pred = fun(Elem) -> bool()

Elem = term()

List = Satisfying = NonSatisfying = [term()]

Partitions `List` into two lists, where the first list contains all elements for which `Pred(Elem)` returns `true`, and the second list contains all elements for which `Pred(Elem)` returns `false`.

Examples:

```
> lists:partition(fun(A) -> A rem 2 == 1 end, [1,2,3,4,5,6,7]).
{[1,3,5,7],[2,4,6]}
> lists:partition(fun(A) -> is_atom(A) end, [a,b,1,c,d,2,3,4,e]).
{[a,b,c,d,e],[1,2,3,4]}
```

See also `splitwith/2` for a different way to partition a list.

prefix(List1, List2) -> bool()

Types:

List1 = List2 = [term()]

Returns `true` if `List1` is a prefix of `List2`, otherwise `false`.

reverse(List1) -> List2

Types:

List1 = List2 = [term()]

Returns a list with the top level elements in `List1` in reverse order.

reverse(List1, Tail) -> List2

Types:

List1 = Tail = List2 = [term()]

Returns a list with the top level elements in `List1` in reverse order, with the tail `Tail` appended. For example:

```
> lists:reverse([1, 2, 3, 4], [a, b, c]).
[4,3,2,1,a,b,c]
```

seq(From, To) -> Seq

seq(From, To, Incr) -> Seq

Types:

From = To = Incr = int()

Seq = [int()]

Returns a sequence of integers which starts with `From` and contains the successive results of adding `Incr` to the previous element, until `To` has been reached or passed (in the latter case, `To` is not an element of the sequence). `Incr` defaults to 1.

Failure: If `To<From-Incr` and `Incr` is positive, or if `To>From-Incr` and `Incr` is negative, or if `Incr==0` and `From/=To`.

The following equalities hold for all sequences:

```
length(lists:seq(From, To)) == To-From+1
length(lists:seq(From, To, Incr)) == (To-From+Incr) div Incr
```

Examples:

```
> lists:seq(1, 10).
[1,2,3,4,5,6,7,8,9,10]
> lists:seq(1, 20, 3).
[1,4,7,10,13,16,19]
> lists:seq(1, 0, 1).
[]
> lists:seq(10, 6, 4).
[]
> lists:seq(1, 1, 0).
[1]
```

sort(List1) -> List2

Types:

List1 = List2 = [term()]

Returns a list containing the sorted elements of `List1`.

sort(Fun, List1) -> List2

Types:

Fun = fun(Elem1, Elem2) -> bool()

Elem1 = Elem2 = term()

List1 = List2 = [term()]

Returns a list containing the sorted elements of `List1`, according to the *ordering function* `Fun`. `Fun(A, B)` should return `true` if `A` compares less than or equal to `B` in the ordering, `false` otherwise.

split(N, List1) -> {List2, List3}

Types:

N = 0..length(List1)

List1 = List2 = List3 = [term()]

Splits `List1` into `List2` and `List3`. `List2` contains the first `N` elements and `List3` the rest of the elements (the `N`th tail).

%分组函数

splitwith(Pred, List) -> {List1, List2}

Types:

Pred = fun(Elem) -> bool()

Elem = term()

List = List1 = List2 = [term()]

Partitions `List` into two lists according to `Pred`. `splitwith/2` behaves as if it is defined as follows:

```
splitwith(Pred, List) ->
    {takewhile(Pred, List), dropwhile(Pred, List)}.
```

Examples:

```
> lists:splitwith(fun(A) -> A rem 2 == 1 end, [1,2,3,4,5,6,7]).
{[1],[2,3,4,5,6,7]}
> lists:splitwith(fun(A) -> is_atom(A) end, [a,b,1,c,d,2,3,4,e]).
{[a,b],[1,c,d,2,3,4,e]}
```

See also `partition/2` for a different way to partition a list.

sublist(List1, Len) -> List2

Types:

List1 = List2 = [term()]

Len = int()

Returns the sub-list of `List1` starting at position 1 and with (max) `Len` elements. It is not an error for `Len` to exceed the length of the list -- in that case the whole list is returned.

sublist(List1, Start, Len) -> List2

Types:

List1 = List2 = [term()]

Start = 1..(length(List1)+1)

Len = int()

Returns the sub-list of `List1` starting at `Start` and with (max) `Len` elements. It is not an error for `Start+Len` to exceed the length of the list.

```
> lists:sublist([1,2,3,4], 2, 2).
[2,3]
> lists:sublist([1,2,3,4], 2, 5).
[2,3,4]
> lists:sublist([1,2,3,4], 5, 2).
[]
```

subtract(List1, List2) -> List3

Types:

List1 = List2 = List3 = [term()]

Returns a new list `List3` which is a copy of `List1`, subjected to the following procedure: for each element in `List2`, its first occurrence in `List1` is deleted. For example:

```
> lists:subtract("123212", "212").
"312".
```

`lists:subtract(A, B)` is equivalent to `A -- B`.

> Warning:
>
> The complexity of `lists:subtract(A, B)` is proportional to `length(A)*length(B)`, meaning that it will be very slow if both `A` and `B` are long lists. (Using ordered lists and *ordsets:subtract/2* is a much better choice if both lists are long.)

`suffix(List1, List2) -> bool()`

Returns `true` if `List1` is a suffix of `List2`, otherwise `false`.

`sum(List) -> number()`

Types:

List = [number()]

Returns the sum of the elements in `List`.

`takewhile(Pred, List1) -> List2`

注意并不是想象中的那样 前面的不符合要求就不执行了

Types:

Pred = fun(Elem) -> bool()

Elem = term()

List1 = List2 = [term()]

Takes elements `Elem` from `List1` while `Pred(Elem)` returns `true`, that is, the function returns the longest prefix of the list for which all elements satisfy the predicate.

`ukeymerge(N, TupleList1, TupleList2) -> TupleList3`

Types:

N = 1..tuple_size(Tuple)

TupleList1 = TupleList2 = TupleList3 = [Tuple]

Tuple = tuple()

Returns the sorted list formed by merging `TupleList1` and `TupleList2`. The merge is performed on the `N`th element of each tuple. Both `TupleList1` and `TupleList2` must be key-sorted without duplicates prior to evaluating this function. When two tuples compare equal, the tuple from `TupleList1` is picked and the one from `TupleList2` deleted.

`ukeysort(N, TupleList1) -> TupleList2`

Types:

N = 1..tuple_size(Tuple)

TupleList1 = TupleList2 = [Tuple]

Tuple = tuple()

Returns a list containing the sorted elements of the list `TupleList1` where all but the first tuple of the tuples comparing equal have been deleted. Sorting is performed on the `N`th element of the tuples.

umerge(ListOfLists) -> List1

Types:

ListOfLists = [List]

List = List1 = [term()]

Returns the sorted list formed by merging all the sub-lists of `ListOfLists`. All sub-lists must be sorted and contain no duplicates prior to evaluating this function. When two elements compare equal, the element from the sub-list with the lowest position in `ListOfLists` is picked and the other one deleted.

umerge(List1, List2) -> List3

Types:

List1 = List2 = List3 = [term()]

Returns the sorted list formed by merging `List1` and `List2`. Both `List1` and `List2` must be sorted and contain no duplicates prior to evaluating this function. When two elements compare equal, the element from `List1` is picked and the one from `List2` deleted.

umerge(Fun, List1, List2) -> List3

Types:

Fun = fun(A, B) -> bool()

List1 = [A]

List2 = [B]

List3 = [A | B]

A = B = term()

Returns the sorted list formed by merging `List1` and `List2`. Both `List1` and `List2` must be sorted according to the *ordering function* `Fun` and contain no duplicates prior to evaluating this function. `Fun(A, B)` should return `true` if `A` compares less than or equal to `B` in the ordering, `false` otherwise. When two elements compare equal, the element from `List1` is picked and the one from `List2` deleted.

umerge3(List1, List2, List3) -> List4

Types:

List1 = List2 = List3 = List4 = [term()]

Returns the sorted list formed by merging `List1`, `List2` and `List3`. All of `List1`, `List2` and `List3` must be sorted and contain no duplicates prior to evaluating this function. When two elements compare equal, the element from `List1` is picked if there is such an element, otherwise the element from `List2` is picked, and the other one deleted.

unzip(List1) -> {List2, List3}

Types:

List1 = [{X, Y}]

List2 = [X]

List3 = [Y]

X = Y = term()

"Unzips" a list of two-tuples into two lists, where the first list contains the first element of each tuple, and the second list contains the second element of each tuple.

unzip3(List1) -> {List2, List3, List4}

Types:

List1 = [{X, Y, Z}]

List2 = [X]

List3 = [Y]

List4 = [Z]

X = Y = Z = term()

"Unzips" a list of three-tuples into three lists, where the first list contains the first element of each tuple, the second list contains the second element of each tuple, and the third list contains the third element of each tuple.

usort(List1) -> List2

Types:

List1 = List2 = [term()]

Returns a list containing the sorted elements of `List1` where all but the first element of the elements comparing equal have been deleted.

usort(Fun, List1) -> List2

Types:

Fun = fun(Elem1, Elem2) -> bool()

Elem1 = Elem2 = term()

List1 = List2 = [term()]

```
1> lists:usort([2,1,3,4,2,3,8,9,a,v,c]).
[1,2,3,4,8,9,a,c,v]
2> lists:usort([2,1,3,4,2,3,8,9,[a,c,q,m,x],a,v,c]).
[1,2,3,4,8,9,a,c,v,[a,c,q,m,x]]
3> lists:usort([2,1,3,4,2,3,8,9,[0,c,q,m,x],a,v,c]).
[1,2,3,4,8,9,a,c,v,[0,c,q,m,x]]
4> lists:usort([2,1,3,{zen},4,2,3,8,9,[0,c,q,m,x],a,v,c]).
[1,2,3,4,8,9,a,c,v,{zen},[0,c,q,m,x]]
```

Returns a list which contains the sorted elements of `List1` where all but the first element of the elements comparing equal according to the *ordering function* `Fun` have been deleted. `Fun(A, B)` should return `true` if `A` compares less than or equal to `B` in the ordering, `false` otherwise.

zip(List1, List2) -> List3

Types:

List1 = [X]

List2 = [Y]

List3 = [{X, Y}]

X = Y = term()

两个等长的list压缩到当前的二元元组的数组中

```
ZL=lists:zip([1,2,4,54,56567,1,2],[1,2,3,4,5,6,67]).
结果:
[{1,1},{2,2},{4,3},{54,4},{56567,5},{1,6},{2,67}]
```

"Zips" two lists of equal length into one list of two-tuples, where the first element of each tuple is taken from the first list and the second element is taken from corresponding element in the second list.

zip3(List1, List2, List3) -> List4

Types:

List1 = [X]

List2 = [Y]

List3 = [Z]

List3 = [{X, Y, Z}]

X = Y = Z = term()

"Zips" three lists of equal length into one list of three-tuples, where the first element of each tuple is taken from the first list, the second element is taken from corresponding element in the second list, and the third element is taken from the corresponding element in the third list.

zipwith(Combine, List1, List2) -> List3

Types:

Combine = fun(X, Y) -> T

List1 = [X]

List2 = [Y]

List3 = [T]

X = Y = T = term()

Combine the elements of two lists of equal length into one list. For each pair `X, Y` of list elements from the two lists, the element in the result list will be `Combine(X, Y)`.

`zipwith(fun(X, Y) -> {X,Y} end, List1, List2)` is equivalent to `zip(List1, List2)`.

Example:

```
> lists:zipwith(fun(X, Y) -> X+Y end, [1,2,3], [4,5,6]).
[5,7,9]
```

zipwith3(Combine, List1, List2, List3) -> List4

Types:

Combine = fun(X, Y, Z) -> T

List1 = [X]

List2 = [Y]

List3 = [Z]

List4 = [T]

X = Y = Z = T = term()

Combine the elements of three lists of equal length into one list. For each triple `X, Y, Z` of list elements from the three lists, the element in the result list will be `Combine(X, Y, Z)`.

`zipwith3(fun(X, Y, Z) -> {X,Y,Z} end, List1, List2, List3)` is equivalent to `zip3(List1, List2, List3)`.

Examples:

```
> lists:zipwith3(fun(X, Y, Z) -> X+Y+Z end, [1,2,3], [4,5,6], [7,8,9]).
[12,15,18]
> lists:zipwith3(fun(X, Y, Z) -> [X,Y,Z] end, [a,b,c], [x,y,z], [1,2,3]).
[[a,x,1],[b,y,2],[c,z,3]]
```

log_mf_h

Erlang module

The `log_mf_h` is a `gen_event` handler module which can be installed in any `gen_event` process. It logs onto disk all events which are sent to an event manager. Each event is written as a binary which makes the logging very fast. However, a tool such as the `Report Browser` (`rb`) must be used in order to read the files. The events are written to multiple files. When all files have been used, the first one is re-used and overwritten. The directory location, the number of files, and the size of each file are configurable. The directory will include one file called `index`, and report files `1, 2, ....`.

Exports

`init(Dir, MaxBytes, MaxFiles)`

`init(Dir, MaxBytes, MaxFiles, Pred) -> Args`

Types:

Dir = string()

MaxBytes = integer()

MaxFiles = 0 < integer() < 256

Pred = fun(Event) -> boolean()

Event = term()

Args = args()

Initiates the event handler. This function returns `Args`, which should be used in a call to `gen_event:add_handler(EventMgr, log_mf_h, Args)`.

`Dir` specifies which directory to use for the log files. `MaxBytes` specifies the size of each individual file. `MaxFiles` specifies how many files are used. `Pred` is a predicate function used to filter the events. If no predicate function is specified, all events are logged.

See Also

gen_event(3), rb(3)

math

Erlang module

This module provides an interface to a number of mathematical functions.

> ### Note:
>
> Not all functions are implemented on all platforms. In particular, the `erf/1` and `erfc/1` functions are not implemented on Windows.

Exports

`pi() -> float()`

A useful number.

sin(X)

cos(X)

tan(X)

asin(X)

acos(X)

atan(X)

atan2(Y, X)

sinh(X)

cosh(X)

tanh(X)

asinh(X)

acosh(X)

atanh(X)

exp(X)

log(X)

log10(X)

pow(X, Y)

sqrt(X)

Types:

X = Y = number()

A collection of math functions which return floats. Arguments are numbers.

erf(X) -> float()

Types:

X = number()

Returns the error function of X, where

```
erf(X) = 2/sqrt(pi)*integral from 0 to X of exp(-t*t) dt.
```

erfc(X) -> float()

Types:

X = number()

`erfc(X)` returns `1.0 - erf(X)`, computed by methods that avoid cancellation for large `X`.

Bugs

As these are the C library, the bugs are the same.

ms_transform

Erlang module

This module implements the parse_transform that makes calls to `ets` and `dbg:fun2ms/1` translate into literal match specifications. It also implements the back end for the same functions when called from the Erlang shell.

The translations from fun's to match_specs is accessed through the two "pseudo functions" `ets:fun2ms/1` and `dbg:fun2ms/1`.

Actually this introduction is more or less an introduction to the whole concept of match specifications. Since everyone trying to use `ets:select` or `dbg` seems to end up reading this page, it seems in good place to explain a little more than just what this module does.

There are some caveats one should be aware of, please read through the whole manual page if it's the first time you're using the transformations.

Match specifications are used more or less as filters. They resemble usual Erlang matching in a list comprehension or in a `fun` used in conjunction with `lists:foldl` etc. The syntax of pure match specifications is somewhat awkward though, as they are made up purely by Erlang terms and there is no syntax in the language to make the match specifications more readable.

As the match specifications execution and structure is quite like that of a fun, it would for most programmers be more straight forward to simply write it using the familiar fun syntax and having that translated into a match specification automatically. Of course a real fun is more powerful than the match specifications allow, but bearing the match specifications in mind, and what they can do, it's still more convenient to write it all as a fun. This module contains the code that simply translates the fun syntax into match_spec terms.

Let's start with an ets example. Using `ets:select` and a match specification, one can filter out rows of a table and construct a list of tuples containing relevant parts of the data in these rows. Of course one could use `ets:foldl` instead, but the select call is far more efficient. Without the translation, one has to struggle with writing match specifications terms to accommodate this, or one has to resort to the less powerful `ets:match(_object)` calls, or simply give up and use the more inefficient method of `ets:foldl`. Using the `ets:fun2ms` transformation, a `ets:select` call is at least as easy to write as any of the alternatives.

As an example, consider a simple table of employees:

```
-record(emp, {empno,     %Employee number as a string, the key
              surname,   %Surname of the employee
              givenname, %Given name of employee
              dept,      %Department one of {dev,sales,prod,adm}
              empyear}). %Year the employee was employed
```

We create the table using:

```
ets:new(emp_tab,[{keypos,#emp.empno},named_table,ordered_set]).
```

Let's also fill it with some randomly chosen data for the examples:

```
[{emp,"011103","Black","Alfred",sales,2000},
 {emp,"041231","Doe","John",prod,2001},
 {emp,"052341","Smith","John",dev,1997},
 {emp,"076324","Smith","Ella",sales,1995},
```

```
 {emp,"122334","Weston","Anna",prod,2002},
 {emp,"535216","Chalker","Samuel",adm,1998},
 {emp,"789789","Harrysson","Joe",adm,1996},
 {emp,"963721","Scott","Juliana",dev,2003},
 {emp,"989891","Brown","Gabriel",prod,1999}]
```

Now, the amount of data in the table is of course to small to justify complicated ets searches, but on real tables, using `select` to get exactly the data you want will increase efficiency remarkably.

Lets say for example that we'd want the employee numbers of everyone in the sales department. One might use `ets:match` in such a situation:

```
1> ets:match(emp_tab, {'_', '$1', '_', '_', sales, '_'}).
[["011103"],["076324"]]
```

Even though `ets:match` does not require a full match specification, but a simpler type, it's still somewhat unreadable, and one has little control over the returned result, it's always a list of lists. OK, one might use `ets:foldl` or `ets:foldr` instead:

```
ets:foldr(fun(#emp{empno = E, dept = sales},Acc) -> [E | Acc];
             (_,Acc) -> Acc
          end,
          [],
          emp_tab).
```

Running that would result in `["011103","076324"]`, which at least gets rid of the extra lists. The fun is also quite straightforward, so the only problem is that all the data from the table has to be transferred from the table to the calling process for filtering. That's inefficient compared to the `ets:match` call where the filtering can be done "inside" the emulator and only the result is transferred to the process. Remember that ets tables are all about efficiency, if it wasn't for efficiency all of ets could be implemented in Erlang, as a process receiving requests and sending answers back. One uses ets because one wants performance, and therefore one wouldn't want all of the table transferred to the process for filtering. OK, let's look at a pure `ets:select` call that does what the `ets:foldr` does:

```
ets:select(emp_tab,[{#emp{empno = '$1', dept = sales, _='_'},[],['$1']}]).
```

Even though the record syntax is used, it's still somewhat hard to read and even harder to write. The first element of the tuple, `#emp{empno = '$1', dept = sales, _='_'}` tells what to match, elements not matching this will not be returned at all, as in the `ets:match` example. The second element, the empty list is a list of guard expressions, which we need none, and the third element is the list of expressions constructing the return value (in ets this almost always is a list containing one single term). In our case `'$1'` is bound to the employee number in the head (first element of tuple), and hence it is the employee number that is returned. The result is `["011103","076324"]`, just as in the `ets:foldr` example, but the result is retrieved much more efficiently in terms of execution speed and memory consumption.

We have one efficient but hardly readable way of doing it and one inefficient but fairly readable (at least to the skilled Erlang programmer) way of doing it. With the use of `ets:fun2ms`, one could have something that is as efficient as possible but still is written as a filter using the fun syntax:

```
-include_lib("stdlib/include/ms_transform.hrl").
```

```
% ...

ets:select(emp_tab, ets:fun2ms(
                      fun(#emp{empno = E, dept = sales}) ->
                              E
                      end)).
```

This may not be the shortest of the expressions, but it requires no special knowledge of match specifications to read. The fun's head should simply match what you want to filter out and the body returns what you want returned. As long as the fun can be kept within the limits of the match specifications, there is no need to transfer all data of the table to the process for filtering as in the `ets:foldr` example. In fact it's even easier to read then the `ets:foldr` example, as the select call in itself discards anything that doesn't match, while the fun of the `foldr` call needs to handle both the elements matching and the ones not matching.

It's worth noting in the above `ets:fun2ms` example that one needs to include `ms_transform.hrl` in the source code, as this is what triggers the parse transformation of the `ets:fun2ms` call to a valid match specification. This also implies that the transformation is done at compile time (except when called from the shell of course) and therefore will take no resources at all in runtime. So although you use the more intuitive fun syntax, it gets as efficient in runtime as writing match specifications by hand.

Let's look at some more `ets` examples. Let's say one wants to get all the employee numbers of any employee hired before the year 2000. Using `ets:match` isn't an alternative here as relational operators cannot be expressed there. Once again, an `ets:foldr` could do it (slowly, but correct):

```
ets:foldr(fun(#emp{empno = E, empyear = Y},Acc) when Y < 2000 -> [E | Acc];
                  (_,Acc) -> Acc
          end,
          [],
          emp_tab).
```

The result will be `["052341","076324","535216","789789","989891"]`, as expected. Now the equivalent expression using a handwritten match specification would look something like this:

```
ets:select(emp_tab,[{#emp{empno = '$1', empyear = '$2', _='_'},
                     [{'<', '$2', 2000}],
                     ['$1']}]).
```

This gives the same result, the `[{'<', '$2', 2000}]` is in the guard part and therefore discards anything that does not have a empyear (bound to '$2' in the head) less than 2000, just as the guard in the `foldl` example. Lets jump on to writing it using `ets:fun2ms`

```
-include_lib("stdlib/include/ms_transform.hrl").

% ...

ets:select(emp_tab, ets:fun2ms(
                      fun(#emp{empno = E, empyear = Y}) when Y < 2000 ->
                              E
                      end)).
```

Obviously readability is gained by using the parse transformation.

I'll show some more examples without the tiresome comparing-to-alternatives stuff. Let's say we'd want the whole object matching instead of only one element. We could of course assign a variable to every part of the record and build it up once again in the body of the `fun`, but it's easier to do like this:

```
ets:select(emp_tab, ets:fun2ms(
                      fun(Obj = #emp{empno = E, empyear = Y})
                         when Y < 2000 ->
                              Obj
                      end)).
```

Just as in ordinary Erlang matching, you can bind a variable to the whole matched object using a "match in then match", i.e. a `=`. Unfortunately this is not general in `fun's` translated to match specifications, only on the "top level", i.e. matching the *whole* object arriving to be matched into a separate variable, is it allowed. For the one's used to writing match specifications by hand, I'll have to mention that the variable A will simply be translated into '$_'. It's not general, but it has very common usage, why it is handled as a special, but useful, case. If this bothers you, the pseudo function `object` also returns the whole matched object, see the part about caveats and limitations below.

Let's do something in the `fun`'s body too: Let's say that someone realizes that there are a few people having an employee number beginning with a zero (`0`), which shouldn't be allowed. All those should have their numbers changed to begin with a one (`1`) instead and one wants the list `[{<Old empno>,<New empno>}]` created:

```
ets:select(emp_tab, ets:fun2ms(
                      fun(#emp{empno = [$0 | Rest] }) ->
                              {[$0|Rest],[$1|Rest]}
                      end)).
```

As a matter of fact, this query hits the feature of partially bound keys in the table type `ordered_set`, so that not the whole table need be searched, only the part of the table containing keys beginning with `0` is in fact looked into.

The fun of course can have several clauses, so that if one could do the following: For each employee, if he or she is hired prior to 1997, return the tuple `{inventory, <employee number>}`, for each hired 1997 or later, but before 2001, return `{rookie, <employee number>}`, for all others return `{newbie, <employee number>}`. All except for the ones named `Smith` as they would be affronted by anything other than the tag `guru` and that is also what's returned for their numbers; `{guru, <employee number>}`:

```
ets:select(emp_tab, ets:fun2ms(
                      fun(#emp{empno = E, surname = "Smith" }) ->
                              {guru,E};
                         (#emp{empno = E, empyear = Y}) when Y < 1997  ->
                              {inventory, E};
                         (#emp{empno = E, empyear = Y}) when Y > 2001  ->
                              {newbie, E};
                         (#emp{empno = E, empyear = Y}) -> % 1997 -- 2001
                              {rookie, E}
                      end)).
```

The result will be:

```
[{rookie,"011103"},
 {rookie,"041231"},
 {guru,"052341"},
 {guru,"076324"},
```

```
{newbie,"122334"},
{rookie,"535216"},
{inventory,"789789"},
{newbie,"963721"},
{rookie,"989891"}]
```

and so the Smith's will be happy...

So, what more can you do? Well, the simple answer would be; look in the documentation of match specifications in ERTS users guide. However let's briefly go through the most useful "built in functions" that you can use when the `fun` is to be translated into a match specification by `ets:fun2ms` (it's worth mentioning, although it might be obvious to some, that calling other functions than the one's allowed in match specifications cannot be done. No "usual" Erlang code can be executed by the `fun` being translated by `fun2ms`, the `fun` is after all limited exactly to the power of the match specifications, which is unfortunate, but the price one has to pay for the execution speed of an `ets:select` compared to `ets:foldl/foldr`).

The head of the `fun` is obviously a head matching (or mismatching) *one* parameter, one object of the table we `select` from. The object is always a single variable (can be _) or a tuple, as that's what's in `ets, dets` and `mnesia` tables (the match specification returned by `ets:fun2ms` can of course be used with `dets:select` and `mnesia:select` as well as with `ets:select`). The use of `=` in the head is allowed (and encouraged) on the top level.

The guard section can contain any guard expression of Erlang. Even the "old" type test are allowed on the toplevel of the guard (`integer(X)` instead of `is_integer(X)`). As the new type tests (the `is_` tests) are in practice just guard bif's they can also be called from within the body of the fun, but so they can in ordinary Erlang code. Also arithmetics is allowed, as well as ordinary guard bif's. Here's a list of bif's and expressions:

- The type tests: is_atom, is_constant, is_float, is_integer, is_list, is_number, is_pid, is_port, is_reference, is_tuple, is_binary, is_function, is_record
- The boolean operators: not, and, or, andalso, orelse
- The relational operators: >, >=, <, =<, =:=, ==, =/=, /=
- Arithmetics: +, -, *, div, rem
- Bitwise operators: band, bor, bxor, bnot, bsl, bsr
- The guard bif's: abs, element, hd, length, node, round, size, tl, trunc, self
- The obsolete type test (only in guards): atom, constant, float, integer, list, number, pid, port, reference, tuple, binary, function, record

Contrary to the fact with "handwritten" match specifications, the `is_record` guard works as in ordinary Erlang code.

Semicolons (`;`) in guards are allowed, the result will be (as expected) one "match_spec-clause" for each semicolon-separated part of the guard. The semantics being identical to the Erlang semantics.

The body of the `fun` is used to construct the resulting value. When selecting from tables one usually just construct a suiting term here, using ordinary Erlang term construction, like tuple parentheses, list brackets and variables matched out in the head, possibly in conjunction with the occasional constant. Whatever expressions are allowed in guards are also allowed here, but there are no special functions except `object` and `bindings` (see further down), which returns the whole matched object and all known variable bindings respectively.

The `dbg` variants of match specifications have an imperative approach to the match specification body, the ets dialect hasn't. The fun body for `ets:fun2ms` returns the result without side effects, and as matching (`=`) in the body of the match specifications is not allowed (for performance reasons) the only thing left, more or less, is term construction...

Let's move on to the `dbg` dialect, the slightly different match specifications translated by `dbg:fun2ms`.

The same reasons for using the parse transformation applies to `dbg`, maybe even more so as filtering using Erlang code is simply not a good idea when tracing (except afterwards, if you trace to file). The concept is similar to that of `ets:fun2ms` except that you usually use it directly from the shell (which can also be done with `ets:fun2ms`).

Let's manufacture a toy module to trace on

```
-module(toy).

-export([start/1, store/2, retrieve/1]).

start(Args) ->
    toy_table = ets:new(toy_table,Args).

store(Key, Value) ->
    ets:insert(toy_table,{Key,Value}).

retrieve(Key) ->
    [{Key, Value}] = ets:lookup(toy_table,Key),
    Value.
```

During model testing, the first test bails out with a `{badmatch,16}` in `{toy,start,1}`, why?

We suspect the ets call, as we match hard on the return value, but want only the particular `new` call with `toy_table` as first parameter. So we start a default tracer on the node:

```
1> dbg:tracer().
{ok,<0.88.0>}
```

And so we turn on call tracing for all processes, we are going to make a pretty restrictive trace pattern, so there's no need to call trace only a few processes (it usually isn't):

```
2> dbg:p(all,call).
{ok,[{matched,nonode@nohost,25}]}
```

It's time to specify the filter. We want to view calls that resemble `ets:new(toy_table,<something>)`:

```
3> dbg:tp(ets,new,dbg:fun2ms(fun([toy_table,_]) -> true end)).
{ok,[{matched,nonode@nohost,1},{saved,1}]}
```

As can be seen, the `fun`'s used with `dbg:fun2ms` takes a single list as parameter instead of a single tuple. The list matches a list of the parameters to the traced function. A single variable may also be used of course. The body of the fun expresses in a more imperative way actions to be taken if the fun head (and the guards) matches. I return `true` here, but it's only because the body of a fun cannot be empty, the return value will be discarded.

When we run the test of our module now, we get the following trace output:

```
(<0.86.0>) call ets:new(toy_table,[ordered_set])
```

Let's play we haven't spotted the problem yet, and want to see what `ets:new` returns. We do a slightly different trace pattern:

```
4> dbg:tp(ets,new,dbg:fun2ms(fun([toy_table,_]) -> return_trace() end)).
```

Resulting in the following trace output when we run the test:

```
(<0.86.0>) call ets:new(toy_table,[ordered_set])
(<0.86.0>) returned from ets:new/2 -> 24
```

The call to `return_trace`, makes a trace message appear when the function returns. It applies only to the specific function call triggering the match specification (and matching the head/guards of the match specification). This is the by far the most common call in the body of a `dbg` match specification.

As the test now fails with `{badmatch,24}`, it's obvious that the badmatch is because the atom `toy_table` does not match the number returned for an unnamed table. So we spotted the problem, the table should be named and the arguments supplied by our test program does not include `named_table`. We rewrite the start function to:

```
start(Args) ->
    toy_table = ets:new(toy_table,[named_table |Args]).
```

And with the same tracing turned on, we get the following trace output:

```
(<0.86.0>) call ets:new(toy_table,[named_table,ordered_set])
(<0.86.0>) returned from ets:new/2 -> toy_table
```

Very well. Let's say the module now passes all testing and goes into the system. After a while someone realizes that the table `toy_table` grows while the system is running and that for some reason there are a lot of elements with atom's as keys. You had expected only integer keys and so does the rest of the system. Well, obviously not all of the system. You turn on call tracing and try to see calls to your module with an atom as the key:

```
1> dbg:tracer().
{ok,<0.88.0>}
2> dbg:p(all,call).
{ok,[{matched,nonode@nohost,25}]}
3> dbg:tpl(toy,store,dbg:fun2ms(fun([A,_]) when is_atom(A) -> true end)).
{ok,[{matched,nonode@nohost,1},{saved,1}]}
```

We use `dbg:tpl` here to make sure to catch local calls (let's say the module has grown since the smaller version and we're not sure this inserting of atoms is not done locally...). When in doubt always use local call tracing.

Let's say nothing happens when we trace in this way. Our function is never called with these parameters. We make the conclusion that someone else (some other module) is doing it and we realize that we must trace on ets:insert and want to see the calling function. The calling function may be retrieved using the match specification function `caller` and to get it into the trace message, one has to use the match spec function `message`. The filter call looks like this (looking for calls to `ets:insert`):

```
4> dbg:tpl(ets,insert,dbg:fun2ms(fun([toy_table,{A,_}]) when is_atom(A) ->
                                     message(caller())
                                 end)).
{ok,[{matched,nonode@nohost,1},{saved,2}]}
```

The caller will now appear in the "additional message" part of the trace output, and so after a while, the following output comes:

```
(<0.86.0>) call ets:insert(toy_table,{garbage,can}) ({evil_mod,evil_fun,2})
```

You have found out that the function `evil_fun` of the module `evil_mod`, with arity `2`, is the one causing all this trouble.

This was just a toy example, but it illustrated the most used calls in match specifications for `dbg` The other, more esotheric calls are listed and explained in the *Users guide of the ERTS application*, they really are beyond the scope of this document.

To end this chatty introduction with something more precise, here follows some parts about caveats and restrictions concerning the fun's used in conjunction with `ets:fun2ms` and `dbg:fun2ms`:

Warning:

To use the pseudo functions triggering the translation, one *has to* include the header file `ms_transform.hrl` in the source code. Failure to do so will possibly result in runtime errors rather than compile time, as the expression may be valid as a plain Erlang program without translation.

Warning:

The `fun` has to be literally constructed inside the parameter list to the pseudo functions. The `fun` cannot be bound to a variable first and then passed to `ets:fun2ms` or `dbg:fun2ms`, i.e this will work: `ets:fun2ms(fun(A) -> A end)` but not this: `F = fun(A) -> A end, ets:fun2ms(F)`. The later will result in a compile time error if the header is included, otherwise a runtime error. Even if the later construction would ever appear to work, it really doesn't, so don't ever use it.

Several restrictions apply to the fun that is being translated into a match_spec. To put it simple you cannot use anything in the fun that you cannot use in a match_spec. This means that, among others, the following restrictions apply to the fun itself:

- Functions written in Erlang cannot be called, neither local functions, global functions or real fun's
- Everything that is written as a function call will be translated into a match_spec call to a builtin function, so that the call `is_list(X)` will be translated to `{'is_list', '$1'}` (`'$1'` is just an example, the numbering may vary). If one tries to call a function that is not a match_spec builtin, it will cause an error.
- Variables occurring in the head of the `fun` will be replaced by match_spec variables in the order of occurrence, so that the fragment `fun({A,B,C})` will be replaced by `{'$1', '$2', '$3'}` etc. Every occurrence of such a variable later in the match_spec will be replaced by a match_spec variable in the same way, so that the fun `fun({A,B}) when is_atom(A) -> B end` will be translated into `[{{'$1','$2'}, [{is_atom,'$1'}],['$2']}]`.
- Variables that are not appearing in the head are imported from the environment and made into match_spec `const` expressions. Example from the shell:

```
1> X = 25.
25
2> ets:fun2ms(fun({A,B}) when A > X -> B end).
[{{'$1','$2'},[{'>','$1',{const,25}}],['$2']}]
```

- Matching with = cannot be used in the body. It can only be used on the top level in the head of the fun. Example from the shell again:

```
1> ets:fun2ms(fun({A,[B|C]} = D) when A > B -> D end).
[{{'$1',['$2'|'$3']},[{'>','$1','$2'}],['$_']}]
2> ets:fun2ms(fun({A,[B|C]=D}) when A > B -> D end).
Error: fun with head matching ('=' in head) cannot be translated into
match_spec
{error,transform_error}
3> ets:fun2ms(fun({A,[B|C]}) when A > B -> D = [B|C], D end).
Error: fun with body matching ('=' in body) is illegal as match_spec
{error,transform_error}
```

  All variables are bound in the head of a match_spec, so the translator can not allow multiple bindings. The special case when matching is done on the top level makes the variable bind to `'$_'` in the resulting match_spec, it is to allow a more natural access to the whole matched object. The pseudo function `object()` could be used instead, see below. The following expressions are translated equally:

```
ets:fun2ms(fun({a,_} = A) -> A end).
ets:fun2ms(fun({a,_}) -> object() end).
```

- The special match_spec variables `'$_'` and `'$*'` can be accessed through the pseudo functions `object()` (for `'$_'`) and `bindings()` (for `'$*'`). as an example, one could translate the following `ets:match_object/2` call to a `ets:select` call:

```
ets:match_object(Table, {'$1',test,'$2'}).
```

  ...is the same as...

```
ets:select(Table, ets:fun2ms(fun({A,test,B}) -> object() end)).
```

  (This was just an example, in this simple case the former expression is probably preferable in terms of readability). The `ets:select/2` call will conceptually look like this in the resulting code:

```
ets:select(Table, [{{'$1',test,'$2'},[],['$_']}]).
```

  Matching on the top level of the fun head might feel like a more natural way to access `'$_'`, see above.
- Term constructions/literals are translated as much as is needed to get them into valid match_specs, so that tuples are made into match_spec tuple constructions (a one element tuple containing the tuple) and constant expressions are used when importing variables from the environment. Records are also translated into plain tuple constructions, calls to element etc. The guard test `is_record/2` is translated into match_spec code using the three parameter version that's built into match_specs, so that `is_record(A,t)` is translated into `{is_record,'$1',t,5}` given that the record size of record type `t` is 5.
- Language constructions like `case`, `if`, `catch` etc that are not present in match_specs are not allowed.
- If the header file `ms_transform.hrl` is not included, the fun won't be translated, which may result in a *runtime error* (depending on if the fun is valid in a pure Erlang context). Be absolutely sure that the header is included when using `ets` and `dbg:fun2ms/1` in compiled code.

- If the pseudo function triggering the translation is `ets:fun2ms/1`, the fun's head must contain a single variable or a single tuple. If the pseudo function is `dbg:fun2ms/1` the fun's head must contain a single variable or a single list.

The translation from fun's to match_specs is done at compile time, so runtime performance is not affected by using these pseudo functions. The compile time might be somewhat longer though.

For more information about match_specs, please read about them in *ERTS users guide*.

Exports

`parse_transform(Forms,_Options) -> Forms`

Types:

Forms = Erlang abstract code format, see the erl_parse module description

_Options = Option list, required but not used

Implements the actual transformation at compile time. This function is called by the compiler to do the source code transformation if and when the `ms_transform.hrl` header file is included in your source code. See the `ets` and `dbg`:`fun2ms/1` function manual pages for documentation on how to use this parse_transform, see the `match_spec` chapter in `ERTS` users guide for a description of match specifications.

`transform_from_shell(Dialect,Clauses,BoundEnvironment) -> term()`

Types:

Dialect = ets | dbg

Clauses = Erlang abstract form for a single fun

BoundEnvironment = [{atom(), term()}, ...], list of variable bindings in the shell environment

Implements the actual transformation when the `fun2ms` functions are called from the shell. In this case the abstract form is for one single fun (parsed by the Erlang shell), and all imported variables should be in the key-value list passed as `BoundEnvironment`. The result is a term, normalized, i.e. not in abstract format.

`format_error(Errcode) -> ErrMessage`

Types:

Errcode = term()

ErrMessage = string()

Takes an error code returned by one of the other functions in the module and creates a textual description of the error. Fairly uninteresting function actually.

orddict

Erlang module

orddict 用按 key 排序的列表表示字典结构，数据量较大时访问开销很大，只适用于数据量较小的情况；
例如：向 orddict 顺序插入 z->1、a->2、c->3 后，orddict 结构为：[{a,2}, {c,3}, {z,1}]

`Orddict` implements a `Key` - `Value` dictionary. An `orddict` is a representation of a dictionary, where a list of pairs is used to store the keys and values. The list is ordered after the keys.

This module provides exactly the same interface as the module `dict` but with a defined representation. One difference is that while `dict` considers two keys as different if they do not match (`=:=`), this module considers two keys as different if and only if they do not compare equal (`==`).

DATA TYPES

```
ordered_dictionary()
  as returned by new/0
```

Exports

append(Key, Value, Orddict1) -> Orddict2

Types:

Key = Value = term()

Orddict1 = Orddict2 = ordered_dictionary()

This function appends a new `Value` to the current list of values associated with `Key`. An exception is generated if the initial value associated with `Key` is not a list of values.

append_list(Key, ValList, Orddict1) -> Orddict2

Types:

ValList = [Value]

Key = Value = term()

Orddict1 = Orddict2 = ordered_dictionary()

This function appends a list of values `ValList` to the current list of values associated with `Key`. An exception is generated if the initial value associated with `Key` is not a list of values.

将ValList追加到Key的Value. 如果Key的原有Value不是List会引发异常.

erase(Key, Orddict1) -> Orddict2

Types:

Key = term()

Orddict1 = Orddict2 = ordered_dictionary()

This function erases all items with a given key from a dictionary.

fetch(Key, Orddict) -> Value

Types:

Key = Value = term()

Orddict = ordered_dictionary()

This function returns the value associated with `Key` in the dictionary `Orddict`. `fetch` assumes that the `Key` is present in the dictionary and an exception is generated if `Key` is not in the dictionary.

`fetch_keys(Orddict) -> Keys`

Types:

Orddict = ordered_dictionary()

Keys = [term()]

This function returns a list of all keys in the dictionary.

`filter(Pred, Orddict1) -> Orddict2`

Types:

Pred = fun(Key, Value) -> bool()

Key = Value = term()

Orddict1 = Orddict2 = ordered_dictionary()

`Orddict2` is a dictionary of all keys and values in `Orddict1` for which `Pred(Key, Value)` is `true`.

`find(Key, Orddict) -> {ok, Value} | error`

Types:

Key = Value = term()

Orddict = ordered_dictionary()

This function searches for a key in a dictionary. Returns `{ok, Value}` where `Value` is the value associated with `Key`, or `error` if the key is not present in the dictionary.

`fold(Fun, Acc0, Orddict) -> Acc1`

Types:

Fun = fun(Key, Value, AccIn) -> AccOut

Key = Value = term()

Acc0 = Acc1 = AccIn = AccOut = term()

Orddict = ordered_dictionary()

Calls `Fun` on successive keys and values of `Orddict` together with an extra argument `Acc` (short for accumulator). `Fun` must return a new accumulator which is passed to the next call. `Acc0` is returned if the list is empty. The evaluation order is undefined.

`from_list(List) -> Orddict`

Types:

List = [{Key, Value}]

Orddict = ordered_dictionary()

This function converts the key/value list `List` to a dictionary.

`is_key(Key, Orddict) -> bool()`

Types:

Key = term()

Orddict = ordered_dictionary()

This function tests if `Key` is contained in the dictionary `Orddict`.

`map(Fun, Orddict1) -> Orddict2`

Types:

Fun = fun(Key, Value1) -> Value2

Key = Value1 = Value2 = term()

Orddict1 = Orddict2 = ordered_dictionary()

`map` calls `Func` on successive keys and values of `Orddict` to return a new value for each key. The evaluation order is undefined.

`merge(Fun, Orddict1, Orddict2) -> Orddict3`

Types:

Fun = fun(Key, Value1, Value2) -> Value

Key = Value1 = Value2 = Value3 = term()

Orddict1 = Orddict2 = Orddict3 = ordered_dictionary()

`merge` merges two dictionaries, `Orddict1` and `Orddict2`, to create a new dictionary. All the `Key - Value` pairs from both dictionaries are included in the new dictionary. If a key occurs in both dictionaries then `Fun` is called with the key and both values to return a new value. `merge` could be defined as:

```
merge(Fun, D1, D2) ->
    fold(fun (K, V1, D) ->
                update(K, fun (V2) -> Fun(K, V1, V2) end, V1, D)
         end, D2, D1).
```

but is faster.

`new() -> ordered_dictionary()`

This function creates a new dictionary.

`size(Orddict) -> int()`

Types:

Orddict = ordered_dictionary()

Returns the number of elements in an `Orddict`.

`store(Key, Value, Orddict1) -> Orddict2`

Types:

Key = Value = term()

Orddict1 = Orddict2 = ordered_dictionary()

This function stores a `Key - Value` pair in a dictionary. If the `Key` already exists in `Orddict1`, the associated value is replaced by `Value`.

`to_list(Orddict) -> List`

Types:

Orddict = ordered_dictionary()

List = [{Key, Value}]

This function converts the dictionary to a list representation.

update(Key, Fun, Orddict1) -> Orddict2

Types:

Key = term()

Fun = fun(Value1) -> Value2

Value1 = Value2 = term()

Orddict1 = Orddict2 = ordered_dictionary()

Update the a value in a dictionary by calling `Fun` on the value to get a new value. An exception is generated if `Key` is not present in the dictionary.

update(Key, Fun, Initial, Orddict1) -> Orddict2

Types:

Key = Initial = term()

Fun = fun(Value1) -> Value2

Value1 = Value2 = term()

Orddict1 = Orddict2 = ordered_dictionary()

Update the a value in a dictionary by calling `Fun` on the value to get a new value. If `Key` is not present in the dictionary then `Initial` will be stored as the first value. For example `append/3` could be defined as:

```
append(Key, Val, D) ->
    update(Key, fun (Old) -> Old ++ [Val] end, [Val], D).
```

update_counter(Key, Increment, Orddict1) -> Orddict2

Types:

Key = term()

Increment = number()

Orddict1 = Orddict2 = ordered_dictionary()

Add `Increment` to the value associated with `Key` and store this value. If `Key` is not present in the dictionary then `Increment` will be stored as the first value.

This could be defined as:

```
update_counter(Key, Incr, D) ->
    update(Key, fun (Old) -> Old + Incr end, Incr, D).
```

but is faster.

Notes

The functions `append` and `append_list` are included so we can store keyed values in a list *accumulator*. For example:

Orddicts are a generally good compromise between complexity and efficiency up to about 75 elements (see my benchmark). After that amount, you should switch to different key-value stores.

There are basically two key-value structures/modules to deal with larger amounts of data: dicts and gb_trees. Dictionaries have the same interface as orddicts: dict:store/3, dict:find/2, dict:fetch/2, dict:erase/2 and every other function, such as dict:map/2 and dict:fold/2 (pretty useful to work on the whole data structure!) Dicts are thus very good choices to scale orddicts up whenever it is needed.

```
> D0 = orddict:new(),
  D1 = orddict:store(files, [], D0),
  D2 = orddict:append(files, f1, D1),
  D3 = orddict:append(files, f2, D2),
  D4 = orddict:append(files, f3, D3),
  orddict:fetch(files, D4).
[f1,f2,f3]
```

This saves the trouble of first fetching a keyed value, appending a new value to the list of stored values, and storing the result.

The function `fetch` should be used if the key is known to be in the dictionary, otherwise `find`.

See Also

dict(3), *gb_trees(3)*

ordsets

Erlang module

ordsets 同 orddict 结构一样，只是存储的是 key 而不是键值对；
例如：向 ordsets 顺序插入 z、a、c 后，ordsets 结构为：[a, c, z]

Sets are collections of elements with no duplicate elements. An `ordset` is a representation of a set, where an ordered list is used to store the elements of the set. An ordered list is more efficient than an unordered list.

This module provides exactly the same interface as the module `sets` but with a defined representation. One difference is that while `sets` considers two elements as different if they do not match (`=:=`), this module considers two elements as different if and only if they do not compare equal (==).

DATA TYPES

```
ordered_set()
  as returned by new/0
```

Exports

`new() -> Ordset`

Types: `new() -> [].`

Ordset = ordered_set()

Returns a new empty ordered set.

```
is_set([E|Es]) -> is_set(Es, E);
is_set([]) -> true;
is_set(_) -> false.

is_set([E2|Es], E1) when E1 < E2 ->
    is_set(Es, E2);
is_set([_|_], _) -> false;
is_set([], _) -> true.
```

`is_set(Ordset) -> bool()`

Types:

Ordset = term()

Returns `true` if `Ordset` is an ordered set of elements, otherwise `false`.

`size(Ordset) -> int()`

Types:

Ordset = term()

Returns the number of elements in `Ordset`.

`to_list(Ordset) -> List`

Types:

Ordset = ordered_set()

List = [term()]

Returns the elements of `Ordset` as a list.

`from_list(List) -> Ordset`

Types:

List = [term()]

Ordset = ordered_set()

Returns an ordered set of the elements in `List`.

```
Eshell V5.9  (abort with ^G)
1> L1=[1,2,3,4,5,6].
[1,2,3,4,5,6]
2> L2=[s,d,f,g,1,2,3,2].
[s,d,f,g,1,2,3,2]
3> S1=ordsets:from_list(L1).
[1,2,3,4,5,6]
4> S2=ordsets:from_list(L2).
[1,2,3,d,f,g,s]
5> ordsets:subtract(S1,S2).
[4,5,6]
6> ordsets:union(S1,S2).
[1,2,3,4,5,6,d,f,g,s]
7> ordsets:is_subset(S1,S2).
false
8> ordsets:is_subset(ordsets:from_list([s,d,f]),S2).
true
9>
```

`is_element(Element, Ordset) -> bool()`

Types:

Element = term()

Ordset = ordered_set()

Returns `true` if `Element` is an element of `Ordset`, otherwise `false`.

`add_element(Element, Ordset1) -> Ordset2`

Types:

Element = term()

Ordset1 = Ordset2 = ordered_set()

Returns a new ordered set formed from `Ordset1` with `Element` inserted.

`del_element(Element, Ordset1) -> Ordset2`

Types:

Element = term()

Ordset1 = Ordset2 = ordered_set()

Returns `Ordset1`, but with `Element` removed.

`union(Ordset1, Ordset2) -> Ordset3`

Types:

Ordset1 = Ordset2 = Ordset3 = ordered_set()

Returns the merged (union) set of `Ordset1` and `Ordset2`.

`union(OrdsetList) -> Ordset`

Types:

OrdsetList = [ordered_set()]

Ordset = ordered_set()

Returns the merged (union) set of the list of sets.

`intersection(Ordset1, Ordset2) -> Ordset3`

Types:

Ordset1 = Ordset2 = Ordset3 = ordered_set()

Returns the intersection of `Ordset1` and `Ordset2`.

`intersection(OrdsetList) -> Ordset`

Types:

OrdsetList = [ordered_set()]

Ordset = ordered_set()

Returns the intersection of the non-empty list of sets.

`is_disjoint(Ordset1, Ordset2) -> bool()`

Types:

Ordset1 = Ordset2 = ordered_set()

Returns `true` if `Ordset1` and `Ordset2` are disjoint (have no elements in common), and `false` otherwise.

`subtract(Ordset1, Ordset2) -> Ordset3`

Types:

Ordset1 = Ordset2 = Ordset3 = ordered_set()

Returns only the elements of `Ordset1` which are not also elements of `Ordset2`.

`is_subset(Ordset1, Ordset2) -> bool()`

Types:

Ordset1 = Ordset2 = ordered_set()

Returns `true` when every element of `Ordset1` is also a member of `Ordset2`, otherwise `false`.

`fold(Function, Acc0, Ordset) -> Acc1`

Types:

Function = fun (E, AccIn) -> AccOut

Acc0 = Acc1 = AccIn = AccOut = term()

Ordset = ordered_set()

Fold `Function` over every element in `Ordset` returning the final value of the accumulator.

`filter(Pred, Ordset1) -> Set2`

Types:

Pred = fun (E) -> bool()

Set1 = Set2 = ordered_set()

Filter elements in `Set1` with boolean function `Fun`.

See Also

gb_sets(3), *sets(3)*

pg

Erlang module

This (experimental) module implements process groups. A process group is a group of processes that can be accessed by a common name. For example, a group named `foobar` can include a set of processes as members of this group and they can be located on different nodes.

When messages are sent to the named group, all members of the group receive the message. The messages are serialized. If the process `P1` sends the message `M1` to the group, and process `P2` simultaneously sends message `M2`, then all members of the group receive the two messages in the same order. If members of a group terminate, they are automatically removed from the group.

This module is not complete. The module is inspired by the ISIS system and the causal order protocol of the ISIS system should also be implemented. At the moment, all messages are serialized by sending them through a group master process.

Exports

`create(PgName) -> ok | {error, Reason}`

Types:

PgName = term()

Reason = already_created | term()

Creates an empty group named `PgName` on the current node.

`create(PgName, Node) -> ok | {error, Reason}`

Types:

PgName = term()

Node = node()

Reason = already_created | term()

Creates an empty group named `PgName` on the node `Node`.

`join(PgName, Pid) -> Members`

Types:

PgName = term()

Pid = pid()

Members = [pid()]

Joins the pid `Pid` to the process group `PgName`. Returns a list of all old members of the group.

`send(PgName, Msg) -> void()`

Types:

PgName = Msg = term()

Sends the tuple `{pg_message, From, PgName, Msg}` to all members of the process group `PgName`.

Failure: `{badarg, {PgName, Msg}}` if `PgName` is not a process group (a globally registered name).

esend(PgName, Msg) -> void()

Types:

PgName = Msg = term()

Sends the tuple `{pg_message, From, PgName, Msg}` to all members of the process group `PgName`, except ourselves.

Failure: `{badarg, {PgName, Msg}}` if `PgName` is not a process group (a globally registered name).

members(PgName) -> Members

Types:

PgName = term()

Members = [pid()]

Returns a list of all members of the process group `PgName`.

pool

Erlang module

`pool` can be used to run a set of Erlang nodes as a pool of computational processors. It is organized as a master and a set of slave nodes and includes the following features:

- The slave nodes send regular reports to the master about their current load.
- Queries can be sent to the master to determine which node will have the least load.

The BIF `statistics(run_queue)` is used for estimating future loads. It returns the length of the queue of ready to run processes in the Erlang runtime system.

The slave nodes are started with the `slave` module. This effects, tty IO, file IO, and code loading.

If the master node fails, the entire pool will exit.

Exports

`start(Name) ->`

`start(Name, Args) -> Nodes`

Types:

Name = atom()

Args = string()

Nodes = [node()]

Starts a new pool. The file `.hosts.erlang` is read to find host names where the pool nodes can be started. See section *Files* below. The start-up procedure fails if the file is not found.

The slave nodes are started with `slave:start/2,3`, passing along `Name` and, if provided, `Args`. `Name` is used as the first part of the node names, `Args` is used to specify command line arguments. See *slave(3)*.

Access rights must be set so that all nodes in the pool have the authority to access each other.

The function is synchronous and all the nodes, as well as all the system servers, are running when it returns a value.

`attach(Node) -> already_attached | attached`

Types:

Node = node()

This function ensures that a pool master is running and includes `Node` in the pool master's pool of nodes.

`stop() -> stopped`

Stops the pool and kills all the slave nodes.

`get_nodes() -> Nodes`

Types:

Nodes = [node()]

Returns a list of the current member nodes of the pool.

pspawn(Mod, Fun, Args) -> pid()

Types:

Mod = Fun = atom()

Args = [term()]

Spawns a process on the pool node which is expected to have the lowest future load.

pspawn_link(Mod, Fun, Args) -> pid()

Types:

Mod = Fun = atom()

Args = [term()]

Spawn links a process on the pool node which is expected to have the lowest future load.

get_node() -> node()

Returns the node with the expected lowest future load.

Files

`.hosts.erlang` is used to pick hosts where nodes can be started. See *net_adm(3)* for information about format and location of this file.

`$HOME/.erlang.slave.out.HOST` is used for all additional IO that may come from the slave nodes on standard IO. If the start-up procedure does not work, this file may indicate the reason.

proc_lib

Erlang module

普通的Erlang进程正常终止的原因只有normal,使用proc_lib启动的进程还有一个原因就是shutdown.

This module is used to start processes adhering to the *OTP Design Principles*. Specifically, the functions in this module are used by the OTP standard behaviors (`gen_server`, `gen_fsm`, ...) when starting new processes. The functions can also be used to start *special processes*, user defined processes which comply to the OTP design principles. See *Sys and Proc_Lib* in OTP Design Principles for an example.

Some useful information is initialized when a process starts. The registered names, or the process identifiers, of the parent process, and the parent ancestors, are stored together with information about the function initially called in the process. 用此方法初始化的时候会包含更多的信息,注册名 Pid 父进程 这些都会在初始化的时候存储在进程中.

While in "plain Erlang" a process is said to terminate normally only for the exit reason `normal`, a process started using `proc_lib` is also said to terminate normally if it exits with reason `shutdown` or `{shutdown,Term}`. `shutdown` is the reason used when an application (supervision tree) is stopped.

When a process started using `proc_lib` terminates abnormally -- that is, with another exit reason than `normal`, `shutdown`, or `{shutdown,Term}` -- a *crash report* is generated, which is written to terminal by the default SASL event handler. That is, the crash report is normally only visible if the SASL application is started. See *sasl(6)* and *SASL User's Guide*. 一个proc_lib创建的进程非正常终止(原因不是normal shutdown) 就会产生一个crash_report,默认会写入SASL的事件处理器. 换句话说,只有在启动了SASL的应用程序才能看到SASL的日志.

The crash report contains the previously stored information such as ancestors and initial function, the termination reason, and information regarding other processes which terminate as a result of this process terminating.

Exports

`spawn(Fun) -> pid()`

`spawn(Node, Fun) -> pid()`

`spawn(Module, Function, Args) -> pid()`

`spawn(Node, Module, Function, Args) -> pid()`

Types:

Node = node()

Fun = fun() -> void()

Module = Function = atom()

Args = [term()]

Spawns a new process and initializes it as described above. The process is spawned using the *spawn* BIFs.

`spawn_link(Fun) -> pid()`

`spawn_link(Node, Fun) -> pid()`

`spawn_link(Module, Function, Args) -> pid()`

`spawn_link(Node, Module, Function, Args) -> pid()`

Types:

Node = node()
Fun = fun() -> void()
Module = Function = atom()
Args = [term()]

Spawns a new process and initializes it as described above. The process is spawned using the *spawn_link* BIFs.

`spawn_opt(Fun, SpawnOpts) -> pid()` 这个函数的意义是:spawn_with_opt

`spawn_opt(Node, Fun, SpawnOpts) -> pid()`

`spawn_opt(Module, Function, Args, SpawnOpts) -> pid()`

`spawn_opt(Node, Module, Func, Args, SpawnOpts) -> pid()`

Types:

Node = node()
Fun = fun() -> void()
Module = Function = atom()
Args = [term()]
SpawnOpts -- see erlang:spawn_opt/2,3,4,5

Spawns a new process and initializes it as described above. The process is spawned using the *spawn_opt* BIFs.

> Note:
> Using the spawn option `monitor` is currently not allowed, but will cause the function to fail with reason `badarg`.

`start(Module, Function, Args) -> Ret`

`start(Module, Function, Args, Time) -> Ret`

`start(Module, Function, Args, Time, SpawnOpts) -> Ret`

`start_link(Module, Function, Args) -> Ret`

`start_link(Module, Function, Args, Time) -> Ret`

`start_link(Module, Function, Args, Time, SpawnOpts) -> Ret`

Types:

Module = Function = atom()
Args = [term()]
Time = int() >= 0 | infinity
SpawnOpts -- see erlang:spawn_opt/2,3,4,5
Ret = term() | {error, Reason}

```
sync_wait(Pid, Timeout) ->
    receive
 {ack, Pid, Return} ->
      Return;
 {'EXIT', Pid, Reason} ->
     {error, Reason}
    after Timeout ->
     unlink(Pid),
     exit(Pid, kill),
     flush(Pid),
     {error, timeout}
    end.
```

同步的方式启动进程,创建一个进程并等待其启动start.当进程启动之后(完成start)必须要调用一下init_ack

如果调用start_link方式创建进程而且创建的进程在调用init_ack之前就死掉了,如果调用进程做了退出捕获(trap_exit)
就会返回{error,Reason}

如果指定了Time参数,这个方法就会等待Time毫秒等待新进程调用init_ack,超时了还没有调用就会返回{error,timeout}并将新进程干掉.

Starts a new process synchronously. Spawns the process and waits for it to start. When the process has started, it *must* call *init_ack(Parent,Ret)* or *init_ack(Ret)*, where `Parent` is the process that evaluates this function. At this time, `Ret` is returned.

If the `start_link/3,4,5` function is used and the process crashes before it has called `init_ack/1,2`, `{error, Reason}` is returned if the calling process traps exits.

If `Time` is specified as an integer, this function waits for `Time` milliseconds for the new process to call `init_ack`, or `{error, timeout}` is returned, and the process is killed.

The `SpawnOpts` argument, if given, will be passed as the last argument to the `spawn_opt/2,3,4,5` BIF.

> ### Note:
>
> Using the spawn option `monitor` is currently not allowed, but will cause the function to fail with reason `badarg`.

`init_ack(Parent, Ret) -> void()`

`init_ack(Ret) -> void()`

Types: 这个方法是初始化的响应方法,告诉父进程当前进程已经完成了初始化,或者没有初始化成功.

Parent = pid()

Ret = term()

This function must used by a process that has been started by a *start[_link]/3,4,5* function. It tells `Parent` that the process has initialized itself, has started, or has failed to initialize itself.

The `init_ack/1` function uses the parent value previously stored by the start function used.

If this function is not called, the start function will return an error tuple (if a link and/or a timeout is used) or hang otherwise. 如果这个方法没有调用,那么start方法就会报错(如果使用了link timeout),或者挂起.

The following example illustrates how this function and `proc_lib:start_link/3` are used.

```
-module(my_proc).
-export([start_link/0]).
-export([init/1]).

start_link() ->
    proc_lib:start_link(my_proc, init, [self()]).

init(Parent) ->
    case do_initialization() of
        ok ->
            proc_lib:init_ack(Parent, {ok, self()});
        {error, Reason} ->
            exit(Reason)
    end,
    loop().

...
```

`format(CrashReport) -> string()`

Types:

CrashReport = term()

This function can be used by a user defined event handler to format a crash report. The crash report is sent using `error_logger:error_report(crash_report, CrashReport)`. That is, the event to be handled is of the format `{error_report, GL, {Pid, crash_report, CrashReport}}` where `GL` is the group leader pid of the process `Pid` which sent the crash report.

`initial_call(Process) -> {Module,Function,Args} | false`

Types:

查看一个进程初始化的方法是什么

Process = pid() | {X,Y,Z} | ProcInfo

X = Y = Z = int()

ProcInfo = term()

Module = Function = atom()

Args = [atom()]

Extracts the initial call of a process that was started using one of the spawn or start functions described above. `Process` can either be a pid, an integer tuple (from which a pid can be created), or the process information of a process `Pid` fetched through an `erlang:process_info(Pid)` function call.

> Note:
>
> The list `Args` no longer contains the actual arguments, but the same number of atoms as the number of arguments; the first atom is always `'Argument__1'`, the second `'Argument__2'`, and so on. The reason is that the argument list could waste a significant amount of memory, and if the argument list contained funs, it could be impossible to upgrade the code for the module.
>
> If the process was spawned using a fun, `initial_call/1` no longer returns the actual fun, but the module, function for the local function implementing the fun, and the arity, for instance `{some_module,-work/3-fun-0-,0}` (meaning that the fun was created in the function `some_module:work/3`). The reason is that keeping the fun would prevent code upgrade for the module, and that a significant amount of memory could be wasted. 如果使用fun并不能原样还原,不保存的原因:保存会影响代码更新,还会浪费很多内存

`translate_initial_call(Process) -> {Module,Function,Arity} | Fun`

Types:

Process = pid() | {X,Y,Z} | ProcInfo

X = Y = Z = int()

ProcInfo = term()

Module = Function = atom()

Arity = int()

Fun = fun() -> void()

This function is used by the `c:i/0` and `c:regs/0` functions in order to present process information.

Extracts the initial call of a process that was started using one of the spawn or start functions described above, and translates it to more useful information. `Process` can either be a pid, an integer tuple (from which a pid can be created), or the process information of a process `Pid` fetched through an `erlang:process_info(Pid)` function call.

If the initial call is to one of the system defined behaviors such as `gen_server` or `gen_event`, it is translated to more useful information. If a `gen_server` is spawned, the returned `Module` is the name of the callback module and `Function` is `init` (the function that initiates the new server).

A `supervisor` and a `supervisor_bridge` are also `gen_server` processes. In order to return information that this process is a supervisor and the name of the call-back module, `Module` is `supervisor` and `Function` is the name of the supervisor callback module. `Arity` is `1` since the `init/1` function is called initially in the callback module.

By default, `{proc_lib,init_p,5}` is returned if no information about the initial call can be found. It is assumed that the caller knows that the process has been spawned with the `proc_lib` module.

`hibernate(Module, Function, Args)`

Types:

Module = Function = atom()

Args = [term()]

This function does the same as (and does call) the BIF *hibernate/3*, but ensures that exception handling and logging continues to work as expected when the process wakes up. Always use this function instead of the BIF for processes started using `proc_lib` functions.

SEE ALSO

相比BIF的hibernate/3 多做的事情就是异常处理和日志
建议使用proc_lib创建的进程都使用这个方法来进行休眠.

error_logger(3)

proplists

本模块适用数据量较少的场景.构成形式是 [{Key,Value}].
配置文件,继承选项的情况常用 操作的粒度就是一个数据项,没有更新和追加数据项的方法,需要lists:replace/4之类的解决

Erlang module

proplists是包含kv的tuple,第一个元素做key用来查询和删除,如果一个key有多个值就会使用第一次出现的值,其它被忽略

Property lists are ordinary lists containing entries in the form of either tuples, whose first elements are keys used for lookup and insertion, or atoms, which work as shorthand for tuples `{Atom, true}`. (Other terms are allowed in the lists, but are ignored by this module.) If there is more than one entry in a list for a certain key, the first occurrence normally overrides any later (irrespective of the arity of the tuples).

Property lists are useful for representing inherited properties, such as options passed to a function where a user may specify options overriding the default settings, object properties, annotations, etc.

Two keys are considered equal if they match (`=:=`). In other words, numbers are compared literally rather than by value, so that, for instance, `1` and `1.0` are different keys.

proplists用来表示继承属性; key值比较使用的是=:=精确等于

Exports

%%将所有key值相同的数据项,value整合在list中.

`append_values(Key, List) -> List`

1> proplists:append_values(a, [{a, [1,2]}, {b, 0}, {a, 3}, {c, -1}, {a, [4]}]).
[1,2,3,4]
2> proplists:append_values(a, [{a, b}, {b, 0}, {a, 3}, {c, -1}, {a, [4]}]).
[b,3,4]
3> proplists:append_values(a, [{a, b}, {b, 0}, {a, 3}, {c, -1}, {a, [[4]]}]).
[b,3,[4]]

Types:

Key = term()

List = [term()]

Similar to `get_all_values/2`, but each value is wrapped in a list unless it is already itself a list, and the resulting list of lists is concatenated. This is often useful for "incremental" options; e.g., `append_values(a, [{a, [1,2]}, {b, 0}, {a, 3}, {c, -1}, {a, [4]}])` will return the list `[1,2,3,4]`.

%%将数据进行压缩

`compact(List) -> List`

Types:

List = [term()]

10> proplists:compact([{a, true}, {b, true}, {a, 3}, {c, true}, {a, [4]}]).
[a,b,{a,3},c,{a,[4]}]

Minimizes the representation of all entries in the list. This is equivalent to `[property(P) || P <- List]`.

See also: `property/1`, `unfold/1`.

`delete(Key, List) -> List`

Types:

Key = term()

List = [term()]

18> proplists:delete(a, [{a, true}, {b, true}, {a, 3}, {c, true}, {a, [4]}]).
[{b,true},{c,true}]

Deletes all entries associated with `Key` from `List`.

`expand(Expansions, List) -> List`

Types:

Key = term()

Expansions = [{Property,[term()]}]

Property = atom() | tuple()

```
expand(Es, Ps) when is_list(Ps) ->
    Es1 = [{property(P), V} || {P, V} <- Es],
    flatten(expand_0(key_uniq(Es1), Ps)).
```

Expands particular properties to corresponding sets of properties (or other terms). For each pair `{Property, Expansion}` in `Expansions`, if `E` is the first entry in `List` with the same key as `Property`, and `E` and `Property` have equivalent normal forms, then `E` is replaced with the terms in `Expansion`, and any following entries with the same key are deleted from `List`.

%%注意:下面的demo排版换行了
1> proplists: expand([{foo, [bar, baz]}],[fie, foo, fum]).
[fie,bar,baz,fum]

For example, the following expressions all return `[fie, bar, baz, fum]`:

```
expand([{foo, [bar, baz]}],
[fie, foo, fum])
expand([{{foo, true}, [bar, baz]}],
[fie, foo, fum])
expand([{{foo, false}, [bar, baz]}],
[fie, {foo, false}, fum])
```

However, no expansion is done in the following call:

```
expand([{{foo, true}, [bar, baz]}],
[{foo, false}, fie, foo, fum])
```

because `{foo, false}` shadows `foo`.

Note that if the original property term is to be preserved in the result when expanded, it must be included in the expansion list. The inserted terms are not expanded recursively. If `Expansions` contains more than one property with the same key, only the first occurrence is used.

See also: `normalize/2`.

`get_all_values(Key, List) -> [term()]`

Types:

Key = term()

List = [term()]

Similar to `get_value/2`, but returns the list of values for *all* entries `{Key, Value}` in `List`. If no such entry exists, the result is the empty list.

See also: `get_value/2`.

`get_bool(Key, List) -> bool()`

Types:

Key = term()

List = [term()]

Returns the value of a boolean key/value option. If `lookup(Key, List)` would yield `{Key, true}`, this function returns `true`; otherwise `false` is returned.

See also: `get_value/2`, `lookup/2`.

```
get_keys(Ps) ->
    sets:to_list(get_keys(Ps, sets:new())).

get_keys([P | Ps], Keys) ->
    if is_atom(P) ->
     get_keys(Ps, sets:add_element(P, Keys));
       tuple_size(P) >= 1 ->
     get_keys(Ps, sets:add_element(element(1, P), Keys));
       true ->
     get_keys(Ps, Keys)
    end;
get_keys([], Keys) ->
    Keys.
```

`get_keys(List) -> [term()]`

Types:

List = [term()]

Returns an unordered list of the keys used in `List`, not containing duplicates.

`get_value(Key, List) -> term()`

Types:

Key = term()

```
3> proplists:get_value([a,b], ["packet",[a,b],"login",22,2,s,f] , "none").
"none"
4> proplists:get_value("login", ["packet",[a,b],"login",22,2,s,f] , "none").
"none"
5> proplists:get_value(login, ["packet",[a,b],"login",22,2,s,f] , "none").
"none"
```


List = [term()]

Equivalent to `get_value(Key, List, undefined)`.

`get_value(Key, List, Default) -> term()`

Types:

Key = term()

Default = term()

List = [term()]

```
get_value(Key, [P | Ps], Default) ->
    if is_atom(P), P =:= Key ->
     true;
       tuple_size(P) >= 1, element(1, P) =:= Key ->
     case P of
      {_, Value} ->
         Value;
      _ ->
         %% Don</code>t continue the search!
         Default
        end;
          true ->
        get_value(Key, Ps, Default)
    end;
get_value(_Key, [], Default) ->
  Default.
```

Returns the value of a simple key/value property in `List`. If `lookup(Key, List)` would yield `{Key, Value}`, this function returns the corresponding `Value`, otherwise `Default` is returned.

See also: `get_all_values/2`, `get_bool/2`, `get_value/2`, `lookup/2`.

`is_defined(Key, List) -> bool()`

Types:

Key = term()

List = [term()]

Returns `true` if `List` contains at least one entry associated with `Key`, otherwise `false` is returned.

`lookup(Key, List) -> none | tuple()`

Types:

Key = term()

List = [term()]

```
6> proplists:lookup(a, [{a, b}, {b, 0}, {a, 3}, {c, -1}, {a, [[4]]}]).
{a,b}
7> proplists:lookup(a, [{a,1},{a, b}, {b, 0}, {a, 3}, {c, -1}, {a, [[4]]}]).
{a,1}
8>
```

Returns the first entry associated with `Key` in `List`, if one exists, otherwise returns `none`. For an atom `A` in the list, the tuple `{A, true}` is the entry associated with `A`.

See also: `get_bool/2`, `get_value/2`, `lookup_all/2`.

`lookup_all(Key, List) -> [tuple()]`

Types:

Key = term()

List = [term()]

```
8> proplists:lookup_all(a, [{a,1},{a, b}, {b, 0}, {a, 3}, {c, -1}, {a, [[4]]}]).
[{a,1},{a,b},{a,3},{a,[[4]]}]
```

Returns the list of all entries associated with `Key` in `List`. If no such entry exists, the result is the empty list.

See also: `lookup/2`.

`normalize(List, Stages) -> List`

Types:

List = [term()]

Stages = [Operation]

Operation = {aliases, Aliases} | {negations, Negations} | {expand, Expansions}

Aliases = [{Key, Key}]

Negations = [{Key, Key}]

Key = term()

Expansions = [{Property, [term()]}]

Property = atom() | tuple()

Passes `List` through a sequence of substitution/expansion stages. For an `aliases` operation, the function `substitute_aliases/2` is applied using the given list of aliases; for a `negations` operation, `substitute_negations/2` is applied using the given negation list; for an `expand` operation, the function `expand/2` is applied using the given list of expansions. The final result is automatically compacted (cf. `compact/1`).

Typically you want to substitute negations first, then aliases, then perform one or more expansions (sometimes you want to pre-expand particular entries before doing the main expansion). You might want to substitute negations and/or aliases repeatedly, to allow such forms in the right-hand side of aliases and expansion lists.

See also: `compact/1`, `expand/2`, `substitute_aliases/2`, `substitute_negations/2`.

`property(Property) -> Property`

Types:

Property = atom() | tuple()

```
16>  proplists:property({a}).
{a}
17>  proplists:property({a,true}).
a
```

Creates a normal form (minimal) representation of a property. If `Property` is `{Key, true}` where `Key` is an atom, this returns `Key`, otherwise the whole term `Property` is returned.

See also: `property/2`.

`property(Key, Value) -> Property`

Types:

Key = term()

Value = term()

Property = atom() | tuple()

```
property(Key, true) when is_atom(Key) ->
    Key;
property(Key, Value) ->
    {Key, Value}.
```

Creates a normal form (minimal) representation of a simple key/value property. Returns `Key` if `Value` is `true` and `Key` is an atom, otherwise a tuple `{Key, Value}` is returned.

See also: `property/1`.

`split(List, Keys) -> {Lists, Rest}`

Types:

List = [term()]

Keys = [term()]

Lists = [[term()]]

Rest = [term()]

Partitions `List` into a list of sublists and a remainder. `Lists` contains one sublist for each key in `Keys`, in the corresponding order. The relative order of the elements in each sublist is preserved from the original `List`. `Rest` contains the elements in `List` that are not associated with any of the given keys, also with their original relative order preserved.

Example: split([{c, 2}, {e, 1}, a, {c, 3, 4}, d, {b, 5}, b], [a, b, c])

returns

{[[a], [{b, 5}, b],[{c, 2}, {c, 3, 4}]], [{e, 1}, d]}

`substitute_aliases(Aliases, List) -> List`

Types:

Aliases = [{Key, Key}]

Key = term()

List = [term()]

Substitutes keys of properties. For each entry in `List`, if it is associated with some key `K1` such that `{K1, K2}` occurs in `Aliases`, the key of the entry is changed to `Key2`. If the same `K1` occurs more than once in `Aliases`, only the first occurrence is used.

Example: `substitute_aliases([{color, colour}], L)` will replace all tuples `{color, ...}` in `L` with `{colour, ...}`, and all atoms `color` with `colour`.

See also: `normalize/2`, `substitute_negations/2`.

`substitute_negations(Negations, List) -> List`

Types:

Negations = [{Key, Key}]

Key = term()

List = [term()]

Substitutes keys of boolean-valued properties and simultaneously negates their values. For each entry in `List`, if it is associated with some key `K1` such that `{K1, K2}` occurs in `Negations`, then if the entry was `{K1, true}` it will be replaced with `{K2, false}`, otherwise it will be replaced with `{K2, true}`, thus changing the name of the option and simultaneously negating the value given by `get_bool(List)`. If the same `K1` occurs more than once in `Negations`, only the first occurrence is used.

Example: `substitute_negations([{no_foo, foo}], L)` will replace any atom `no_foo` or tuple `{no_foo, true}` in `L` with `{foo, false}`, and any other tuple `{no_foo, ...}` with `{foo, true}`.

See also: `get_bool/2`, `normalize/2`, `substitute_aliases/2`.

`unfold(List) -> List`

Types:

List = [term()]

Unfolds all occurrences of atoms in `List` to tuples `{Atom, true}`.

```
12> proplists:unfold([foo,bar,test,haha]).
[{foo,true},{bar,true},{test,true},{haha,true}]
13> proplists:unfold([foo,bar,{test,false},haha]).
[{foo,true},{bar,true},{test,false},{haha,true}]
14>
```

qlc

Erlang module

The `qlc` module provides a query interface to Mnesia, ETS, Dets and other data structures that implement an iterator style traversal of objects.

Overview

The `qlc` module implements a query interface to *QLC tables*. Typical QLC tables are ETS, Dets, and Mnesia tables. There is also support for user defined tables, see the *Implementing a QLC table* section. A *query* is stated using *Query List Comprehensions* (QLCs). The answers to a query are determined by data in QLC tables that fulfill the constraints expressed by the QLCs of the query. QLCs are similar to ordinary list comprehensions as described in the Erlang Reference Manual and Programming Examples except that variables introduced in patterns cannot be used in list expressions. In fact, in the absence of optimizations and options such as `cache` and `unique` (see below), every QLC free of QLC tables evaluates to the same list of answers as the identical ordinary list comprehension.

While ordinary list comprehensions evaluate to lists, calling *qlc:q/1,2* returns a *Query Handle*. To obtain all the answers to a query, *qlc:eval/1,2* should be called with the query handle as first argument. Query handles are essentially functional objects ("funs") created in the module calling `q/1,2`. As the funs refer to the module's code, one should be careful not to keep query handles too long if the module's code is to be replaced. Code replacement is described in the *Erlang Reference Manual*. The list of answers can also be traversed in chunks by use of a *Query Cursor*. Query cursors are created by calling *qlc:cursor/1,2* with a query handle as first argument. Query cursors are essentially Erlang processes. One answer at a time is sent from the query cursor process to the process that created the cursor.

Syntax

Syntactically QLCs have the same parts as ordinary list comprehensions:

```
[Expression || Qualifier1, Qualifier2, ...]
```

`Expression` (the *template*) is an arbitrary Erlang expression. Qualifiers are either *filters* or *generators*. Filters are Erlang expressions returning `bool()`. Generators have the form `Pattern <- ListExpression`, where `ListExpression` is an expression evaluating to a query handle or a list. Query handles are returned from `qlc:table/2`, `qlc:append/1,2`, `qlc:sort/1,2`, `qlc:keysort/2,3`, `qlc:q/1,2`, and `qlc:string_to_handle/1,2,3`.

Evaluation

The evaluation of a query handle begins by the inspection of options and the collection of information about tables. As a result qualifiers are modified during the optimization phase. Next all list expressions are evaluated. If a cursor has been created evaluation takes place in the cursor process. For those list expressions that are QLCs, the list expressions of the QLCs' generators are evaluated as well. One has to be careful if list expressions have side effects since the order in which list expressions are evaluated is unspecified. Finally the answers are found by evaluating the qualifiers from left to right, backtracking when some filter returns `false`, or collecting the template when all filters return `true`.

Filters that do not return `bool()` but fail are handled differently depending on their syntax: if the filter is a guard it returns `false`, otherwise the query evaluation fails. This behavior makes it possible for the `qlc` module to do some optimizations without affecting the meaning of a query. For example, when testing some position of a table and one or more constants for equality, only the objects with equal values are candidates for further evaluation. The other objects are guaranteed to make the filter return `false`, but never fail. The (small) set of candidate objects can often be found by looking up some key values of the table or by traversing the table using a match specification. It is necessary to place the guard filters immediately after the table's generator, otherwise the candidate objects will not be restricted

to a small set. The reason is that objects that could make the query evaluation fail must not be excluded by looking up a key or running a match specification.

Join

The `qlc` module supports fast join of two query handles. Fast join is possible if some position `P1` of one query handler and some position `P2` of another query handler are tested for equality. Two fast join methods have been implemented:

- Lookup join traverses all objects of one query handle and finds objects of the other handle (a QLC table) such that the values at `P1` and `P2` match or compare equal. The `qlc` module does not create any indices but looks up values using the key position and the indexed positions of the QLC table.
- Merge join sorts the objects of each query handle if necessary and filters out objects where the values at `P1` and `P2` do not compare equal. If there are many objects with the same value of `P2` a temporary file will be used for the equivalence classes.

The `qlc` module warns at compile time if a QLC combines query handles in such a way that more than one join is possible. In other words, there is no query planner that can choose a good order between possible join operations. It is up to the user to order the joins by introducing query handles.

The join is to be expressed as a guard filter. The filter must be placed immediately after the two joined generators, possibly after guard filters that use variables from no other generators but the two joined generators. The `qlc` module inspects the operands of `=:=/2`, `==/2`, `is_record/2`, `element/2`, and logical operators (`and/2`, `or/2`, `andalso/2`, `orelse/2`, `xor/2`) when determining which joins to consider.

Common options

The following options are accepted by `cursor/2`, `eval/2`, `fold/4`, and `info/2`:

- `{cache_all, Cache}` where `Cache` is equal to `ets` or `list` adds a `{cache, Cache}` option to every list expression of the query except tables and lists. Default is `{cache_all, no}`. The option `cache_all` is equivalent to `{cache_all, ets}`.
- `{max_list_size, MaxListSize}` where `MaxListSize` is the size in bytes of terms on the external format. If the accumulated size of collected objects exceeds `MaxListSize` the objects are written onto a temporary file. This option is used by the `{cache, list}` option as well as by the merge join method. Default is 512*1024 bytes.
- `{tmpdir_usage, TmpFileUsage}` determines the action taken when `qlc` is about to create temporary files on the directory set by the `tmpdir` option. If the value is `not_allowed` an error tuple is returned, otherwise temporary files are created as needed. Default is `allowed` which means that no further action is taken. The values `info_msg`, `warning_msg`, and `error_msg` mean that the function with the corresponding name in the module `error_logger` is called for printing some information (currently the stacktrace).
- `{tmpdir, TempDirectory}` sets the directory used by merge join for temporary files and by the `{cache, list}` option. The option also overrides the `tmpdir` option of `keysort/3` and `sort/2`. The default value is `""` which means that the directory returned by `file:get_cwd()` is used.
- `{unique_all, true}` adds a `{unique, true}` option to every list expression of the query. Default is `{unique_all, false}`. The option `unique_all` is equivalent to `{unique_all, true}`.

Common data types

- `QueryCursor = {qlc_cursor, term()}`
- `QueryHandle = {qlc_handle, term()}`
- `QueryHandleOrList = QueryHandle | list()`
- `Answers = [Answer]`
- `Answer = term()`

- `AbstractExpression =` - parse trees for Erlang expressions, see the *abstract format* documentation in the ERTS User's Guide -
- `MatchExpression =` - match specifications, see the *match specification* documentation in the ERTS User's Guide and *ms_transform(3)* -
- `SpawnOptions = default | spawn_options()`
- `SortOptions = [SortOption] | SortOption`
- `SortOption = {compressed, bool()} | {no_files, NoFiles} | {order, Order} | {size, Size} | {tmpdir, TempDirectory} | {unique, bool()}` - see *file_sorter(3)* -
- `Order = ascending | descending | OrderFun`
- `OrderFun = fun(term(), term()) -> bool()`
- `TempDirectory = "" | filename()`
- `Size = int() > 0`
- `NoFiles = int() > 1`
- `KeyPos = int() > 0 | [int() > 0]`
- `MaxListSize = int() >= 0`
- `bool() = true | false`
- `Cache = ets | list | no`
- `TmpFileUsage = allowed | not_allowed | info_msg | warning_msg | error_msg`
- `filename() =` - see *filename(3)* -
- `spawn_options() =` - see *erlang(3)* -

Getting started

As already mentioned queries are stated in the list comprehension syntax as described in the *Erlang Reference Manual*. In the following some familiarity with list comprehensions is assumed. There are examples in *Programming Examples* that can get you started. It should be stressed that list comprehensions do not add any computational power to the language; anything that can be done with list comprehensions can also be done without them. But they add a syntax for expressing simple search problems which is compact and clear once you get used to it.

Many list comprehension expressions can be evaluated by the `qlc` module. Exceptions are expressions such that variables introduced in patterns (or filters) are used in some generator later in the list comprehension. As an example consider an implementation of lists:append(L): `[X ||Y <- L, X <- Y]`. Y is introduced in the first generator and used in the second. The ordinary list comprehension is normally to be preferred when there is a choice as to which to use. One difference is that `qlc:eval/1,2` collects answers in a list which is finally reversed, while list comprehensions collect answers on the stack which is finally unwound.

What the `qlc` module primarily adds to list comprehensions is that data can be read from QLC tables in small chunks. A QLC table is created by calling `qlc:table/2`. Usually `qlc:table/2` is not called directly from the query but via an interface function of some data structure. There are a few examples of such functions in Erlang/OTP: `mnesia:table/1,2`, `ets:table/1,2`, and `dets:table/1,2`. For a given data structure there can be several functions that create QLC tables, but common for all these functions is that they return a query handle created by `qlc:table/2`. Using the QLC tables provided by OTP is probably sufficient in most cases, but for the more advanced user the section *Implementing a QLC table* describes the implementation of a function calling `qlc:table/2`.

Besides `qlc:table/2` there are other functions that return query handles. They might not be used as often as tables, but are useful from time to time. `qlc:append` traverses objects from several tables or lists after each other. If, for instance, you want to traverse all answers to a query QH and then finish off by a term `{finished}`, you can do that by calling `qlc:append(QH, [{finished}])`. `append` first returns all objects of QH, then `{finished}`. If there is one tuple `{finished}` among the answers to QH it will be returned twice from `append`.

As another example, consider concatenating the answers to two queries QH1 and QH2 while removing all duplicates. The means to accomplish this is to use the `unique` option:

```
qlc:q([X || X <- qlc:append(QH1, QH2)], {unique, true})
```

The cost is substantial: every returned answer will be stored in an ETS table. Before returning an answer it is looked up in the ETS table to check if it has already been returned. Without the `unique` options all answers to QH1 would be returned followed by all answers to QH2. The `unique` options keeps the order between the remaining answers.

If the order of the answers is not important there is the alternative to sort the answers uniquely:

```
qlc:sort(qlc:q([X || X <- qlc:append(QH1, QH2)], {unique, true})).
```

This query also removes duplicates but the answers will be sorted. If there are many answers temporary files will be used. Note that in order to get the first unique answer all answers have to be found and sorted. Both alternatives find duplicates by comparing answers, that is, if A1 and A2 are answers found in that order, then A2 is a removed if A1 == A2.

To return just a few answers cursors can be used. The following code returns no more than five answers using an ETS table for storing the unique answers:

```
C = qlc:cursor(qlc:q([X || X <- qlc:append(QH1, QH2)],{unique,true})),
R = qlc:next_answers(C, 5),
ok = qlc:delete_cursor(C),
R.
```

Query list comprehensions are convenient for stating constraints on data from two or more tables. An example that does a natural join on two query handles on position 2:

```
qlc:q([{X1,X2,X3,Y1} ||
          {X1,X2,X3} <- QH1,
          {Y1,Y2} <- QH2,
          X2 =:= Y2])
```

The `qlc` module will evaluate this differently depending on the query handles `QH1` and `QH2`. If, for example, `X2` is matched against the key of a QLC table the lookup join method will traverse the objects of `QH2` while looking up key values in the table. On the other hand, if neither `X2` nor `Y2` is matched against the key or an indexed position of a QLC table, the merge join method will make sure that `QH1` and `QH2` are both sorted on position 2 and next do the join by traversing the objects one by one.

The `join` option can be used to force the `qlc` module to use a certain join method. For the rest of this section it is assumed that the excessively slow join method called "nested loop" has been chosen:

```
qlc:q([{X1,X2,X3,Y1} ||
          {X1,X2,X3} <- QH1,
          {Y1,Y2} <- QH2,
          X2 =:= Y2],
      {join, nested_loop})
```

In this case the filter will be applied to every possible pair of answers to QH1 and QH2, one at a time. If there are M answers to QH1 and N answers to QH2 the filter will be run M*N times.

If QH2 is a call to the function for `gb_trees` as defined in the *Implementing a QLC table* section, `gb_table:table/1`, the iterator for the gb-tree will be initiated for each answer to QH1 after which the objects of the gb-tree will be returned one by one. This is probably the most efficient way of traversing the table in that case since it takes minimal computational power to get the following object. But if QH2 is not a table but a more complicated QLC, it can be more efficient use some RAM memory for collecting the answers in a cache, particularly if there are only a few answers. It must then be assumed that evaluating QH2 has no side effects so that the meaning of the query does not change if QH2 is evaluated only once. One way of caching the answers is to evaluate QH2 first of all and substitute the list of answers for QH2 in the query. Another way is to use the `cache` option. It is stated like this:

```
QH2' = qlc:q([X || X <- QH2], {cache, ets})
```

or just

```
QH2' = qlc:q([X || X <- QH2], cache)
```

The effect of the `cache` option is that when the generator QH2' is run the first time every answer is stored in an ETS table. When next answer of QH1 is tried, answers to QH2' are copied from the ETS table which is very fast. As for the `unique` option the cost is a possibly substantial amount of RAM memory. The `{cache, list}` option offers the possibility to store the answers in a list on the process heap. While this has the potential of being faster than ETS tables since there is no need to copy answers from the table it can often result in slower evaluation due to more garbage collections of the process' heap as well as increased RAM memory consumption due to larger heaps. Another drawback with cache lists is that if the size of the list exceeds a limit a temporary file will be used. Reading the answers from a file is very much slower than copying them from an ETS table. But if the available RAM memory is scarce setting the *limit* to some low value is an alternative.

There is an option `cache_all` that can be set to `ets` or `list` when evaluating a query. It adds a `cache` or `{cache, list}` option to every list expression except QLC tables and lists on all levels of the query. This can be used for testing if caching would improve efficiency at all. If the answer is yes further testing is needed to pinpoint the generators that should be cached.

Implementing a QLC table

As an example of how to use the *qlc:table/2* function the implementation of a QLC table for the *gb_trees* module is given:

```
-module(gb_table).

-export([table/1]).

table(T) ->
    TF = fun() -> qlc_next(gb_trees:next(gb_trees:iterator(T))) end,
    InfoFun = fun(num_of_objects) -> gb_trees:size(T);
                 (keypos) -> 1;
                 (is_sorted_key) -> true;
                 (is_unique_objects) -> true;
                 (_) -> undefined
              end,
    LookupFun =
        fun(1, Ks) ->
                lists:flatmap(fun(K) ->
```

```
                                    case gb_trees:lookup(K, T) of
                                        {value, V} -> [{K,V}];
                                        none -> []
                                    end
                            end, Ks)
        end,
    FormatFun =
        fun({all, NElements, ElementFun}) ->
                ValsS = io_lib:format("gb_trees:from_orddict(~w)",
                                    [gb_nodes(T, NElements, ElementFun)]),
                io_lib:format("gb_table:table(~s)", [ValsS]);
           ({lookup, 1, KeyValues, _NElements, ElementFun}) ->
                ValsS = io_lib:format("gb_trees:from_orddict(~w)",
                                    [gb_nodes(T, infinity, ElementFun)]),
                io_lib:format("lists:flatmap(fun(K) -> "
                            "case gb_trees:lookup(K, ~s) of "
                            "{value, V} -> [{K,V}];none -> [] end "
                            "end, ~w)",
                            [ValsS, [ElementFun(KV) || KV <- KeyValues]])
        end,
    qlc:table(TF, [{info_fun, InfoFun}, {format_fun, FormatFun},
                   {lookup_fun, LookupFun},{key_equality,'=='}]).

qlc_next({X, V, S}) ->
    [{X,V} | fun() -> qlc_next(gb_trees:next(S)) end];
qlc_next(none) ->
    [].

gb_nodes(T, infinity, ElementFun) ->
    gb_nodes(T, -1, ElementFun);
gb_nodes(T, NElements, ElementFun) ->
    gb_iter(gb_trees:iterator(T), NElements, ElementFun).

gb_iter(_I, 0, _EFun) ->
    '...';
gb_iter(I0, N, EFun) ->
    case gb_trees:next(I0) of
        {X, V, I} ->
            [EFun({X,V}) | gb_iter(I, N-1, EFun)];
        none ->
            []
    end.
```

`TF` is the traversal function. The `qlc` module requires that there is a way of traversing all objects of the data structure; in `gb_trees` there is an iterator function suitable for that purpose. Note that for each object returned a new fun is created. As long as the list is not terminated by `[]` it is assumed that the tail of the list is a nullary function and that calling the function returns further objects (and functions).

The lookup function is optional. It is assumed that the lookup function always finds values much faster than it would take to traverse the table. The first argument is the position of the key. Since `qlc_next` returns the objects as {Key, Value} pairs the position is 1. Note that the lookup function should return {Key, Value} pairs, just as the traversal function does.

The format function is also optional. It is called by `qlc:info` to give feedback at runtime of how the query will be evaluated. One should try to give as good feedback as possible without showing too much details. In the example at most 7 objects of the table are shown. The format function handles two cases: `all` means that all objects of the table will be traversed; `{lookup, 1, KeyValues}` means that the lookup function will be used for looking up key values.

Whether the whole table will be traversed or just some keys looked up depends on how the query is stated. If the query has the form

```
qlc:q([T || P <- LE, F])
```

and P is a tuple, the `qlc` module analyzes P and F in compile time to find positions of the tuple P that are tested for equality to constants. If such a position at runtime turns out to be the key position, the lookup function can be used, otherwise all objects of the table have to be traversed. It is the info function `InfoFun` that returns the key position. There can be indexed positions as well, also returned by the info function. An index is an extra table that makes lookup on some position fast. Mnesia maintains indices upon request, thereby introducing so called secondary keys. The `qlc` module prefers to look up objects using the key before secondary keys regardless of the number of constants to look up.

Key equality

In Erlang there are two operators for testing term equality, namely `==/2` and `=:=/2`. The difference between them is all about the integers that can be represented by floats. For instance, `2 == 2.0` evaluates to `true` while `2 =:= 2.0` evaluates to `false`. Normally this is a minor issue, but the `qlc` module cannot ignore the difference, which affects the user's choice of operators in QLCs.

If the `qlc` module can find out at compile time that some constant is free of integers, it does not matter which one of `==/2` or `=:=/2` is used:

```
1> E1 = ets:new(t, [set]), % uses =:=/2 for key equality
Q1 = qlc:q([K ||
{K} <- ets:table(E1),
K == 2.71 orelse K == a]),
io:format("~s~n", [qlc:info(Q1)]).
ets:match_spec_run(lists:flatmap(fun(V) ->
                                     ets:lookup(20493, V)
                                 end,
                                 [a,2.71]),
                   ets:match_spec_compile([{{'$1'},[],['$1']}]))
```

In the example the `==/2` operator has been handled exactly as `=:=/2` would have been handled. On the other hand, if it cannot be determined at compile time that some constant is free of integers and the table uses `=:=/2` when comparing keys for equality (see the option *key_equality*), the `qlc` module will not try to look up the constant. The reason is that there is in the general case no upper limit on the number of key values that can compare equal to such a constant; every combination of integers and floats has to be looked up:

```
2> E2 = ets:new(t, [set]),
true = ets:insert(E2, [{{2,2},a},{{2,2.0},b},{{2.0,2},c}]),
F2 = fun(I) ->
qlc:q([V || {K,V} <- ets:table(E2), K == I])
end,
Q2 = F2({2,2}),
io:format("~s~n", [qlc:info(Q2)]).
ets:table(53264,
          [{traverse,
            {select,[{{'$1','$2'},[{'==','$1',{const,{2,2}}}],['$2']}]}}])
3> lists:sort(qlc:e(Q2)).
[a,b,c]
```

Looking up just `{2,2}` would not return `b` and `c`.

If the table uses `==/2` when comparing keys for equality, the `qlc` module will look up the constant regardless of which operator is used in the QLC. However, `==/2` is to be preferred:

```
4> E3 = ets:new(t, [ordered_set]), % uses ==/2 for key equality
true = ets:insert(E3, [{{2,2.0},b}]),
F3 = fun(I) ->
qlc:q([V || {K,V} <- ets:table(E3), K == I])
end,
Q3 = F3({2,2}),
io:format("~s~n", [qlc:info(Q3)]).
ets:match_spec_run(ets:lookup(86033, {2,2}),
                   ets:match_spec_compile([{{'$1','$2'},[],['$2']}]))
5> qlc:e(Q3).
[b]
```

Lookup join is handled analogously to lookup of constants in a table: if the join operator is `==/2` and the table where constants are to be looked up uses `=:=/2` when testing keys for equality, the `qlc` module will not consider lookup join for that table.

Exports

append(QHL) -> QH

Types:

QHL = **[QueryHandleOrList]**

QH = **QueryHandle**

Returns a query handle. When evaluating the query handle `QH` all answers to the first query handle in `QHL` is returned followed by all answers to the rest of the query handles in `QHL`.

append(QH1, QH2) -> QH3

Types:

QH1 = **QH2** = **QueryHandleOrList**

QH3 = **QueryHandle**

Returns a query handle. When evaluating the query handle `QH3` all answers to `QH1` are returned followed by all answers to `QH2`.

`append(QH1, QH2)` is equivalent to `append([QH1, QH2])`.

cursor(QueryHandleOrList [, Options]) -> QueryCursor

Types:

Options = **[Option] | Option**

Option = **{cache_all, Cache} | cache_all | {max_list_size, MaxListSize} | {spawn_options, SpawnOptions} | {tmpdir_usage, TmpFileUsage} | {tmpdir, TempDirectory} | {unique_all, bool()} | unique_all**

Creates a query cursor and makes the calling process the owner of the cursor. The cursor is to be used as argument to `next_answers/1,2` and (eventually) `delete_cursor/1`. Calls `erlang:spawn_opt` to spawn and link a process which will evaluate the query handle. The value of the option `spawn_options` is used as last argument when calling `spawn_opt`. The default value is `[link]`.

```
1> QH = qlc:q([{X,Y} || X <- [a,b], Y <- [1,2]]),
QC = qlc:cursor(QH),
qlc:next_answers(QC, 1).
[{a,1}]
2> qlc:next_answers(QC, 1).
```

```
[{a,2}]
3> qlc:next_answers(QC, all_remaining).
[{b,1},{b,2}]
4> qlc:delete_cursor(QC).
ok
```

delete_cursor(QueryCursor) -> ok

Deletes a query cursor. Only the owner of the cursor can delete the cursor.

eval(QueryHandleOrList [, Options]) -> Answers | Error

e(QueryHandleOrList [, Options]) -> Answers

Types:

Options = [Option] | Option

Option = {cache_all, Cache} | cache_all | {max_list_size, MaxListSize} | {tmpdir_usage, TmpFileUsage} | {tmpdir, TempDirectory} | {unique_all, bool()} | unique_all

Error = {error, module(), Reason}

Reason = - as returned by file_sorter(3) -

Evaluates a query handle in the calling process and collects all answers in a list.

```
1> QH = qlc:q([{X,Y} || X <- [a,b], Y <- [1,2]]),
qlc:eval(QH).
[{a,1},{a,2},{b,1},{b,2}]
```

fold(Function, Acc0, QueryHandleOrList [, Options]) -> Acc1 | Error

Types:

Function = fun(Answer, AccIn) -> AccOut

Acc0 = Acc1 = AccIn = AccOut = term()

Options = [Option] | Option

Option = {cache_all, Cache} | cache_all | {max_list_size, MaxListSize} | {tmpdir_usage, TmpFileUsage} | {tmpdir, TempDirectory} | {unique_all, bool()} | unique_all

Error = {error, module(), Reason}

Reason = - as returned by file_sorter(3) -

Calls `Function` on successive answers to the query handle together with an extra argument `AccIn`. The query handle and the function are evaluated in the calling process. `Function` must return a new accumulator which is passed to the next call. `Acc0` is returned if there are no answers to the query handle.

```
1> QH = [1,2,3,4,5,6],
qlc:fold(fun(X, Sum) -> X + Sum end, 0, QH).
21
```

format_error(Error) -> Chars

Types:

Error = {error, module(), term()}

Chars = [char() | Chars]

Returns a descriptive string in English of an error tuple returned by some of the functions of the `qlc` module or the parse transform. This function is mainly used by the compiler invoking the parse transform.

`info(QueryHandleOrList [, Options]) -> Info`

Types:

Options = [Option] | Option

Option = EvalOption | ReturnOption

EvalOption = {cache_all, Cache} | cache_all | {max_list_size, MaxListSize} | {tmpdir_usage, TmpFileUsage} | {tmpdir, TempDirectory} | {unique_all, bool()} | unique_all

ReturnOption = {depth, Depth} | {flat, bool()} | {format, Format} | {n_elements, NElements}

Depth = infinity | int() >= 0

Format = abstract_code | string

NElements = infinity | int() > 0

Info = AbstractExpression | string()

Returns information about a query handle. The information describes the simplifications and optimizations that are the results of preparing the query for evaluation. This function is probably useful mostly during debugging.

The information has the form of an Erlang expression where QLCs most likely occur. Depending on the format functions of mentioned QLC tables it may not be absolutely accurate.

The default is to return a sequence of QLCs in a block, but if the option `{flat, false}` is given, one single QLC is returned. The default is to return a string, but if the option `{format, abstract_code}` is given, abstract code is returned instead. In the abstract code port identifiers, references, and pids are represented by strings. The default is to return all elements in lists, but if the `{n_elements, NElements}` option is given, only a limited number of elements are returned. The default is to show all of objects and match specifications, but if the `{depth, Depth}` option is given, parts of terms below a certain depth are replaced by `'...'`.

```
1> QH = qlc:q([{X,Y} || X <- [x,y], Y <- [a,b]]),
io:format("~s~n", [qlc:info(QH, unique_all)]).
begin
    V1 =
        qlc:q([
               SQV ||
                   SQV <- [x,y]
              ],
              [{unique,true}]),
    V2 =
        qlc:q([
               SQV ||
                   SQV <- [a,b]
              ],
              [{unique,true}]),
    qlc:q([
           {X,Y} ||
               X <- V1,
               Y <- V2
          ],
          [{unique,true}])
end
```

In this example two simple QLCs have been inserted just to hold the `{unique, true}` option.

```
1> E1 = ets:new(e1, []),
E2 = ets:new(e2, []),
true = ets:insert(E1, [{1,a},{2,b}]),
true = ets:insert(E2, [{a,1},{b,2}]),
Q = qlc:q([{X,Z,W} ||
{X, Z} <- ets:table(E1),
{W, Y} <- ets:table(E2),
X =:= Y]),
io:format("~s~n", [qlc:info(Q)]).
begin
    V1 =
        qlc:q([
               P0 ||
                   P0 = {W,Y} <- ets:table(17)
              ]),
    V2 =
        qlc:q([
               [G1|G2] ||
                   G2 <- V1,
                   G1 <- ets:table(16),
                   element(2, G1) =:= element(1, G2)
              ],
              [{join,lookup}]),
    qlc:q([
           {X,Z,W} ||
               [{X,Z}|{W,Y}] <- V2
          ])
end
```

In this example the query list comprehension `V2` has been inserted to show the joined generators and the join method chosen. A convention is used for lookup join: the first generator (`G2`) is the one traversed, the second one (`G1`) is the table where constants are looked up.

keysort(KeyPos, QH1 [, SortOptions]) -> QH2

Types:

QH1 = QueryHandleOrList

QH2 = QueryHandle

Returns a query handle. When evaluating the query handle `QH2` the answers to the query handle `QH1` are sorted by *file_sorter:keysort/4* according to the options.

The sorter will use temporary files only if `QH1` does not evaluate to a list and the size of the binary representation of the answers exceeds `Size` bytes, where `Size` is the value of the `size` option.

next_answers(QueryCursor [, NumberOfAnswers]) -> Answers | Error

Types:

NumberOfAnswers = all_remaining | int() > 0

Error = {error, module(), Reason}

Reason = - as returned by file_sorter(3) -

Returns some or all of the remaining answers to a query cursor. Only the owner of `Cursor` can retrieve answers.

The optional argument `NumberOfAnswers`determines the maximum number of answers returned. The default value is `10`. If less than the requested number of answers is returned, subsequent calls to `next_answers` will return `[]`.

q(QueryListComprehension [, Options]) -> QueryHandle

Types:

QueryListComprehension = - literal query listcomprehension -

Options = [Option] | Option

Option = {max_lookup, MaxLookup} | {cache, Cache} | cache | {join, Join} | {lookup, Lookup} | {unique, bool()} | unique

MaxLookup = int() >= 0 | infinity

Join = any | lookup | merge | nested_loop

Lookup = bool() | any

Returns a query handle for a query list comprehension. The query list comprehension must be the first argument to `qlc:q/1,2` or it will be evaluated as an ordinary list comprehension. It is also necessary to add the line

```
-include_lib("stdlib/include/qlc.hrl").
```

to the source file. This causes a parse transform to substitute a fun for the query list comprehension. The (compiled) fun will be called when the query handle is evaluated.

When calling `qlc:q/1,2` from the Erlang shell the parse transform is automatically called. When this happens the fun substituted for the query list comprehension is not compiled but will be evaluated by `erl_eval(3)`. This is also true when expressions are evaluated by means of `file:eval/1,2` or in the debugger.

To be very explicit, this will not work:

```
...
A = [X || {X} <- [{1},{2}]],
QH = qlc:q(A),
...
```

The variable `A` will be bound to the evaluated value of the list comprehension (`[1,2]`). The compiler complains with an error message ("argument is not a query list comprehension"); the shell process stops with a `badarg` reason.

The `{cache, ets}` option can be used to cache the answers to a query list comprehension. The answers are stored in one ETS table for each cached query list comprehension. When a cached query list comprehension is evaluated again, answers are fetched from the table without any further computations. As a consequence, when all answers to a cached query list comprehension have been found, the ETS tables used for caching answers to the query list comprehension's qualifiers can be emptied. The option `cache` is equivalent to `{cache, ets}`.

The `{cache, list}` option can be used to cache the answers to a query list comprehension just like `{cache, ets}`. The difference is that the answers are kept in a list (on the process heap). If the answers would occupy more than a certain amount of RAM memory a temporary file is used for storing the answers. The option `max_list_size` sets the limit in bytes and the temporary file is put on the directory set by the `tmpdir` option.

The `cache` option has no effect if it is known that the query list comprehension will be evaluated at most once. This is always true for the top-most query list comprehension and also for the list expression of the first generator in a list of qualifiers. Note that in the presence of side effects in filters or callback functions the answers to query list comprehensions can be affected by the `cache` option.

The `{unique, true}` option can be used to remove duplicate answers to a query list comprehension. The unique answers are stored in one ETS table for each query list comprehension. The table is emptied every time it is known that there are no more answers to the query list comprehension. The option `unique` is equivalent to `{unique, true}`. If the `unique` option is combined with the `{cache, ets}` option, two ETS tables are used, but the full answers

are stored in one table only. If the `unique` option is combined with the `{cache, list}` option the answers are sorted twice using `keysort/3`; once to remove duplicates, and once to restore the order.

The `cache` and `unique` options apply not only to the query list comprehension itself but also to the results of looking up constants, running match specifications, and joining handles.

```
1> Q = qlc:q([{A,X,Z,W} ||
A <- [a,b,c],
{X,Z} <- [{a,1},{b,4},{c,6}],
{W,Y} <- [{2,a},{3,b},{4,c}],
X =:= Y],
{cache, list}),
io:format("~s~n", [qlc:info(Q)]).
begin
    V1 =
        qlc:q([
               P0 ||
                   P0 = {X,Z} <-
                       qlc:keysort(1, [{a,1},{b,4},{c,6}], [])
               ]),
    V2 =
        qlc:q([
               P0 ||
                   P0 = {W,Y} <-
                       qlc:keysort(2, [{2,a},{3,b},{4,c}], [])
               ]),
    V3 =
        qlc:q([
               [G1|G2] ||
                   G1 <- V1,
                   G2 <- V2,
                   element(1, G1) == element(2, G2)
               ],
               [{join,merge},{cache,list}]),
    qlc:q([
           {A,X,Z,W} ||
               A <- [a,b,c],
               [{X,Z}|{W,Y}] <- V3,
               X =:= Y
          ])
end
```

In this example the cached results of the merge join are traversed for each value of `A`. Note that without the `cache` option the join would have been carried out three times, once for each value of `A`

`sort/1,2` and `keysort/2,3` can also be used for caching answers and for removing duplicates. When sorting answers are cached in a list, possibly stored on a temporary file, and no ETS tables are used.

Sometimes (see *qlc:table/2* below) traversal of tables can be done by looking up key values, which is assumed to be fast. Under certain (rare) circumstances it could happen that there are too many key values to look up. The `{max_lookup, MaxLookup}` option can then be used to limit the number of lookups: if more than `MaxLookup` lookups would be required no lookups are done but the table traversed instead. The default value is `infinity` which means that there is no limit on the number of keys to look up.

```
1> T = gb_trees:empty(),
QH = qlc:q([X || {{X,Y},_} <- gb_table:table(T),
((X == 1) or (X == 2)) andalso
((Y == a) or (Y == b) or (Y == c))]),
io:format("~s~n", [qlc:info(QH)]).
```

```
ets:match_spec_run(
       lists:flatmap(fun(K) ->
                             case
                                 gb_trees:lookup(K,
                                                 gb_trees:from_orddict([]))
                             of
                                 {value,V} ->
                                     [{K,V}];
                                 none ->
                                     []
                             end
                     end,
                     [{1,a},{1,b},{1,c},{2,a},{2,b},{2,c}]),
       ets:match_spec_compile([{{{'$1','$2'},'_'},[],['$1']}]))
```

In this example using the `gb_table` module from the *Implementing a QLC table* section there are six keys to look up: `{1,a}`, `{1,b}`, `{1,c}`, `{2,a}`, `{2,b}`, and `{2,c}`. The reason is that the two elements of the key {X, Y} are compared separately.

The `{lookup, true}` option can be used to ensure that the `qlc` module will look up constants in some QLC table. If there are more than one QLC table among the generators' list expressions, constants have to be looked up in at least one of the tables. The evaluation of the query fails if there are no constants to look up. This option is useful in situations when it would be unacceptable to traverse all objects in some table. Setting the `lookup` option to `false` ensures that no constants will be looked up (`{max_lookup, 0}` has the same effect). The default value is `any` which means that constants will be looked up whenever possible.

The `{join, Join}` option can be used to ensure that a certain join method will be used: `{join, lookup}` invokes the lookup join method; `{join, merge}` invokes the merge join method; and `{join, nested_loop}` invokes the method of matching every pair of objects from two handles. The last method is mostly very slow. The evaluation of the query fails if the `qlc` module cannot carry out the chosen join method. The default value is `any` which means that some fast join method will be used if possible.

`sort(QH1 [, SortOptions]) -> QH2`

Types:

QH1 = QueryHandleOrList

QH2 = QueryHandle

Returns a query handle. When evaluating the query handle `QH2` the answers to the query handle `QH1` are sorted by *file_sorter:sort/3* according to the options.

The sorter will use temporary files only if `QH1` does not evaluate to a list and the size of the binary representation of the answers exceeds `Size` bytes, where `Size` is the value of the `size` option.

`string_to_handle(QueryString [, Options [, Bindings]]) -> QueryHandle | Error`

Types:

QueryString = string()

Options = [Option] | Option

Option = {max_lookup, MaxLookup} | {cache, Cache} | cache | {join, Join} | {lookup, Lookup} | {unique, bool()} | unique

MaxLookup = int() >= 0 | infinity

Join = any | lookup | merge | nested_loop

Lookup = bool() | any

Bindings = - as returned by erl_eval:bindings/1 -

Error = {error, module(), Reason}

Reason = - ErrorInfo as returned by erl_scan:string/1 or erl_parse:parse_exprs/1 -

A string version of `qlc:q/1,2`. When the query handle is evaluated the fun created by the parse transform is interpreted by `erl_eval(3)`. The query string is to be one single query list comprehension terminated by a period.

```
1> L = [1,2,3],
Bs = erl_eval:add_binding('L', L, erl_eval:new_bindings()),
QH = qlc:string_to_handle("[X+1 || X <- L].", [], Bs),
qlc:eval(QH).
[2,3,4]
```

This function is probably useful mostly when called from outside of Erlang, for instance from a driver written in C.

`table(TraverseFun, Options) -> QueryHandle`

Types:

TraverseFun = TraverseFun0 | TraverseFun1

TraverseFun0 = fun() -> TraverseResult

TraverseFun1 = fun(MatchExpression) -> TraverseResult

TraverseResult = Objects | term()

Objects = [] | [term() | ObjectList]

ObjectList = TraverseFun0 | Objects

Options = [Option] | Option

Option = {format_fun, FormatFun} | {info_fun, InfoFun} | {lookup_fun, LookupFun} | {parent_fun, ParentFun} | {post_fun, PostFun} | {pre_fun, PreFun} | {key_equality, KeyComparison}

FormatFun = undefined | fun(SelectedObjects) -> FormatedTable

SelectedObjects = all | {all, NElements, DepthFun} | {match_spec, MatchExpression} | {lookup, Position, Keys} | {lookup, Position, Keys, NElements, DepthFun}

NElements = infinity | int() > 0

DepthFun = fun(term()) -> term()

FormatedTable = {Mod, Fun, Args} | AbstractExpression | character_list()

InfoFun = undefined | fun(InfoTag) -> InfoValue

InfoTag = indices | is_unique_objects | keypos | num_of_objects

InfoValue = undefined | term()

LookupFun = undefined | fun(Position, Keys) -> LookupResult

LookupResult = [term()] | term()

ParentFun = undefined | fun() -> ParentFunValue

PostFun = undefined | fun() -> void()

PreFun = undefined | fun([PreArg]) -> void()

PreArg = {parent_value, ParentFunValue} | {stop_fun, StopFun}

ParentFunValue = undefined | term()

StopFun = undefined | fun() -> void()

KeyComparison = '=:=' | '=='

Position = int() > 0

Keys = [term()]

Mod = Fun = atom()

Args = [term()]

Returns a query handle for a QLC table. In Erlang/OTP there is support for ETS, Dets and Mnesia tables, but it is also possible to turn many other data structures into QLC tables. The way to accomplish this is to let function(s) in the module implementing the data structure create a query handle by calling `qlc:table/2`. The different ways to traverse the table as well as properties of the table are handled by callback functions provided as options to `qlc:table/2`.

The callback function `TraverseFun` is used for traversing the table. It is to return a list of objects terminated by either `[]` or a nullary fun to be used for traversing the not yet traversed objects of the table. Any other return value is immediately returned as value of the query evaluation. Unary `TraverseFun`s are to accept a match specification as argument. The match specification is created by the parse transform by analyzing the pattern of the generator calling `qlc:table/2` and filters using variables introduced in the pattern. If the parse transform cannot find a match specification equivalent to the pattern and filters, `TraverseFun` will be called with a match specification returning every object. Modules that can utilize match specifications for optimized traversal of tables should call `qlc:table/2` with a unary `TraverseFun` while other modules can provide a nullary `TraverseFun`. `ets:table/2` is an example of the former; `gb_table:table/1` in the *Implementing a QLC table* section is an example of the latter.

`PreFun` is a unary callback function that is called once before the table is read for the first time. If the call fails, the query evaluation fails. Similarly, the nullary callback function `PostFun` is called once after the table was last read. The return value, which is caught, is ignored. If `PreFun` has been called for a table, `PostFun` is guaranteed to be called for that table, even if the evaluation of the query fails for some reason. The order in which pre (post) functions for different tables are evaluated is not specified. Other table access than reading, such as calling `InfoFun`, is assumed to be OK at any time. The argument `PreArgs` is a list of tagged values. Currently there are two tags, `parent_value` and `stop_fun`, used by Mnesia for managing transactions. The value of `parent_value` is the value returned by `ParentFun`, or `undefined` if there is no `ParentFun`. `ParentFun` is called once just before the call of `PreFun` in the context of the process calling `eval`, `fold`, or `cursor`. The value of `stop_fun` is a nullary fun that deletes the cursor if called from the parent, or `undefined` if there is no cursor.

The binary callback function `LookupFun` is used for looking up objects in the table. The first argument `Position` is the key position or an indexed position and the second argument `Keys` is a sorted list of unique values. The return value is to be a list of all objects (tuples) such that the element at `Position` is a member of `Keys`. Any other return value is immediately returned as value of the query evaluation. `LookupFun` is called instead of traversing the table if the parse transform at compile time can find out that the filters match and compare the element at `Position` in such a way that only `Keys` need to be looked up in order to find all potential answers. The key position is obtained by calling `InfoFun(keypos)` and the indexed positions by calling `InfoFun(indices)`. If the key position can be used for lookup it is always chosen, otherwise the indexed position requiring the least number of lookups is chosen. If there is a tie between two indexed positions the one occurring first in the list returned by `InfoFun` is chosen. Positions requiring more than *max_lookup* lookups are ignored.

The unary callback function `InfoFun` is to return information about the table. `undefined` should be returned if the value of some tag is unknown:

- `indices`. Returns a list of indexed positions, a list of positive integers.
- `is_unique_objects`. Returns `true` if the objects returned by `TraverseFun` are unique.
- `keypos`. Returns the position of the table's key, a positive integer.
- `is_sorted_key`. Returns `true` if the objects returned by `TraverseFun` are sorted on the key.
- `num_of_objects`. Returns the number of objects in the table, a non-negative integer.

The unary callback function `FormatFun` is used by *qlc:info/1,2* for displaying the call that created the table's query handle. The default value, `undefined`, means that `info/1,2` displays a call to `'$MOD':'$FUN'/0`. It is up to `FormatFun` to present the selected objects of the table in a suitable way. However, if a character list is chosen for presentation it must be an Erlang expression that can be scanned and parsed (a trailing dot will be added by `qlc:info` though). `FormatFun` is called with an argument that describes the selected objects based on optimizations done as a result of analyzing the filters of the QLC where the call to `qlc:table/2` occurs. The possible values of the argument are:

- `{lookup, Position, Keys, NElements, DepthFun}`. `LookupFun` is used for looking up objects in the table.
- `{match_spec, MatchExpression}`. No way of finding all possible answers by looking up keys was found, but the filters could be transformed into a match specification. All answers are found by calling `TraverseFun(MatchExpression)`.
- `{all, NElements, DepthFun}`. No optimization was found. A match specification matching all objects will be used if `TraverseFun` is unary.

`NElements` is the value of the `info/1,2` option `n_elements`, and `DepthFun` is a function that can be used for limiting the size of terms; calling `DepthFun(Term)` substitutes `'...'` for parts of `Term` below the depth specified by the `info/1,2` option `depth`. If calling `FormatFun` with an argument including `NElements` and `DepthFun` fails, `FormatFun` is called once again with an argument excluding `NElements` and `DepthFun` (`{lookup, Position, Keys}` or `all`).

The value of `key_equality` is to be `'=:='` if the table considers two keys equal if they match, and to be `'=='` if two keys are equal if they compare equal. The default is `'=:='`.

See *ets(3)*, *dets(3)* and *mnesia(3)* for the various options recognized by `table/1,2` in respective module.

See Also

dets(3), *Erlang Reference Manual*, *erl_eval(3)*, *erlang(3)*, *ets(3)*, *file(3)*, *error_logger(3)*, *file_sorter(3)*, *mnesia(3)*, *Programming Examples*, *shell(3)*

queue

Erlang module

This module implements (double ended) FIFO queues in an efficient manner.

All functions fail with reason `badarg` if arguments are of wrong type, for example queue arguments are not queues, indexes are not integers, list arguments are not lists. Improper lists cause internal crashes. An index out of range for a queue also causes a failure with reason `badarg`.

Some functions, where noted, fail with reason `empty` for an empty queue.

The data representing a queue as used by this module should be regarded as opaque by other modules. Any code assuming knowledge of the format is running on thin ice.

All operations has an amortized O(1) running time, except `len/1`, `join/2`, `split/2`, `filter/2` and `member/2` that have O(n). To minimize the size of a queue minimizing the amount of garbage built by queue operations, the queues do not contain explicit length information, and that is why `len/1` is O(n). If better performance for this particular operation is essential, it is easy for the caller to keep track of the length.

Queues are double ended. The mental picture of a queue is a line of people (items) waiting for their turn. The queue front is the end with the item that has waited the longest. The queue rear is the end an item enters when it starts to wait. If instead using the mental picture of a list, the front is called head and the rear is called tail.

Entering at the front and exiting at the rear are reverse operations on the queue.

The module has several sets of interface functions. The "Original API", the "Extended API" and the "Okasaki API".

The "Original API" and the "Extended API" both use the mental picture of a waiting line of items. Both also have reverse operations suffixed "_r".

The "Original API" item removal functions return compound terms with both the removed item and the resulting queue. The "Extended API" contain alternative functions that build less garbage as well as functions for just inspecting the queue ends. Also the "Okasaki API" functions build less garbage.

The "Okasaki API" is inspired by "Purely Functional Data structures" by Chris Okasaki. It regards queues as lists. The API is by many regarded as strange and avoidable. For example many reverse operations have lexically reversed names, some with more readable but perhaps less understandable aliases.

Original API

Exports

`new() -> Q`

Types:

Q = queue()

Returns an empty queue.

`is_queue(Term) -> true | false`

Types:

Term = term()

Tests if `Q` is a queue and returns `true` if so and `false` otherwise.

is_empty(Q) -> true | false

Types:

Q = queue()

Tests if `Q` is empty and returns `true` if so and `false` otherwise.

len(Q) -> N

Types:

Q = queue()

N = integer()

Calculates and returns the length of queue `Q`.

in(Item, Q1) -> Q2

Types:

Item = term()

Q1 = Q2 = queue()

Inserts `Item` at the rear of queue `Q1`. Returns the resulting queue `Q2`.

in_r(Item, Q1) -> Q2

Types:

Item = term()

Q1 = Q2 = queue()

Inserts `Item` at the front of queue `Q1`. Returns the resulting queue `Q2`.

out(Q1) -> Result

Types:

Result = {{value, Item}, Q2} | {empty, Q1}

Q1 = Q2 = queue()

Removes the item at the front of queue `Q1`. Returns the tuple `{{value, Item}, Q2}`, where `Item` is the item removed and `Q2` is the resulting queue. If `Q1` is empty, the tuple `{empty, Q1}` is returned.

out_r(Q1) -> Result

Types:

Result = {{value, Item}, Q2} | {empty, Q1}

Q1 = Q2 = queue()

Removes the item at the rear of the queue `Q1`. Returns the tuple `{{value, Item}, Q2}`, where `Item` is the item removed and `Q2` is the new queue. If `Q1` is empty, the tuple `{empty, Q1}` is returned.

from_list(L) -> queue()

Types:

L = list()

Returns a queue containing the items in `L` in the same order; the head item of the list will become the front item of the queue.

to_list(Q) -> list()

Types:

Q = queue()

Returns a list of the items in the queue in the same order; the front item of the queue will become the head of the list.

reverse(Q1) -> Q2

Types:

Q1 = Q2 = queue()

Returns a queue `Q2` that contains the items of `Q1` in the reverse order.

split(N, Q1) -> {Q2,Q3}

Types:

N = integer()

Q1 = Q2 = Q3 = queue()

Splits `Q1` in two. The `N` front items are put in `Q2` and the rest in `Q3`

join(Q1, Q2) -> Q3

Types:

Q1 = Q2 = Q3 = queue()

Returns a queue `Q3` that is the result of joining `Q1` and `Q2` with `Q1` in front of `Q2`.

filter(Fun, Q1) -> Q2

Types:

Fun = fun(Item) -> bool() | list()

Q1 = Q2 = queue()

Returns a queue `Q2` that is the result of calling `Fun(Item)` on all items in `Q1`, in order from front to rear.

If `Fun(Item)` returns `true`, `Item` is copied to the result queue. If it returns `false`, `Item` is not copied. If it returns a list the list elements are inserted instead of `Item` in the result queue.

So, `Fun(Item)` returning `[Item]` is thereby semantically equivalent to returning `true`, just as returning `[]` is semantically equivalent to returning `false`. But returning a list builds more garbage than returning an atom.

member(Item, Q) -> bool()

Types:

Item = term()

Q = queue()

Returns `true` if `Item` matches some element in `Q`, otherwise `false`.

Extended API

Exports

get(Q) -> Item

Types:

Item = **term()**

Q = **queue()**

Returns `Item` at the front of queue `Q`.

Fails with reason `empty` if `Q` is empty.

`get_r(Q) -> Item`

Types:

Item = **term()**

Q = **queue()**

Returns `Item` at the rear of queue `Q`.

Fails with reason `empty` if `Q` is empty.

`drop(Q1) -> Q2`

Types:

Item = **term()**

Q1 = **Q2** = **queue()**

Returns a queue `Q2` that is the result of removing the front item from `Q1`.

Fails with reason `empty` if `Q1` is empty.

`drop_r(Q1) -> Q2`

Types:

Item = **term()**

Q1 = **Q2** = **queue()**

Returns a queue `Q2` that is the result of removing the rear item from `Q1`.

Fails with reason `empty` if `Q1` is empty.

`peek(Q) -> {value,Item} | empty`

Types:

Item = **term()**

Q = **queue()**

Returns the tuple `{value, Item}` where `Item` is the front item of `Q`, or `empty` if `Q1` is empty.

`peek_r(Q) -> {value,Item} | empty`

Types:

Item = **term()**

Q = **queue()**

Returns the tuple `{value, Item}` where `Item` is the rear item of `Q`, or `empty` if `Q1` is empty.

Okasaki API

Exports

`cons(Item, Q1) -> Q2`

Types:

Item = **term()**

Q1 = **Q2** = **queue()**

Inserts `Item` at the head of queue `Q1`. Returns the new queue `Q2`.

`head(Q) -> Item`

Types:

Item = **term()**

Q = **queue()**

Returns `Item` from the head of queue `Q`.

Fails with reason `empty` if `Q` is empty.

`tail(Q1) -> Q2`

Types:

Item = **term()**

Q1 = **Q2** = **queue()**

Returns a queue `Q2` that is the result of removing the head item from `Q1`.

Fails with reason `empty` if `Q1` is empty.

`snoc(Q1, Item) -> Q2`

Types:

Item = **term()**

Q1 = **Q2** = **queue()**

Inserts `Item` as the tail item of queue `Q1`. Returns the new queue `Q2`.

`daeh(Q) -> Item`

`last(Q) -> Item`

Types:

Item = **term()**

Q = **queue()**

Returns the tail item of queue `Q`.

Fails with reason `empty` if `Q` is empty.

liat(Q1) -> Q2

init(Q1) -> Q2

lait(Q1) -> Q2

Types:

Item = term()

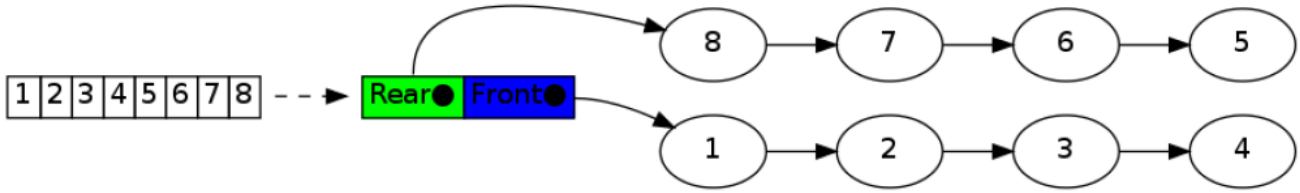
Q1 = Q2 = queue()

Returns a queue `Q2` that is the result of removing the tail item from `Q1`.

Fails with reason `empty` if `Q1` is empty.

The name `lait/1` is a misspelling - do not use it anymore.

- 为了实现快速的双向入队、出队操作，queue 基于论文 **Purely Functional Data Structures** 中描述的算法构建；
- queue 被表示为 2 元 tuple，第一个元素是反向排列的尾部元素列表，第二个元素是正向排列的头部元素列表，二者合在一起就是完整的队列；
- queue 内部结构如下图所示：



random

Erlang module

Random number generator. The method is attributed to B.A. Wichmann and I.D.Hill, in 'An efficient and portable pseudo-random number generator', Journal of Applied Statistics. AS183. 1982. Also Byte March 1987.

The current algorithm is a modification of the version attributed to Richard A O'Keefe in the standard Prolog library.

Every time a random number is requested, a state is used to calculate it, and a new state produced. The state can either be implicit (kept in the process dictionary) or be an explicit argument and return value. In this implementation, the state (the type `ran()`) consists of a tuple of three integers.

It should be noted that this random number generator is not cryptographically strong. If a strong cryptographic random number generator is needed for example `crypto:rand_bytes/1` could be used instead.

Exports

`seed() -> ran()`

Seeds random number generation with default (fixed) values in the process dictionary, and returns the old state.

`seed(A1, A2, A3) -> undefined | ran()`

Types:

A1 = **A2** = **A3** = **integer()**

Seeds random number generation with integer values in the process dictionary, and returns the old state.

One way of obtaining a seed is to use the BIF `now/0`:

```
...
{A1,A2,A3} = now(),
random:seed(A1, A2, A3),
...
```

`seed({A1, A2, A3}) -> undefined | ran()`

Types:

A1 = **A2** = **A3** = **integer()**

`seed({A1, A2, A3})` is equivalent to `seed(A1, A2, A3)`.

`seed0() -> ran()`

Returns the default state.

`uniform()-> float()`

Returns a random float uniformly distributed between `0.0` and `1.0`, updating the state in the process dictionary.

`uniform(N) -> integer()`

Types:

N = **integer()**

Given an integer `N >= 1`, `uniform/1` returns a random integer uniformly distributed between `1` and `N`, updating the state in the process dictionary.

`uniform_s(State0) -> {float(), State1}`

Types:

State0 = **State1** = **ran()**

Given a state, `uniform_s/1`returns a random float uniformly distributed between `0.0` and `1.0`, and a new state.

`uniform_s(N, State0) -> {integer(), State1}`

Types:

N = **integer()**

State0 = **State1** = **ran()**

Given an integer `N >= 1` and a state, `uniform_s/2` returns a random integer uniformly distributed between `1` and `N`, and a new state.

Note

Some of the functions use the process dictionary variable `random_seed` to remember the current seed.

If a process calls `uniform/0` or `uniform/1` without setting a seed first, `seed/0` is called automatically.

re

Erlang module

This module contains regular expression matching functions for strings and binaries.

The regular expression syntax and semantics resemble that of Perl. This library replaces the deprecated pure-Erlang regexp library; it has a richer syntax, more options and is faster.

The library's matching algorithms are currently based on the PCRE library, but not all of the PCRE library is interfaced and some parts of the library go beyond what PCRE offers. The sections of the PCRE documentation which are relevant to this module are included here.

> ### Note:
>
> The Erlang literal syntax for strings uses the "\" (backslash) character as an escape code. You need to escape backslashes in literal strings, both in your code and in the shell, with an additional backslash, i.e.: "\\".

DATA TYPES

```
iodata() = iolist() | binary()
iolist() = [char() | binary() | iolist()]
  - a binary is allowed as the tail of the list
```

```
unicode_binary() = binary() with characters encoded in UTF-8 coding standard
unicode_char() = integer() representing a valid unicode codepoint

chardata() = charlist() | unicode_binary()

charlist() = [unicode_char() | unicode_binary() | charlist()]
  - a unicode_binary is allowed as the tail of the list
```

```
mp() = Opaque datatype containing a compiled regular expression.
  - The mp() is guaranteed to be a tuple() having the atom
're_pattern' as its first element, to allow for matching in
     guards. The arity of the tuple() or the content of the other fields
may change in future releases.
```

Exports

`compile(Regexp) -> {ok, MP} | {error, ErrSpec}`

Types:

Regexp = iodata()

The same as `compile(Regexp,[])`

`compile(Regexp,Options) -> {ok, MP} | {error, ErrSpec}`

Types:

Regexp = iodata() | charlist()

Options = [Option]

Option = unicode | anchored | caseless | dollar_endonly | dotall | extended | firstline | multiline | no_auto_capture | dupnames | ungreedy | {newline, NLSpec}| bsr_anycrlf | bsr_unicode

NLSpec = cr | crlf | lf | anycrlf | any

MP = mp()

ErrSpec = {ErrString, Position}

ErrString = string()

Position = int()

This function compiles a regular expression with the syntax described below into an internal format to be used later as a parameter to the run/2,3 functions.

Compiling the regular expression before matching is useful if the same expression is to be used in matching against multiple subjects during the program's lifetime. Compiling once and executing many times is far more efficient than compiling each time one wants to match.

When the unicode option is given, the regular expression should be given as a valid unicode `charlist()`, otherwise as any valid `iodata()`.

The options have the following meanings:

`unicode`
: The regular expression is given as a unicode `charlist()` and the resulting regular expression code is to be run against a valid unicode `charlist()` subject.

`anchored`
: The pattern is forced to be "anchored", that is, it is constrained to match only at the first matching point in the string that is being searched (the "subject string"). This effect can also be achieved by appropriate constructs in the pattern itself.

`caseless`
: Letters in the pattern match both upper and lower case letters. It is equivalent to Perl's /i option, and it can be changed within a pattern by a (?i) option setting. Uppercase and lowercase letters are defined as in the ISO-8859-1 character set.

`dollar_endonly`
: A dollar metacharacter in the pattern matches only at the end of the subject string. Without this option, a dollar also matches immediately before a newline at the end of the string (but not before any other newlines). The `dollar_endonly` option is ignored if `multiline` is given. There is no equivalent option in Perl, and no way to set it within a pattern.

`dotall`
: A dot in the pattern matches all characters, including those that indicate newline. Without it, a dot does not match when the current position is at a newline. This option is equivalent to Perl's /s option, and it can be changed within a pattern by a (?s) option setting. A negative class such as [^a] always matches newline characters, independent of this option's setting.

`extended`
: Whitespace data characters in the pattern are ignored except when escaped or inside a character class. Whitespace does not include the VT character (ASCII 11). In addition, characters between an unescaped # outside a character class and the next newline, inclusive, are also ignored. This is equivalent to Perl's /x option, and it can be changed within a pattern by a (?x) option setting. This option makes it possible to include comments inside complicated patterns. Note, however, that this applies only to data characters. Whitespace characters may never appear within special character sequences in a pattern, for example within the sequence `(?(` which introduces a conditional subpattern.

`firstline`
An unanchored pattern is required to match before or at the first newline in the subject string, though the matched text may continue over the newline.

`multiline`
By default, PCRE treats the subject string as consisting of a single line of characters (even if it actually contains newlines). The "start of line" metacharacter (^) matches only at the start of the string, while the "end of line" metacharacter ($) matches only at the end of the string, or before a terminating newline (unless `dollar_endonly` is given). This is the same as Perl.

When `multiline` it is given, the "start of line" and "end of line" constructs match immediately following or immediately before internal newlines in the subject string, respectively, as well as at the very start and end. This is equivalent to Perl's /m option, and it can be changed within a pattern by a (?m) option setting. If there are no newlines in a subject string, or no occurrences of ^ or $ in a pattern, setting `multiline` has no effect.

`no_auto_capture`
Disables the use of numbered capturing parentheses in the pattern. Any opening parenthesis that is not followed by ? behaves as if it were followed by ?: but named parentheses can still be used for capturing (and they acquire numbers in the usual way). There is no equivalent of this option in Perl.

`dupnames`
Names used to identify capturing subpatterns need not be unique. This can be helpful for certain types of pattern when it is known that only one instance of the named subpattern can ever be matched. There are more details of named subpatterns below

`ungreedy`
This option inverts the "greediness" of the quantifiers so that they are not greedy by default, but become greedy if followed by "?". It is not compatible with Perl. It can also be set by a (?U) option setting within the pattern.

`{newline, NLSpec}`
Override the default definition of a newline in the subject string, which is LF (ASCII 10) in Erlang.

`cr`
Newline is indicated by a single character CR (ASCII 13)

`lf`
Newline is indicated by a single character LF (ASCII 10), the default

`crlf`
Newline is indicated by the two-character CRLF (ASCII 13 followed by ASCII 10) sequence.

`anycrlf`
Any of the three preceding sequences should be recognized.

`any`
Any of the newline sequences above, plus the Unicode sequences VT (vertical tab, U+000B), FF (formfeed, U+000C), NEL (next line, U+0085), LS (line separator, U+2028), and PS (paragraph separator, U+2029).

`bsr_anycrlf`
Specifies specifically that \R is to match only the cr, lf or crlf sequences, not the Unicode specific newline characters.

`bsr_unicode`
Specifies specifically that \R is to match all the Unicode newline characters (including crlf etc, the default).

`run(Subject,RE) -> {match, Captured} | nomatch`

Types:

Subject = iodata() | charlist()

RE = mp() | iodata() | charlist()

Captured = [CaptureData]

CaptureData = {int(),int()}

The same as `run(Subject,RE,[])`.

`run(Subject,RE,Options) -> {match, Captured} | match | nomatch`

Types:

Subject = iodata() | charlist()

RE = mp() | iodata() | charlist()

Options = [Option]

Option = anchored | global | notbol | noteol | notempty | {offset, int()} | {newline, NLSpec} | bsr_anycrlf | bsr_unicode | {capture, ValueSpec} | {capture, ValueSpec, Type} | CompileOpt

Type = index | list | binary

ValueSpec = all | all_but_first | first | none | ValueList

ValueList = [ValueID]

ValueID = int() | string() | atom()

CompileOpt = see compile/2 above

NLSpec = cr | crlf | lf | anycrlf | any

Captured = [CaptureData] | [[CaptureData] ...]

CaptureData = {int(),int()} | ListConversionData | binary()

ListConversionData = string() | {error, string(), binary()} | {incomplete, string(), binary()}

Executes a regexp matching, returning `match/{match, Captured}` or `nomatch`. The regular expression can be given either as `iodata()` in which case it is automatically compiled (as by `re:compile/2`) and executed, or as a pre compiled `mp()` in which case it is executed against the subject directly.

When compilation is involved, the exception `badarg` is thrown if a compilation error occurs. Call `re:compile/2` to get information about the location of the error in the regular expression.

If the regular expression is previously compiled, the option list can only contain the options `anchored`, `global`, `notbol`, `noteol`, `notempty`, `{offset, int()}`, `{newline, NLSpec}` and `{capture, ValueSpec}/{capture, ValueSpec, Type}`. Otherwise all options valid for the `re:compile/2` function are allowed as well. Options allowed both for compilation and execution of a match, namely `anchored` and `{newline, NLSpec}`, will affect both the compilation and execution if present together with a non pre-compiled regular expression.

If the regular expression was previously compiled with the option `unicode`, the `Subject` should be provided as a valid Unicode `charlist()`, otherwise any `iodata()` will do. If compilation is involved and the option `unicode` is given, both the `Subject` and the regular expression should be given as valid Unicode `charlists()`.

The `{capture, ValueSpec}/{capture, ValueSpec, Type}` defines what to return from the function upon successful matching. The `capture` tuple may contain both a value specification telling which of the captured substrings are to be returned, and a type specification, telling how captured substrings are to be returned (as index tuples, lists or binaries). The `capture` option makes the function quite flexible and powerful. The different options are described in detail below.

If the capture options describe that no substring capturing at all is to be done (`{capture, none}`), the function will return the single atom `match` upon successful matching, otherwise the tuple `{match, ValueList}` is returned. Disabling capturing can be done either by specifying `none` or an empty list as `ValueSpec`.

The options relevant for execution are:

`anchored`

Limits `re:run/3` to matching at the first matching position. If a pattern was compiled with `anchored`, or turned out to be anchored by virtue of its contents, it cannot be made unanchored at matching time, hence there is no `unanchored` option.

`global`

Implements global (repetitive) search (the `g` flag in Perl). Each match is returned as a separate `list()` containing the specific match as well as any matching subexpressions (or as specified by the `capture option`). The `Captured` part of the return value will hence be a `list()` of `list()`s when this option is given.

The interaction of the global option with a regular expression which matches an empty string surprises some users. When the global option is given, `re:run/3` handles empty matches in the same way as Perl: a zero-length match at any point will be retried with the options `[anchored, notempty]` as well. If that search gives a result of length > 0, the result is included. For example:

```
re:run("cat","(|at)",[global]).
```

The following matching will be performed:

At offset `0`
The regexp `(|at)` will first match at the initial position of the string `cat`, giving the result set `[{0,0}, {0,0}]` (the second `{0,0}` is due to the subexpression marked by the parentheses). As the length of the match is 0, we don't advance to the next position yet.
At offset `0` with `[anchored, notempty]`
The search is retried with the options `[anchored, notempty]` at the same position, which does not give any interesting result of longer length, so the search position is now advanced to the next character (`a`).
At offset `1`
This time, the search results in `[{1,0},{1,0}]`, so this search will also be repeated with the extra options.
At offset `1` with `[anchored, notempty]`
Now the `ab` alternative is found and the result will be [{1,2},{1,2}]. The result is added to the list of results and the position in the search string is advanced two steps.
At offset `3`
The search now once again matches the empty string, giving `[{3,0},{3,0}]`.
At offset `1` with `[anchored, notempty]`
This will give no result of length > 0 and we are at the last position, so the global search is complete.

The result of the call is:

```
{match,[[{0,0},{0,0}],[{1,0},{1,0}],[{1,2},{1,2}],[{3,0},{3,0}]]}
```

`notempty`

An empty string is not considered to be a valid match if this option is given. If there are alternatives in the pattern, they are tried. If all the alternatives match the empty string, the entire match fails. For example, if the pattern

```
a?b?
```

is applied to a string not beginning with "a" or "b", it would normally match the empty string at the start of the subject. With the `notempty` option, this match is not valid, so re:run/3 searches further into the string for occurrences of "a" or "b".

Perl has no direct equivalent of `notempty`, but it does make a special case of a pattern match of the empty string within its split() function, and when using the /g modifier. It is possible to emulate Perl's behavior after matching a null string by first trying the match again at the same offset with `notempty` and `anchored`, and then, if that fails, by advancing the starting offset (see below) and trying an ordinary match again.

`notbol`

This option specifies that the first character of the subject string is not the beginning of a line, so the circumflex metacharacter should not match before it. Setting this without `multiline` (at compile time) causes

circumflex never to match. This option only affects the behavior of the circumflex metacharacter. It does not affect \A.

`noteol`

This option specifies that the end of the subject string is not the end of a line, so the dollar metacharacter should not match it nor (except in multiline mode) a newline immediately before it. Setting this without `multiline` (at compile time) causes dollar never to match. This option affects only the behavior of the dollar metacharacter. It does not affect \Z or \z.

`{offset, int()}`

Start matching at the offset (position) given in the subject string. The offset is zero-based, so that the default is `{offset,0}` (all of the subject string).

`{newline, NLSpec}`

Override the default definition of a newline in the subject string, which is LF (ASCII 10) in Erlang.

`cr`

Newline is indicated by a single character CR (ASCII 13)

`lf`

Newline is indicated by a single character LF (ASCII 10), the default

`crlf`

Newline is indicated by the two-character CRLF (ASCII 13 followed by ASCII 10) sequence.

`anycrlf`

Any of the three preceding sequences should be recognized.

`any`

Any of the newline sequences above, plus the Unicode sequences VT (vertical tab, U+000B), FF (formfeed, U+000C), NEL (next line, U+0085), LS (line separator, U+2028), and PS (paragraph separator, U+2029).

`bsr_anycrlf`

Specifies specifically that \R is to match only the cr, lf or crlf sequences, not the Unicode specific newline characters. (overrides compilation option)

`bsr_unicode`

Specifies specifically that \R is to match all the Unicode newline characters (including crlf etc, the default). (overrides compilation option)

`{capture, ValueSpec}`/`{capture, ValueSpec, Type}`

Specifies which captured substrings are returned and in what format. By default, `re:run/3` captures all of the matching part of the substring as well as all capturing subpatterns (all of the pattern is automatically captured). The default return type is (zero-based) indexes of the captured parts of the string, given as `{Offset,Length}` pairs (the `index` `Type` of capturing).

As an example of the default behavior, the following call:

```
re:run("ABCabcdABC","abcd",[]).
```

returns, as first and only captured string the matching part of the subject ("abcd" in the middle) as a index pair `{3,4}`, where character positions are zero based, just as in offsets. The return value of the call above would then be:

```
{match,[{3,4}]}
```

Another (and quite common) case is where the regular expression matches all of the subject, as in:

```
re:run("ABCabcdABC",".*abcd.*",[]).
```

where the return value correspondingly will point out all of the string, beginning at index 0 and being 10 characters long:

```
{match,[{0,10}]}
```

If the regular expression contains capturing subpatterns, like in the following case:

```
re:run("ABCabcdABC",".*(abcd).*",[]).
```

all of the matched subject is captured, as well as the captured substrings:

```
{match,[{0,10},{3,4}]}
```

the complete matching pattern always giving the first return value in the list and the rest of the subpatterns being added in the order they occurred in the regular expression.

The capture tuple is built up as follows:

`ValueSpec`

Specifies which captured (sub)patterns are to be returned. The ValueSpec can either be an atom describing a predefined set of return values, or a list containing either the indexes or the names of specific subpatterns to return.

The predefined sets of subpatterns are:

`all`
All captured subpatterns including the complete matching string. This is the default.
`first`
Only the first captured subpattern, which is always the complete matching part of the subject. All explicitly captured subpatterns are discarded.
`all_but_first`
All but the first matching subpattern, i.e. all explicitly captured subpatterns, but not the complete matching part of the subject string. This is useful if the regular expression as a whole matches a large part of the subject, but the part you're interested in is in an explicitly captured subpattern. If the return type is `list` or `binary`, not returning subpatterns you're not interested in is a good way to optimize.
`none`
Do not return matching subpatterns at all, yielding the single atom `match` as the return value of the function when matching successfully instead of the `{match, list()}` return. Specifying an empty list gives the same behavior.

The value list is a list of indexes for the subpatterns to return, where index 0 is for all of the pattern, and 1 is for the first explicit capturing subpattern in the regular expression, and so forth. When using named captured subpatterns (see below) in the regular expression, one can use `atom()`s or `string()`s to specify the subpatterns to be returned. For example, consider the regular expression:

```
".*(abcd).*"
```

matched against the string ""ABCabcdABC", capturing only the "abcd" part (the first explicit subpattern):

```
re:run("ABCabcdABC",".*(abcd).*",[{capture,[1]}]).
```

The call will yield the following result:

```
{match,[{3,4}]}
```

as the first explicitly captured subpattern is "(abcd)", matching "abcd" in the subject, at (zero-based) position 3, of length 4.

Now consider the same regular expression, but with the subpattern explicitly named 'FOO':

```
".*(?<FOO>abcd).*"
```

With this expression, we could still give the index of the subpattern with the following call:

```
re:run("ABCabcdABC",".*(?<FOO>abcd).*",[{capture,[1]}]).
```

giving the same result as before. But, since the subpattern is named, we can also specify its name in the value list:

```
re:run("ABCabcdABC",".*(?<FOO>abcd).*",[{capture,['FOO']}]).
```

which would yield the same result as the earlier examples, namely:

```
{match,[{3,4}]}
```

The values list might specify indexes or names not present in the regular expression, in which case the return values vary depending on the type. If the type is `index`, the tuple `{-1,0}` is returned for values having no corresponding subpattern in the regexp, but for the other types (`binary` and `list`), the values are the empty binary or list respectively.

`Type`

Optionally specifies how captured substrings are to be returned. If omitted, the default of `index` is used. The `Type` can be one of the following:

`index`
Return captured substrings as pairs of byte indexes into the subject string and length of the matching string in the subject (as if the subject string was flattened with `iolist_to_binary/1` or `unicode:characters_to_binary/2` prior to matching). Note that the `unicode` option results in *byte-oriented* indexes in a (possibly virtual) *UTF-8 encoded* binary. A byte index tuple `{0,2}` might therefore represent one or two characters when `unicode` is in effect. This might seem counter-intuitive, but has been deemed the most effective and useful way to way to do it. To return lists instead might result in simpler code if that is desired. This return type is the default.

`list`
Return matching substrings as lists of characters (Erlang `string()`s). It the `unicode` option is used in combination with the \C sequence in the regular expression, a captured subpattern can contain bytes that are not valid UTF-8 (\C matches bytes regardless of character encoding). In that case the `list` capturing may result in the same types of tuples that `unicode:characters_to_list/2` can return, namely three-tuples with the tag `incomplete` or `error`, the successfully converted characters and the invalid UTF-8 tail of the conversion as a binary. The best strategy is to avoid using the \C sequence when capturing lists.

`binary`
Return matching substrings as binaries. If the `unicode` option is used, these binaries are in UTF-8. If the \C sequence is used together with `unicode` the binaries may be invalid UTF-8.

In general, subpatterns that were not assigned a value in the match are returned as the tuple `{-1,0}` when `type` is `index`. Unassigned subpatterns are returned as the empty binary or list, respectively, for other return types. Consider the regular expression:

```
".*((?<FOO>abdd)|a(..d)).*"
```

There are three explicitly capturing subpatterns, where the opening parenthesis position determines the order in the result, hence `((?<FOO>abdd)|a(..d))` is subpattern index 1, `(?<FOO>abdd)` is subpattern index 2 and `(..d)` is subpattern index 3. When matched against the following string:

```
"ABCabcdABC"
```

the subpattern at index 2 won't match, as "abdd" is not present in the string, but the complete pattern matches (due to the alternative `a(..d)`. The subpattern at index 2 is therefore unassigned and the default return value will be:

```
{match,[{0,10},{3,4},{-1,0},{4,3}]}
```

Setting the capture `Type` to `binary` would give the following:

```
{match,[<<"ABCabcdABC">>,<<"abcd">>,<<>>,<<"bcd">>]}
```

where the empty binary (`<<>>`) represents the unassigned subpattern. In the `binary` case, some information about the matching is therefore lost, the `<<>>` might just as well be an empty string captured.

If differentiation between empty matches and non existing subpatterns is necessary, use the `type index` and do the conversion to the final type in Erlang code.

When the option `global` is given, the `capture` specification affects each match separately, so that:

```
re:run("cacb","c(a|b)",[global,{capture,[1],list}]).
```

gives the result:

```
{match,[["a"],["b"]]}
```

The options solely affecting the compilation step are described in the `re:compile/2` function.

replace(Subject,RE,Replacement) -> iodata() | charlist()

Types:

Subject = iodata() | charlist()

RE = mp() | iodata()

Replacement = iodata() | charlist()

The same as `replace(Subject,RE,Replacement,[])`.

replace(Subject,RE,Replacement,Options) -> iodata() | charlist() | binary() | list()

Types:

Subject = iodata() | charlist()

RE = mp() | iodata() | charlist()

Replacement = iodata() | charlist()

Options = [Option]

Option = anchored | global | notbol | noteol | notempty | {offset, int()} | {newline, NLSpec} | bsr_anycrlf | bsr_unicode | {return, ReturnType} | CompileOpt

ReturnType = iodata | list | binary

CompileOpt = see compile/2 above

NLSpec = cr | crlf | lf | anycrlf | any

Replaces the matched part of the `Subject` string with the contents of `Replacement`.

The permissible options are the same as for `re:run/3`, except that the `capture` option is not allowed. Instead a `{return, ReturnType}` is present. The default return type is `iodata`, constructed in a way to minimize copying. The `iodata` result can be used directly in many i/o-operations. If a flat `list()` is desired, specify `{return, list}` and if a binary is preferred, specify `{return, binary}`.

As in the `re:run/3` function, an `mp()` compiled with the `unicode` option requires the `Subject` to be a Unicode `charlist()`. If compilation is done implicitly and the `unicode` compilation option is given to this function, both the regular expression and the `Subject` should be given as valid Unicode `charlist()`s.

The replacement string can contain the special character &, which inserts the whole matching expression in the result, and the special sequence \N (where N is an integer > 0), resulting in the subexpression number N will be inserted in the result. If no subexpression with that number is generated by the regular expression, nothing is inserted.

To insert an & or \ in the result, precede it with a \. Note that Erlang already gives a special meaning to \ in literal strings, so a single \ has to be written as `"\\"` and therefore a double \ as `"\\\\"`. Example:

```
re:replace("abcd","c","[&]",[{return,list}]).
```

gives

```
"ab[c]d"
```

while

```
re:replace("abcd","c","[\\&]",[{return,list}]).
```

gives

```
"ab[&]d"
```

As with `re:run/3`, compilation errors raise the `badarg` exception, `re:compile/2` can be used to get more information about the error.

`split(Subject,RE) -> SplitList`

Types:

Subject = iodata() | charlist()

RE = mp() | iodata()

SplitList = [iodata() | charlist()]

The same as `split(Subject,RE,[])`.

`split(Subject,RE,Options) -> SplitList`

Types:

Subject = iodata() | charlist()

RE = mp() | iodata() | charlist()

Options = [Option]

Option = anchored | global | notbol | noteol | notempty | {offset, int()} | {newline, NLSpec} | bsr_anycrlf | bsr_unicode | {return, ReturnType} | {parts, NumParts} | group | trim | CompileOpt

NumParts = int() | infinity

ReturnType = iodata | list | binary

CompileOpt = see compile/2 above

NLSpec = cr | crlf | lf | anycrlf | any

SplitList = [RetData] | [GroupedRetData]

GroupedRetData = [RetData]

RetData = iodata() | charlist() | binary() | list()

This function splits the input into parts by finding tokens according to the regular expression supplied.

The splitting is done basically by running a global regexp match and dividing the initial string wherever a match occurs. The matching part of the string is removed from the output.

As in the `re:run/3` function, an `mp()` compiled with the `unicode` option requires the `Subject` to be a Unicode `charlist()`. If compilation is done implicitly and the `unicode` compilation option is given to this function, both the regular expression and the `Subject` should be given as valid Unicode `charlist()`s.

The result is given as a list of "strings", the preferred datatype given in the `return` option (default iodata).

If subexpressions are given in the regular expression, the matching subexpressions are returned in the resulting list as well. An example:

```
re:split("Erlang","[ln]",[{return,list}]).
```

will yield the result:

```
["Er","a","g"]
```

while

```
re:split("Erlang","([ln])",[{return,list}]).
```

will yield

```
["Er","l","a","n","g"]
```

The text matching the subexpression (marked by the parentheses in the regexp) is inserted in the result list where it was found. In effect this means that concatenating the result of a split where the whole regexp is a single subexpression (as in the example above) will always result in the original string.

As there is no matching subexpression for the last part in the example (the "g"), there is nothing inserted after that. To make the group of strings and the parts matching the subexpressions more obvious, one might use the `group` option, which groups together the part of the subject string with the parts matching the subexpressions when the string was split:

```
re:split("Erlang","([ln])",[{return,list},group]).
```

gives:

```
[["Er","l"],["a","n"],["g"]]
```

Here the regular expression matched first the "l", causing "Er" to be the first part in the result. When the regular expression matched, the (only) subexpression was bound to the "l", so the "l" is inserted in the group together with "Er". The next match is of the "n", making "a" the next part to be returned. Since the subexpression is bound to the substring "n" in this case, the "n" is inserted into this group. The last group consists of the rest of the string, as no more matches are found.

By default, all parts of the string, including the empty strings, are returned from the function. For example:

```
re:split("Erlang","[lg]",[{return,list}]).
```

will return:

```
["Er","an",[]]
```

since the matching of the "g" in the end of the string leaves an empty rest which is also returned. This behaviour differs from the default behaviour of the split function in Perl, where empty strings at the end are by default removed. To get the "trimming" default behavior of Perl, specify `trim` as an option:

```
re:split("Erlang","[lg]",[{return,list},trim]).
```

The result will be:

```
["Er","an"]
```

The "trim" option in effect says; "give me as many parts as possible except the empty ones", which might be useful in some circumstances. You can also specify how many parts you want, by specifying `{parts,`N`}`:

```
re:split("Erlang","[lg]",[{return,list},{parts,2}]).
```

This will give:

```
["Er","ang"]
```

Note that the last part is "ang", not "an", as we only specified splitting into two parts, and the splitting stops when enough parts are given, which is why the result differs from that of `trim`.

More than three parts are not possible with this indata, so

```
re:split("Erlang","[lg]",[{return,list},{parts,4}]).
```

will give the same result as the default, which is to be viewed as "an infinite number of parts".

Specifying `0` as the number of parts gives the same effect as the option `trim`. If subexpressions are captured, empty subexpression matches at the end are also stripped from the result if `trim` or `{parts,0}` is specified.

If you are familiar with Perl, the `trim` behaviour corresponds exactly to the Perl default, the `{parts,N}` where N is a positive integer corresponds exactly to the Perl behaviour with a positive numerical third parameter and the default behaviour of `re:split/3` corresponds to that when the Perl routine is given a negative integer as the third parameter.

Summary of options not previously described for the `re:run/3` function:

{return,ReturnType}

Specifies how the parts of the original string are presented in the result list. The possible types are:

iodata
The variant of `iodata()` that gives the least copying of data with the current implementation (often a binary, but don't depend on it).
binary
All parts returned as binaries.
list
All parts returned as lists of characters ("strings").

group

Groups together the part of the string with the parts of the string matching the subexpressions of the regexp.

The return value from the function will in this case be a `list()` of `list()`s. Each sublist begins with the string picked out of the subject string, followed by the parts matching each of the subexpressions in order of occurrence in the regular expression.

{parts,N}

Specifies the number of parts the subject string is to be split into.

The number of parts should be a positive integer for a specific maximum on the number of parts and `infinity` for the maximum number of parts possible (the default). Specifying `{parts,0}` gives as many parts as possible disregarding empty parts at the end, the same as specifying `trim`

trim

Specifies that empty parts at the end of the result list are to be disregarded. The same as specifying `{parts,0}`. This corresponds to the default behaviour of the `split` built in function in Perl.

PERL LIKE REGULAR EXPRESSIONS SYNTAX

The following sections contain reference material for the regular expressions used by this module. The regular expression reference is based on the PCRE documentation, with changes in cases where the re module behaves differently to the PCRE library.

PCRE regular expression details

The syntax and semantics of the regular expressions that are supported by PCRE are described in detail below. Perl's regular expressions are described in its own documentation, and regular expressions in general are covered in a number of books, some of which have copious examples. Jeffrey Friedl's "Mastering Regular Expressions", published by O'Reilly, covers regular expressions in great detail. This description of PCRE's regular expressions is intended as reference material.

The reference material is divided into the following sections:

- *Newline conventions*
- *Characters and metacharacters*
- *Backslash*
- *Circumflex and dollar*
- *Full stop (period, dot)*

- *Matching a single byte*
- *Square brackets and character classes*
- *POSIX character classes*
- *Vertical bar*
- *Internal option setting*
- *Subpatterns*
- *Duplicate subpattern numbers*
- *Named subpatterns*
- *Repetition*
- *Atomic grouping and possessive quantifiers*
- *Back references*
- *Assertions*
- *Conditional subpatterns*
- *Comments*
- *Recursive patterns*
- *Subpatterns as subroutines*
- *Backtracking control*

Newline conventions

PCRE supports five different conventions for indicating line breaks in strings: a single CR (carriage return) character, a single LF (linefeed) character, the two-character sequence CRLF , any of the three preceding, or any Unicode newline sequence.

It is also possible to specify a newline convention by starting a pattern string with one of the following five sequences:

(*CR)
: carriage return

(*LF)
: linefeed

(*CRLF)
: carriage return, followed by linefeed

(*ANYCRLF)
: any of the three above

(*ANY)
: all Unicode newline sequences

These override the default and the options given to `re:compile/2`. For example, the pattern:

(*CR)a.b

changes the convention to CR. That pattern matches "a\nb" because LF is no longer a newline. Note that these special settings, which are not Perl-compatible, are recognized only at the very start of a pattern, and that they must be in upper case. If more than one of them is present, the last one is used.

The newline convention does not affect what the \R escape sequence matches. By default, this is any Unicode newline sequence, for Perl compatibility. However, this can be changed; see the description of \R in the section entitled "Newline sequences" below. A change of \R setting can be combined with a change of newline convention.

Characters and metacharacters

A regular expression is a pattern that is matched against a subject string from left to right. Most characters stand for themselves in a pattern, and match the corresponding characters in the subject. As a trivial example, the pattern

The quick brown fox

matches a portion of a subject string that is identical to itself. When caseless matching is specified (the `caseless` option), letters are matched independently of case.

The power of regular expressions comes from the ability to include alternatives and repetitions in the pattern. These are encoded in the pattern by the use of *metacharacters*, which do not stand for themselves but instead are interpreted in some special way.

There are two different sets of metacharacters: those that are recognized anywhere in the pattern except within square brackets, and those that are recognized within square brackets. Outside square brackets, the metacharacters are as follows:

\
: general escape character with several uses

^
: assert start of string (or line, in multiline mode)

$
: assert end of string (or line, in multiline mode)

.
: match any character except newline (by default)

[
: start character class definition

|
: start of alternative branch

(
: start subpattern

)
: end subpattern

?
: extends the meaning of (, also 0 or 1 quantifier, also quantifier minimizer

*
: 0 or more quantifier

+
: 1 or more quantifier, also "possessive quantifier"

{
: start min/max quantifier

Part of a pattern that is in square brackets is called a "character class". In a character class the only metacharacters are:

\
: general escape character

^
: negate the class, but only if the first character

-
: indicates character range

[
: POSIX character class (only if followed by POSIX syntax)

]
: terminates the character class

The following sections describe the use of each of the metacharacters.

Backslash

The backslash character has several uses. Firstly, if it is followed by a non-alphanumeric character, it takes away any special meaning that character may have. This use of backslash as an escape character applies both inside and outside character classes.

For example, if you want to match a * character, you write * in the pattern. This escaping action applies whether or not the following character would otherwise be interpreted as a metacharacter, so it is always safe to precede a non-alphanumeric with backslash to specify that it stands for itself. In particular, if you want to match a backslash, you write \\.

If a pattern is compiled with the `extended` option, whitespace in the pattern (other than in a character class) and characters between a # outside a character class and the next newline are ignored. An escaping backslash can be used to include a whitespace or # character as part of the pattern.

If you want to remove the special meaning from a sequence of characters, you can do so by putting them between \Q and \E. This is different from Perl in that $ and @ are handled as literals in \Q...\E sequences in PCRE, whereas in Perl, $ and @ cause variable interpolation. Note the following examples:

```
Pattern           PCRE matches   Perl matches

\Qabc$xyz\E       abc$xyz        abc followed by the contents of $xyz
\Qabc\$xyz\E      abc\$xyz       abc\$xyz
\Qabc\E\$\Qxyz\E  abc$xyz        abc$xyz
```

The \Q...\E sequence is recognized both inside and outside character classes.

Non-printing characters

A second use of backslash provides a way of encoding non-printing characters in patterns in a visible manner. There is no restriction on the appearance of non-printing characters, apart from the binary zero that terminates a pattern, but when a pattern is being prepared by text editing, it is usually easier to use one of the following escape sequences than the binary character it represents:

\a
: alarm, that is, the BEL character (hex 07)

\cx
: "control-x", where x is any character

\e
: escape (hex 1B)

\f
: formfeed (hex 0C)

\n
: linefeed (hex 0A)

\r
: carriage return (hex 0D)

\t
: tab (hex 09)

\ddd
: character with octal code ddd, or backreference

\xhh
: character with hex code hh

\x{hhh..}
: character with hex code hhh..

The precise effect of \cx is as follows: if x is a lower case letter, it is converted to upper case. Then bit 6 of the character (hex 40) is inverted. Thus \cz becomes hex 1A, but \c{ becomes hex 3B, while \c; becomes hex 7B.

After \x, from zero to two hexadecimal digits are read (letters can be in upper or lower case). Any number of hexadecimal digits may appear between \x{ and }, but the value of the character code must be less than 256 in non-UTF-8 mode, and less than 2**31 in UTF-8 mode. That is, the maximum value in hexadecimal is 7FFFFFFF. Note that this is bigger than the largest Unicode code point, which is 10FFFF.

If characters other than hexadecimal digits appear between \x{ and }, or if there is no terminating }, this form of escape is not recognized. Instead, the initial \x will be interpreted as a basic hexadecimal escape, with no following digits, giving a character whose value is zero.

Characters whose value is less than 256 can be defined by either of the two syntaxes for \x. There is no difference in the way they are handled. For example, \xdc is exactly the same as \x{dc}.

After \0 up to two further octal digits are read. If there are fewer than two digits, just those that are present are used. Thus the sequence \0\x\07 specifies two binary zeros followed by a BEL character (code value 7). Make sure you supply two digits after the initial zero if the pattern character that follows is itself an octal digit.

The handling of a backslash followed by a digit other than 0 is complicated. Outside a character class, PCRE reads it and any following digits as a decimal number. If the number is less than 10, or if there have been at least that many previous capturing left parentheses in the expression, the entire sequence is taken as a *back reference*. A description of how this works is given later, following the discussion of parenthesized subpatterns.

Inside a character class, or if the decimal number is greater than 9 and there have not been that many capturing subpatterns, PCRE re-reads up to three octal digits following the backslash, and uses them to generate a data character. Any subsequent digits stand for themselves. The value of a character specified in octal must be less than \400. In non-UTF-8 mode, the value of a character specified in octal must be less than \400. In UTF-8 mode, values up to \777 are permitted. For example:

\040
: is another way of writing a space

\40
: is the same, provided there are fewer than 40 previous capturing subpatterns

\7
: is always a back reference

\11
: might be a back reference, or another way of writing a tab

\011
: is always a tab

\0113
: is a tab followed by the character "3"

\113
: might be a back reference, otherwise the character with octal code 113

\377
: might be a back reference, otherwise the byte consisting entirely of 1 bits

\81
: is either a back reference, or a binary zero followed by the two characters "8" and "1"

Note that octal values of 100 or greater must not be introduced by a leading zero, because no more than three octal digits are ever read.

All the sequences that define a single character value can be used both inside and outside character classes. In addition, inside a character class, the sequence \b is interpreted as the backspace character (hex 08), and the sequences \R and \X are interpreted as the characters "R" and "X", respectively. Outside a character class, these sequences have different meanings (see below).

Absolute and relative back references

The sequence \g followed by an unsigned or a negative number, optionally enclosed in braces, is an absolute or relative back reference. A named back reference can be coded as \g{name}. Back references are discussed later, following the discussion of parenthesized subpatterns.

Generic character types

Another use of backslash is for specifying generic character types. The following are always recognized:

\d
: any decimal digit

\D
: any character that is not a decimal digit

\h
: any horizontal whitespace character

\H
: any character that is not a horizontal whitespace character

\s
: any whitespace character

\S
: any character that is not a whitespace character

\v
: any vertical whitespace character

\V
: any character that is not a vertical whitespace character

\w
: any "word" character

\W
: any "non-word" character

Each pair of escape sequences partitions the complete set of characters into two disjoint sets. Any given character matches one, and only one, of each pair.

These character type sequences can appear both inside and outside character classes. They each match one character of the appropriate type. If the current matching point is at the end of the subject string, all of them fail, since there is no character to match.

For compatibility with Perl, \s does not match the VT character (code 11). This makes it different from the POSIX "space" class. The \s characters are HT (9), LF (10), FF (12), CR (13), and space (32). If "use locale;" is included in a Perl script, \s may match the VT character. In PCRE, it never does.

In UTF-8 mode, characters with values greater than 128 never match \d, \s, or \w, and always match \D, \S, and \W. This is true even when Unicode character property support is available. These sequences retain their original meanings from before UTF-8 support was available, mainly for efficiency reasons.

The sequences \h, \H, \v, and \V are Perl 5.10 features. In contrast to the other sequences, these do match certain high-valued codepoints in UTF-8 mode. The horizontal space characters are:

U+0009
: Horizontal tab

U+0020
: Space

U+00A0
: Non-break space

U+1680
: Ogham space mark

U+180E
: Mongolian vowel separator

U+2000
: En quad

U+2001
: Em quad

U+2002
: En space

U+2003
: Em space

U+2004
: Three-per-em space

U+2005
: Four-per-em space

U+2006
: Six-per-em space

U+2007
: Figure space

U+2008
: Punctuation space

U+2009
: Thin space

U+200A
: Hair space

U+202F
: Narrow no-break space

U+205F
: Medium mathematical space

U+3000
: Ideographic space

The vertical space characters are:

U+000A
: Linefeed

U+000B
: Vertical tab

U+000C
: Formfeed

U+000D
: Carriage return

U+0085
: Next line

U+2028
: Line separator

U+2029
: Paragraph separator

A "word" character is an underscore or any character less than 256 that is a letter or digit. The definition of letters and digits is controlled by PCRE's low-valued character tables, which are always ISO-8859-1.

Newline sequences

Outside a character class, by default, the escape sequence \R matches any Unicode newline sequence. This is a Perl 5.10 feature. In non-UTF-8 mode \R is equivalent to the following:

(?>\r\n|\n|\x0b|\f|\r|\x85)

This is an example of an "atomic group", details of which are given below.

This particular group matches either the two-character sequence CR followed by LF, or one of the single characters LF (linefeed, U+000A), VT (vertical tab, U+000B), FF (formfeed, U+000C), CR (carriage return, U+000D), or NEL (next line, U+0085). The two-character sequence is treated as a single unit that cannot be split.

In UTF-8 mode, two additional characters whose codepoints are greater than 255 are added: LS (line separator, U +2028) and PS (paragraph separator, U+2029). Unicode character property support is not needed for these characters to be recognized.

It is possible to restrict \R to match only CR, LF, or CRLF (instead of the complete set of Unicode line endings) by setting the option `bsr_anycrlf` either at compile time or when the pattern is matched. (BSR is an abbreviation for "backslash R".) This can be made the default when PCRE is built; if this is the case, the other behaviour can be requested via the `bsr_unicode` option. It is also possible to specify these settings by starting a pattern string with one of the following sequences:

(*BSR_ANYCRLF) CR, LF, or CRLF only (*BSR_UNICODE) any Unicode newline sequence

These override the default and the options given to `re:compile/2`, but they can be overridden by options given to `re:run/3`. Note that these special settings, which are not Perl-compatible, are recognized only at the very start of a pattern, and that they must be in upper case. If more than one of them is present, the last one is used. They can be combined with a change of newline convention, for example, a pattern can start with:

(*ANY)(*BSR_ANYCRLF)

Inside a character class, \R matches the letter "R".

Unicode character properties

When PCRE is built with Unicode character property support, three additional escape sequences that match characters with specific properties are available. When not in UTF-8 mode, these sequences are of course limited to testing characters whose codepoints are less than 256, but they do work in this mode. The extra escape sequences are:

\p{*xx*} a character with the *xx* property \P{*xx*} a character without the *xx* property \X an extended Unicode sequence

The property names represented by *xx* above are limited to the Unicode script names, the general category properties, and "Any", which matches any character (including newline). Other properties such as "InMusicalSymbols" are not currently supported by PCRE. Note that \P{Any} does not match any characters, so always causes a match failure.

Sets of Unicode characters are defined as belonging to certain scripts. A character from one of these sets can be matched using a script name. For example:

\p{Greek} \P{Han}

Those that are not part of an identified script are lumped together as "Common". The current list of scripts is:

- Arabic
- Armenian
- Balinese
- Bengali
- Bopomofo
- Braille
- Buginese
- Buhid
- Canadian_Aboriginal
- Cherokee
- Common
- Coptic

- Cuneiform
- Cypriot
- Cyrillic
- Deseret
- Devanagari
- Ethiopic
- Georgian
- Glagolitic
- Gothic
- Greek
- Gujarati
- Gurmukhi
- Han
- Hangul
- Hanunoo
- Hebrew
- Hiragana
- Inherited
- Kannada
- Katakana
- Kharoshthi
- Khmer
- Lao
- Latin
- Limbu
- Linear_B
- Malayalam
- Mongolian
- Myanmar
- New_Tai_Lue
- Nko
- Ogham
- Old_Italic
- Old_Persian
- Oriya
- Osmanya
- Phags_Pa
- Phoenician
- Runic
- Shavian
- Sinhala
- Syloti_Nagri
- Syriac

- Tagalog
- Tagbanwa
- Tai_Le
- Tamil
- Telugu
- Thaana
- Thai
- Tibetan
- Tifinagh
- Ugaritic
- Yi

Each character has exactly one general category property, specified by a two-letter abbreviation. For compatibility with Perl, negation can be specified by including a circumflex between the opening brace and the property name. For example, \p{^Lu} is the same as \P{Lu}.

If only one letter is specified with \p or \P, it includes all the general category properties that start with that letter. In this case, in the absence of negation, the curly brackets in the escape sequence are optional; these two examples have the same effect:

- \p{L}
- \pL

The following general category property codes are supported:

C
: Other

Cc
: Control

Cf
: Format

Cn
: Unassigned

Co
: Private use

Cs
: Surrogate

L
: Letter

Ll
: Lower case letter

Lm
: Modifier letter

Lo
: Other letter

Lt
: Title case letter

Lu
: Upper case letter

M
: Mark

Mc
: Spacing mark

Me
: Enclosing mark

Mn
: Non-spacing mark

N
: Number

Nd
: Decimal number

Nl
: Letter number

No
: Other number

P
: Punctuation

Pc
: Connector punctuation

Pd
: Dash punctuation

Pe
: Close punctuation

Pf
: Final punctuation

Pi
: Initial punctuation

Po
: Other punctuation

Ps
: Open punctuation

S
: Symbol

Sc
: Currency symbol

Sk
: Modifier symbol

Sm
: Mathematical symbol

So
: Other symbol

Z
: Separator

Zl
: Line separator

Zp
: Paragraph separator

Zs
: Space separator

The special property L& is also supported: it matches a character that has the Lu, Ll, or Lt property, in other words, a letter that is not classified as a modifier or "other".

The Cs (Surrogate) property applies only to characters in the range U+D800 to U+DFFF. Such characters are not valid in UTF-8 strings (see RFC 3629) and so cannot be tested by PCRE, unless UTF-8 validity checking has been turned off (see the discussion of `no_utf8_check` in the *pcreapi* page).

The long synonyms for these properties that Perl supports (such as \p{Letter}) are not supported by PCRE, nor is it permitted to prefix any of these properties with "Is".

No character that is in the Unicode table has the Cn (unassigned) property. Instead, this property is assumed for any code point that is not in the Unicode table.

Specifying caseless matching does not affect these escape sequences. For example, \p{Lu} always matches only upper case letters.

The \X escape matches any number of Unicode characters that form an extended Unicode sequence. \X is equivalent to

(?>\PM\pM*)

That is, it matches a character without the "mark" property, followed by zero or more characters with the "mark" property, and treats the sequence as an atomic group (see below). Characters with the "mark" property are typically accents that affect the preceding character. None of them have codepoints less than 256, so in non-UTF-8 mode \X matches any one character.

Matching characters by Unicode property is not fast, because PCRE has to search a structure that contains data for over fifteen thousand characters. That is why the traditional escape sequences such as \d and \w do not use Unicode properties in PCRE.

Resetting the match start

The escape sequence \K, which is a Perl 5.10 feature, causes any previously matched characters not to be included in the final matched sequence. For example, the pattern:

foo\Kbar

matches "foobar", but reports that it has matched "bar". This feature is similar to a lookbehind assertion (described below). However, in this case, the part of the subject before the real match does not have to be of fixed length, as lookbehind assertions do. The use of \K does not interfere with the setting of captured substrings. For example, when the pattern

(foo)\Kbar

matches "foobar", the first substring is still set to "foo".

Simple assertions

The final use of backslash is for certain simple assertions. An assertion specifies a condition that has to be met at a particular point in a match, without consuming any characters from the subject string. The use of subpatterns for more complicated assertions is described below. The backslashed assertions are:

\b
: matches at a word boundary

\B
: matches when not at a word boundary

\A
: matches at the start of the subject

\Z
: matches at the end of the subject also matches before a newline at the end of the subject

\z
: matches only at the end of the subject

\G
: matches at the first matching position in the subject

These assertions may not appear in character classes (but note that \b has a different meaning, namely the backspace character, inside a character class).

A word boundary is a position in the subject string where the current character and the previous character do not both match \w or \W (i.e. one matches \w and the other matches \W), or the start or end of the string if the first or last character matches \w, respectively.

The \A, \Z, and \z assertions differ from the traditional circumflex and dollar (described in the next section) in that they only ever match at the very start and end of the subject string, whatever options are set. Thus, they are independent of multiline mode. These three assertions are not affected by the `notbol` or `noteol` options, which affect only the behaviour of the circumflex and dollar metacharacters. However, if the *startoffset* argument of `re:run/3` is non-zero, indicating that matching is to start at a point other than the beginning of the subject, \A can never match. The difference between \Z and \z is that \Z matches before a newline at the end of the string as well as at the very end, whereas \z matches only at the end.

The \G assertion is true only when the current matching position is at the start point of the match, as specified by the *startoffset* argument of `re:run/3`. It differs from \A when the value of *startoffset* is non-zero. By calling `re:run/3` multiple times with appropriate arguments, you can mimic Perl's /g option, and it is in this kind of implementation where \G can be useful.

Note, however, that PCRE's interpretation of \G, as the start of the current match, is subtly different from Perl's, which defines it as the end of the previous match. In Perl, these can be different when the previously matched string was empty. Because PCRE does just one match at a time, it cannot reproduce this behaviour.

If all the alternatives of a pattern begin with \G, the expression is anchored to the starting match position, and the "anchored" flag is set in the compiled regular expression.

Circumflex and dollar

Outside a character class, in the default matching mode, the circumflex character is an assertion that is true only if the current matching point is at the start of the subject string. If the *startoffset* argument of `re:run/3` is non-zero, circumflex can never match if the `multiline` option is unset. Inside a character class, circumflex has an entirely different meaning (see below).

Circumflex need not be the first character of the pattern if a number of alternatives are involved, but it should be the first thing in each alternative in which it appears if the pattern is ever to match that branch. If all possible alternatives start with a circumflex, that is, if the pattern is constrained to match only at the start of the subject, it is said to be an "anchored" pattern. (There are also other constructs that can cause a pattern to be anchored.)

A dollar character is an assertion that is true only if the current matching point is at the end of the subject string, or immediately before a newline at the end of the string (by default). Dollar need not be the last character of the pattern if a number of alternatives are involved, but it should be the last item in any branch in which it appears. Dollar has no special meaning in a character class.

The meaning of dollar can be changed so that it matches only at the very end of the string, by setting the `dollar_endonly` option at compile time. This does not affect the \Z assertion.

The meanings of the circumflex and dollar characters are changed if the `multiline` option is set. When this is the case, a circumflex matches immediately after internal newlines as well as at the start of the subject string. It does not match after a newline that ends the string. A dollar matches before any newlines in the string, as well as at the very end, when `multiline` is set. When newline is specified as the two-character sequence CRLF, isolated CR and LF characters do not indicate newlines.

For example, the pattern /^abc$/ matches the subject string "def\nabc" (where \n represents a newline) in multiline mode, but not otherwise. Consequently, patterns that are anchored in single line mode because all branches start with ^ are not anchored in multiline mode, and a match for circumflex is possible when the *startoffset* argument of `re:run/3` is non-zero. The `dollar_endonly` option is ignored if `multiline` is set.

Note that the sequences \A, \Z, and \z can be used to match the start and end of the subject in both modes, and if all branches of a pattern start with \A it is always anchored, whether or not `multiline` is set.

Full stop (period, dot)

Outside a character class, a dot in the pattern matches any one character in the subject string except (by default) a character that signifies the end of a line. In UTF-8 mode, the matched character may be more than one byte long.

When a line ending is defined as a single character, dot never matches that character; when the two-character sequence CRLF is used, dot does not match CR if it is immediately followed by LF, but otherwise it matches all characters (including isolated CRs and LFs). When any Unicode line endings are being recognized, dot does not match CR or LF or any of the other line ending characters.

The behaviour of dot with regard to newlines can be changed. If the `dotall` option is set, a dot matches any one character, without exception. If the two-character sequence CRLF is present in the subject string, it takes two dots to match it.

The handling of dot is entirely independent of the handling of circumflex and dollar, the only relationship being that they both involve newlines. Dot has no special meaning in a character class.

Matching a single byte

Outside a character class, the escape sequence \C matches any one byte, both in and out of UTF-8 mode. Unlike a dot, it always matches any line-ending characters. The feature is provided in Perl in order to match individual bytes in UTF-8 mode. Because it breaks up UTF-8 characters into individual bytes, what remains in the string may be a malformed UTF-8 string. For this reason, the \C escape sequence is best avoided.

PCRE does not allow \C to appear in lookbehind assertions (described below), because in UTF-8 mode this would make it impossible to calculate the length of the lookbehind.

Square brackets and character classes

An opening square bracket introduces a character class, terminated by a closing square bracket. A closing square bracket on its own is not special. If a closing square bracket is required as a member of the class, it should be the first data character in the class (after an initial circumflex, if present) or escaped with a backslash.

A character class matches a single character in the subject. In UTF-8 mode, the character may occupy more than one byte. A matched character must be in the set of characters defined by the class, unless the first character in the class definition is a circumflex, in which case the subject character must not be in the set defined by the class. If a circumflex is actually required as a member of the class, ensure it is not the first character, or escape it with a backslash.

For example, the character class [aeiou] matches any lower case vowel, while [^aeiou] matches any character that is not a lower case vowel. Note that a circumflex is just a convenient notation for specifying the characters that are in the class by enumerating those that are not. A class that starts with a circumflex is not an assertion: it still consumes a character from the subject string, and therefore it fails if the current pointer is at the end of the string.

In UTF-8 mode, characters with values greater than 255 can be included in a class as a literal string of bytes, or by using the \x{ escaping mechanism.

When caseless matching is set, any letters in a class represent both their upper case and lower case versions, so for example, a caseless [aeiou] matches "A" as well as "a", and a caseless [^aeiou] does not match "A", whereas a caseful version would. In UTF-8 mode, PCRE always understands the concept of case for characters whose values are less than 128, so caseless matching is always possible. For characters with higher values, the concept of case is supported if PCRE is compiled with Unicode property support, but not otherwise. If you want to use caseless matching for characters 128 and above, you must ensure that PCRE is compiled with Unicode property support as well as with UTF-8 support.

Characters that might indicate line breaks are never treated in any special way when matching character classes, whatever line-ending sequence is in use, and whatever setting of the `dotall` and `multiline` options is used. A class such as [^a] always matches one of these characters.

The minus (hyphen) character can be used to specify a range of characters in a character class. For example, [d-m] matches any letter between d and m, inclusive. If a minus character is required in a class, it must be escaped with a backslash or appear in a position where it cannot be interpreted as indicating a range, typically as the first or last character in the class.

It is not possible to have the literal character "]" as the end character of a range. A pattern such as [W-]46] is interpreted as a class of two characters ("W" and "-") followed by a literal string "46]", so it would match "W46]" or "-46]". However, if the "]" is escaped with a backslash it is interpreted as the end of range, so [W-\]46] is interpreted as a class containing a range followed by two other characters. The octal or hexadecimal representation of "]" can also be used to end a range.

Ranges operate in the collating sequence of character values. They can also be used for characters specified numerically, for example [\000-\037]. In UTF-8 mode, ranges can include characters whose values are greater than 255, for example [\x{100}-\x{2ff}].

If a range that includes letters is used when caseless matching is set, it matches the letters in either case. For example, [W-c] is equivalent to [][\\^_`wxyzabc], matched caselessly , and in non-UTF-8 mode, if character tables for a French locale are in use, [\xc8-\xcb] matches accented E characters in both cases. In UTF-8 mode, PCRE supports the concept of case for characters with values greater than 128 only when it is compiled with Unicode property support.

The character types \d, \D, \p, \P, \s, \S, \w, and \W may also appear in a character class, and add the characters that they match to the class. For example, [\dABCDEF] matches any hexadecimal digit. A circumflex can conveniently be used with the upper case character types to specify a more restricted set of characters than the matching lower case type. For example, the class [^\W_] matches any letter or digit, but not underscore.

The only metacharacters that are recognized in character classes are backslash, hyphen (only where it can be interpreted as specifying a range), circumflex (only at the start), opening square bracket (only when it can be interpreted as introducing a POSIX class name - see the next section), and the terminating closing square bracket. However, escaping other non-alphanumeric characters does no harm.

POSIX character classes

Perl supports the POSIX notation for character classes. This uses names enclosed by [: and :] within the enclosing square brackets. PCRE also supports this notation. For example,

[01[:alpha:]%]

matches "0", "1", any alphabetic character, or "%". The supported class names are

alnum
: letters and digits

alpha
: letters

ascii
: character codes 0 - 127

blank
: space or tab only

cntrl
: control characters

digit
: decimal digits (same as \d)

graph
: printing characters, excluding space

lower
: lower case letters

print
: printing characters, including space

punct
: printing characters, excluding letters and digits

space
: whitespace (not quite the same as \s)

upper
: upper case letters

word
: "word" characters (same as \w)

xdigit
: hexadecimal digits

The "space" characters are HT (9), LF (10), VT (11), FF (12), CR (13), and space (32). Notice that this list includes the VT character (code 11). This makes "space" different to \s, which does not include VT (for Perl compatibility).

The name "word" is a Perl extension, and "blank" is a GNU extension from Perl 5.8. Another Perl extension is negation, which is indicated by a ^ character after the colon. For example,

[12[:^digit:]]

matches "1", "2", or any non-digit. PCRE (and Perl) also recognize the POSIX syntax [.ch.] and [=ch=] where "ch" is a "collating element", but these are not supported, and an error is given if they are encountered.

In UTF-8 mode, characters with values greater than 128 do not match any of the POSIX character classes.

Vertical bar

Vertical bar characters are used to separate alternative patterns. For example, the pattern

gilbert|sullivan

matches either "gilbert" or "sullivan". Any number of alternatives may appear, and an empty alternative is permitted (matching the empty string). The matching process tries each alternative in turn, from left to right, and the first one that succeeds is used. If the alternatives are within a subpattern (defined below), "succeeds" means matching the rest of the main pattern as well as the alternative in the subpattern.

Internal option setting

The settings of the `caseless`, `multiline`, `dotall`, and `extended` options (which are Perl-compatible) can be changed from within the pattern by a sequence of Perl option letters enclosed between "(?" and ")". The option letters are

i
: for `caseless`

m
: for `multiline`

s
: for `dotall`

x
: for `extended`

For example, (?im) sets caseless, multiline matching. It is also possible to unset these options by preceding the letter with a hyphen, and a combined setting and unsetting such as (?im-sx), which sets `caseless` and `multiline` while unsetting `dotall` and `extended`, is also permitted. If a letter appears both before and after the hyphen, the option is unset.

The PCRE-specific options `dupnames`, `ungreedy`, and `extra` can be changed in the same way as the Perl-compatible options by using the characters J, U and X respectively.

When an option change occurs at top level (that is, not inside subpattern parentheses), the change applies to the remainder of the pattern that follows. If the change is placed right at the start of a pattern, PCRE extracts it into the global options

An option change within a subpattern (see below for a description of subpatterns) affects only that part of the current pattern that follows it, so

(a(?i)b)c

matches abc and aBc and no other strings (assuming `caseless` is not used). By this means, options can be made to have different settings in different parts of the pattern. Any changes made in one alternative do carry on into subsequent branches within the same subpattern. For example,

(a(?i)b|c)

matches "ab", "aB", "c", and "C", even though when matching "C" the first branch is abandoned before the option setting. This is because the effects of option settings happen at compile time. There would be some very weird behaviour otherwise.

Note: There are other PCRE-specific options that can be set by the application when the compile or match functions are called. In some cases the pattern can contain special leading sequences to override what the application has set or what has been defaulted. Details are given in the section entitled "Newline sequences" above.

Subpatterns

Subpatterns are delimited by parentheses (round brackets), which can be nested. Turning part of a pattern into a subpattern does two things:

1. It localizes a set of alternatives. For example, the pattern

cat(aract|erpillar|)

matches one of the words "cat", "cataract", or "caterpillar". Without the parentheses, it would match "cataract", "erpillar" or an empty string.

2. It sets up the subpattern as a capturing subpattern. This means that, when the complete pattern matches, that portion of the subject string that matched the subpattern is passed back to the caller via the return value of `re:run/3`. Opening parentheses are counted from left to right (starting from 1) to obtain numbers for the capturing subpatterns.

For example, if the string "the red king" is matched against the pattern

the ((red|white) (king|queen))

the captured substrings are "red king", "red", and "king", and are numbered 1, 2, and 3, respectively.

The fact that plain parentheses fulfil two functions is not always helpful. There are often times when a grouping subpattern is required without a capturing requirement. If an opening parenthesis is followed by a question mark and a colon, the subpattern does not do any capturing, and is not counted when computing the number of any subsequent capturing subpatterns. For example, if the string "the white queen" is matched against the pattern

the ((?:red|white) (king|queen))

the captured substrings are "white queen" and "queen", and are numbered 1 and 2. The maximum number of capturing subpatterns is 65535.

As a convenient shorthand, if any option settings are required at the start of a non-capturing subpattern, the option letters may appear between the "?" and the ":". Thus the two patterns

- (?i:saturday|sunday)
- (?:(?i)saturday|sunday)

match exactly the same set of strings. Because alternative branches are tried from left to right, and options are not reset until the end of the subpattern is reached, an option setting in one branch does affect subsequent branches, so the above patterns match "SUNDAY" as well as "Saturday".

Duplicate subpattern numbers

Perl 5.10 introduced a feature whereby each alternative in a subpattern uses the same numbers for its capturing parentheses. Such a subpattern starts with (?| and is itself a non-capturing subpattern. For example, consider this pattern:

(?|(Sat)ur|(Sun))day

Because the two alternatives are inside a (?| group, both sets of capturing parentheses are numbered one. Thus, when the pattern matches, you can look at captured substring number one, whichever alternative matched. This construct is useful when you want to capture part, but not all, of one of a number of alternatives. Inside a (?| group, parentheses are numbered as usual, but the number is reset at the start of each branch. The numbers of any capturing buffers that follow the subpattern start after the highest number used in any branch. The following example is taken from the Perl documentation. The numbers underneath show in which buffer the captured content will be stored.

```
# before  ---------------branch-reset----------- after
/ ( a )  (?| x ( y ) z | (p (q) r) | (t) u (v) ) ( z ) /x
# 1           2         2  3        2     3     4
```

A backreference or a recursive call to a numbered subpattern always refers to the first one in the pattern with the given number.

An alternative approach to using this "branch reset" feature is to use duplicate named subpatterns, as described in the next section.

Named subpatterns

Identifying capturing parentheses by number is simple, but it can be very hard to keep track of the numbers in complicated regular expressions. Furthermore, if an expression is modified, the numbers may change. To help with this difficulty, PCRE supports the naming of subpatterns. This feature was not added to Perl until release 5.10. Python had the feature earlier, and PCRE introduced it at release 4.0, using the Python syntax. PCRE now supports both the Perl and the Python syntax.

In PCRE, a subpattern can be named in one of three ways: (?<name>...) or (?'name'...) as in Perl, or (?P<name>...) as in Python. References to capturing parentheses from other parts of the pattern, such as backreferences, recursion, and conditions, can be made by name as well as by number.

Names consist of up to 32 alphanumeric characters and underscores. Named capturing parentheses are still allocated numbers as well as names, exactly as if the names were not present. The `capture` specification to `re:run/3` can use named values if they are present in the regular expression.

By default, a name must be unique within a pattern, but it is possible to relax this constraint by setting the `dupnames` option at compile time. This can be useful for patterns where only one instance of the named parentheses can match. Suppose you want to match the name of a weekday, either as a 3-letter abbreviation or as the full name, and in both cases you want to extract the abbreviation. This pattern (ignoring the line breaks) does the job:

```
(?<DN>Mon|Fri|Sun)(?:day)?|
(?<DN>Tue)(?:sday)?|
(?<DN>Wed)(?:nesday)?|
(?<DN>Thu)(?:rsday)?|
(?<DN>Sat)(?:urday)?
```

There are five capturing substrings, but only one is ever set after a match. (An alternative way of solving this problem is to use a "branch reset" subpattern, as described in the previous section.)

In case of capturing named subpatterns which are not unique, the first occurrence is returned from `re:exec/3`, if the name is specified int the `values` part of the `capture` statement.

Repetition

Repetition is specified by quantifiers, which can follow any of the following items:

- a literal data character
- the dot metacharacter
- the \C escape sequence
- the \X escape sequence (in UTF-8 mode with Unicode properties)
- the \R escape sequence
- an escape such as \d that matches a single character
- a character class
- a back reference (see next section)
- a parenthesized subpattern (unless it is an assertion)

The general repetition quantifier specifies a minimum and maximum number of permitted matches, by giving the two numbers in curly brackets (braces), separated by a comma. The numbers must be less than 65536, and the first must be less than or equal to the second. For example:

z{2,4}

matches "zz", "zzz", or "zzzz". A closing brace on its own is not a special character. If the second number is omitted, but the comma is present, there is no upper limit; if the second number and the comma are both omitted, the quantifier specifies an exact number of required matches. Thus

[aeiou]{3,}

matches at least 3 successive vowels, but may match many more, while

\d{8}

matches exactly 8 digits. An opening curly bracket that appears in a position where a quantifier is not allowed, or one that does not match the syntax of a quantifier, is taken as a literal character. For example, {,6} is not a quantifier, but a literal string of four characters.

In UTF-8 mode, quantifiers apply to UTF-8 characters rather than to individual bytes. Thus, for example, \x{100}{2} matches two UTF-8 characters, each of which is represented by a two-byte sequence. Similarly, when Unicode property support is available, \X{3} matches three Unicode extended sequences, each of which may be several bytes long (and they may be of different lengths).

The quantifier {0} is permitted, causing the expression to behave as if the previous item and the quantifier were not present.

For convenience, the three most common quantifiers have single-character abbreviations:

*
: is equivalent to {0,}

+
: is equivalent to {1,}

?
: is equivalent to {0,1}

It is possible to construct infinite loops by following a subpattern that can match no characters with a quantifier that has no upper limit, for example:

(a?)*

Earlier versions of Perl and PCRE used to give an error at compile time for such patterns. However, because there are cases where this can be useful, such patterns are now accepted, but if any repetition of the subpattern does in fact match no characters, the loop is forcibly broken.

By default, the quantifiers are "greedy", that is, they match as much as possible (up to the maximum number of permitted times), without causing the rest of the pattern to fail. The classic example of where this gives problems is in trying to match comments in C programs. These appear between /* and */ and within the comment, individual * and / characters may appear. An attempt to match C comments by applying the pattern

/*.**/

to the string

/* first comment */ not comment /* second comment */

fails, because it matches the entire string owing to the greediness of the .* item.

However, if a quantifier is followed by a question mark, it ceases to be greedy, and instead matches the minimum number of times possible, so the pattern

/*.*?*/

does the right thing with the C comments. The meaning of the various quantifiers is not otherwise changed, just the preferred number of matches. Do not confuse this use of question mark with its use as a quantifier in its own right. Because it has two uses, it can sometimes appear doubled, as in

\d??\d

which matches one digit by preference, but can match two if that is the only way the rest of the pattern matches.

If the `ungreedy` option is set (an option that is not available in Perl), the quantifiers are not greedy by default, but individual ones can be made greedy by following them with a question mark. In other words, it inverts the default behaviour.

When a parenthesized subpattern is quantified with a minimum repeat count that is greater than 1 or with a limited maximum, more memory is required for the compiled pattern, in proportion to the size of the minimum or maximum.

If a pattern starts with .* or .{0,} and the `dotall` option (equivalent to Perl's /s) is set, thus allowing the dot to match newlines, the pattern is implicitly anchored, because whatever follows will be tried against every character position in the subject string, so there is no point in retrying the overall match at any position after the first. PCRE normally treats such a pattern as though it were preceded by \A.

In cases where it is known that the subject string contains no newlines, it is worth setting `dotall` in order to obtain this optimization, or alternatively using ^ to indicate anchoring explicitly.

However, there is one situation where the optimization cannot be used. When .* is inside capturing parentheses that are the subject of a backreference elsewhere in the pattern, a match at the start may fail where a later one succeeds. Consider, for example:

(.*)abc\1

If the subject is "xyz123abc123" the match point is the fourth character. For this reason, such a pattern is not implicitly anchored.

When a capturing subpattern is repeated, the value captured is the substring that matched the final iteration. For example, after

(tweedle[dume]{3}\s*)+

has matched "tweedledum tweedledee" the value of the captured substring is "tweedledee". However, if there are nested capturing subpatterns, the corresponding captured values may have been set in previous iterations. For example, after

/(a|(b))+/

matches "aba" the value of the second captured substring is "b".

Atomic grouping and possessive quantifiers

With both maximizing ("greedy") and minimizing ("ungreedy" or "lazy") repetition, failure of what follows normally causes the repeated item to be re-evaluated to see if a different number of repeats allows the rest of the pattern to match. Sometimes it is useful to prevent this, either to change the nature of the match, or to cause it fail earlier than it otherwise might, when the author of the pattern knows there is no point in carrying on.

Consider, for example, the pattern \d+foo when applied to the subject line

123456bar

After matching all 6 digits and then failing to match "foo", the normal action of the matcher is to try again with only 5 digits matching the \d+ item, and then with 4, and so on, before ultimately failing. "Atomic grouping" (a term taken from Jeffrey Friedl's book) provides the means for specifying that once a subpattern has matched, it is not to be re-evaluated in this way.

If we use atomic grouping for the previous example, the matcher gives up immediately on failing to match "foo" the first time. The notation is a kind of special parenthesis, starting with (?> as in this example:

(?>\d+)foo

This kind of parenthesis "locks up" the part of the pattern it contains once it has matched, and a failure further into the pattern is prevented from backtracking into it. Backtracking past it to previous items, however, works as normal.

An alternative description is that a subpattern of this type matches the string of characters that an identical standalone pattern would match, if anchored at the current point in the subject string.

Atomic grouping subpatterns are not capturing subpatterns. Simple cases such as the above example can be thought of as a maximizing repeat that must swallow everything it can. So, while both \d+ and \d+? are prepared to adjust the number of digits they match in order to make the rest of the pattern match, (?>\d+) can only match an entire sequence of digits.

Atomic groups in general can of course contain arbitrarily complicated subpatterns, and can be nested. However, when the subpattern for an atomic group is just a single repeated item, as in the example above, a simpler notation, called a "possessive quantifier" can be used. This consists of an additional + character following a quantifier. Using this notation, the previous example can be rewritten as

\d++foo

Note that a possessive quantifier can be used with an entire group, for example:

(abc|xyz){2,3}+

Possessive quantifiers are always greedy; the setting of the `ungreedy` option is ignored. They are a convenient notation for the simpler forms of atomic group. However, there is no difference in the meaning of a possessive quantifier and the equivalent atomic group, though there may be a performance difference; possessive quantifiers should be slightly faster.

The possessive quantifier syntax is an extension to the Perl 5.8 syntax. Jeffrey Friedl originated the idea (and the name) in the first edition of his book. Mike McCloskey liked it, so implemented it when he built Sun's Java package, and PCRE copied it from there. It ultimately found its way into Perl at release 5.10.

PCRE has an optimization that automatically "possessifies" certain simple pattern constructs. For example, the sequence A+B is treated as A++B because there is no point in backtracking into a sequence of A's when B must follow.

When a pattern contains an unlimited repeat inside a subpattern that can itself be repeated an unlimited number of times, the use of an atomic group is the only way to avoid some failing matches taking a very long time indeed. The pattern

(\D+|<\d+>)*[!?]

matches an unlimited number of substrings that either consist of non-digits, or digits enclosed in <>, followed by either ! or ?. When it matches, it runs quickly. However, if it is applied to

aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa

it takes a long time before reporting failure. This is because the string can be divided between the internal \D+ repeat and the external * repeat in a large number of ways, and all have to be tried. (The example uses [!?] rather than a single character at the end, because both PCRE and Perl have an optimization that allows for fast failure when a single character is used. They remember the last single character that is required for a match, and fail early if it is not present in the string.) If the pattern is changed so that it uses an atomic group, like this:

((?>\D+)|<\d+>)*[!?]

sequences of non-digits cannot be broken, and failure happens quickly.

Back references

Outside a character class, a backslash followed by a digit greater than 0 (and possibly further digits) is a back reference to a capturing subpattern earlier (that is, to its left) in the pattern, provided there have been that many previous capturing left parentheses.

However, if the decimal number following the backslash is less than 10, it is always taken as a back reference, and causes an error only if there are not that many capturing left parentheses in the entire pattern. In other words, the parentheses that are referenced need not be to the left of the reference for numbers less than 10. A "forward back reference" of this type can make sense when a repetition is involved and the subpattern to the right has participated in an earlier iteration.

It is not possible to have a numerical "forward back reference" to a subpattern whose number is 10 or more using this syntax because a sequence such as \50 is interpreted as a character defined in octal. See the subsection entitled "Non-printing characters" above for further details of the handling of digits following a backslash. There is no such problem when named parentheses are used. A back reference to any subpattern is possible using named parentheses (see below).

Another way of avoiding the ambiguity inherent in the use of digits following a backslash is to use the \g escape sequence, which is a feature introduced in Perl 5.10. This escape must be followed by an unsigned number or a negative number, optionally enclosed in braces. These examples are all identical:

- (ring), \1
- (ring), \g1
- (ring), \g{1}

An unsigned number specifies an absolute reference without the ambiguity that is present in the older syntax. It is also useful when literal digits follow the reference. A negative number is a relative reference. Consider this example:

(abc(def)ghi)\g{-1}

The sequence \g{-1} is a reference to the most recently started capturing subpattern before \g, that is, is it equivalent to \2. Similarly, \g{-2} would be equivalent to \1. The use of relative references can be helpful in long patterns, and also in patterns that are created by joining together fragments that contain references within themselves.

A back reference matches whatever actually matched the capturing subpattern in the current subject string, rather than anything matching the subpattern itself (see "Subpatterns as subroutines" below for a way of doing that). So the pattern

(sens|respons)e and \1ibility

matches "sense and sensibility" and "response and responsibility", but not "sense and responsibility". If caseful matching is in force at the time of the back reference, the case of letters is relevant. For example,

((?i)rah)\s+\1

matches "rah rah" and "RAH RAH", but not "RAH rah", even though the original capturing subpattern is matched caselessly.

There are several different ways of writing back references to named subpatterns. The .NET syntax \k{name} and the Perl syntax \k<name> or \k'name' are supported, as is the Python syntax (?P=name). Perl 5.10's unified back reference syntax, in which \g can be used for both numeric and named references, is also supported. We could rewrite the above example in any of the following ways:

- (?<p1>(?i)rah)\s+\k<p1>
- (?'p1'(?i)rah)\s+\k{p1}
- (?P<p1>(?i)rah)\s+(?P=p1)
- (?<p1>(?i)rah)\s+\g{p1}

A subpattern that is referenced by name may appear in the pattern before or after the reference.

There may be more than one back reference to the same subpattern. If a subpattern has not actually been used in a particular match, any back references to it always fail. For example, the pattern

(a|(bc))\2

always fails if it starts to match "a" rather than "bc". Because there may be many capturing parentheses in a pattern, all digits following the backslash are taken as part of a potential back reference number. If the pattern continues with a digit character, some delimiter must be used to terminate the back reference. If the `extended` option is set, this can be whitespace. Otherwise an empty comment (see "Comments" below) can be used.

A back reference that occurs inside the parentheses to which it refers fails when the subpattern is first used, so, for example, (a\1) never matches. However, such references can be useful inside repeated subpatterns. For example, the pattern

(a|b\1)+

matches any number of "a"s and also "aba", "ababbaa" etc. At each iteration of the subpattern, the back reference matches the character string corresponding to the previous iteration. In order for this to work, the pattern must be such that the first iteration does not need to match the back reference. This can be done using alternation, as in the example above, or by a quantifier with a minimum of zero.

Assertions

An assertion is a test on the characters following or preceding the current matching point that does not actually consume any characters. The simple assertions coded as \b, \B, \A, \G, \Z, \z, ^ and $ are described above.

More complicated assertions are coded as subpatterns. There are two kinds: those that look ahead of the current position in the subject string, and those that look behind it. An assertion subpattern is matched in the normal way, except that it does not cause the current matching position to be changed.

Assertion subpatterns are not capturing subpatterns, and may not be repeated, because it makes no sense to assert the same thing several times. If any kind of assertion contains capturing subpatterns within it, these are counted for the purposes of numbering the capturing subpatterns in the whole pattern. However, substring capturing is carried out only for positive assertions, because it does not make sense for negative assertions.

Lookahead assertions

Lookahead assertions start with (?= for positive assertions and (?! for negative assertions. For example,

\w+(?=;)

matches a word followed by a semicolon, but does not include the semicolon in the match, and

foo(?!bar)

matches any occurrence of "foo" that is not followed by "bar". Note that the apparently similar pattern

(?!foo)bar

does not find an occurrence of "bar" that is preceded by something other than "foo"; it finds any occurrence of "bar" whatsoever, because the assertion (?!foo) is always true when the next three characters are "bar". A lookbehind assertion is needed to achieve the other effect.

If you want to force a matching failure at some point in a pattern, the most convenient way to do it is with (?!) because an empty string always matches, so an assertion that requires there not to be an empty string must always fail.

Lookbehind assertions

Lookbehind assertions start with (?<= for positive assertions and (?<! for negative assertions. For example,

(?<!foo)bar

does find an occurrence of "bar" that is not preceded by "foo". The contents of a lookbehind assertion are restricted such that all the strings it matches must have a fixed length. However, if there are several top-level alternatives, they do not all have to have the same fixed length. Thus

(?<=bullock|donkey)

is permitted, but

(?<!dogs?|cats?)

causes an error at compile time. Branches that match different length strings are permitted only at the top level of a lookbehind assertion. This is an extension compared with Perl (at least for 5.8), which requires all branches to match the same length of string. An assertion such as

(?<=ab(c|de))

is not permitted, because its single top-level branch can match two different lengths, but it is acceptable if rewritten to use two top-level branches:

(?<=abc|abde)

In some cases, the Perl 5.10 escape sequence \K (see above) can be used instead of a lookbehind assertion; this is not restricted to a fixed-length.

The implementation of lookbehind assertions is, for each alternative, to temporarily move the current position back by the fixed length and then try to match. If there are insufficient characters before the current position, the assertion fails.

PCRE does not allow the \C escape (which matches a single byte in UTF-8 mode) to appear in lookbehind assertions, because it makes it impossible to calculate the length of the lookbehind. The \X and \R escapes, which can match different numbers of bytes, are also not permitted.

Possessive quantifiers can be used in conjunction with lookbehind assertions to specify efficient matching at the end of the subject string. Consider a simple pattern such as

abcd$

when applied to a long string that does not match. Because matching proceeds from left to right, PCRE will look for each "a" in the subject and then see if what follows matches the rest of the pattern. If the pattern is specified as

^.*abcd$

the initial .* matches the entire string at first, but when this fails (because there is no following "a"), it backtracks to match all but the last character, then all but the last two characters, and so on. Once again the search for "a" covers the entire string, from right to left, so we are no better off. However, if the pattern is written as

^.*+(?<=abcd)

there can be no backtracking for the .*+ item; it can match only the entire string. The subsequent lookbehind assertion does a single test on the last four characters. If it fails, the match fails immediately. For long strings, this approach makes a significant difference to the processing time.

Using multiple assertions

Several assertions (of any sort) may occur in succession. For example,

(?<=\d{3})(?<!999)foo

matches "foo" preceded by three digits that are not "999". Notice that each of the assertions is applied independently at the same point in the subject string. First there is a check that the previous three characters are all digits, and then there is a check that the same three characters are not "999". This pattern does *not* match "foo" preceded by six characters, the first of which are digits and the last three of which are not "999". For example, it doesn't match "123abcfoo". A pattern to do that is

(?<=\d{3}...)(?<!999)foo

This time the first assertion looks at the preceding six characters, checking that the first three are digits, and then the second assertion checks that the preceding three characters are not "999".

Assertions can be nested in any combination. For example,

(?<=(?<!foo)bar)baz

matches an occurrence of "baz" that is preceded by "bar" which in turn is not preceded by "foo", while

(?<=\d{3}(?!999)...)foo

is another pattern that matches "foo" preceded by three digits and any three characters that are not "999".

Conditional subpatterns

It is possible to cause the matching process to obey a subpattern conditionally or to choose between two alternative subpatterns, depending on the result of an assertion, or whether a previous capturing subpattern matched or not. The two possible forms of conditional subpattern are

- (?(condition)yes-pattern)
- (?(condition)yes-pattern|no-pattern)

If the condition is satisfied, the yes-pattern is used; otherwise the no-pattern (if present) is used. If there are more than two alternatives in the subpattern, a compile-time error occurs.

There are four kinds of condition: references to subpatterns, references to recursion, a pseudo-condition called DEFINE, and assertions.

Checking for a used subpattern by number

If the text between the parentheses consists of a sequence of digits, the condition is true if the capturing subpattern of that number has previously matched. An alternative notation is to precede the digits with a plus or minus sign. In this case, the subpattern number is relative rather than absolute. The most recently opened parentheses can be referenced by (?(-1), the next most recent by (?(-2), and so on. In looping constructs it can also make sense to refer to subsequent groups with constructs such as (?(+2).

Consider the following pattern, which contains non-significant whitespace to make it more readable (assume the `extended` option) and to divide it into three parts for ease of discussion:

(\( )? [^()]+ (?(1) \))

The first part matches an optional opening parenthesis, and if that character is present, sets it as the first captured substring. The second part matches one or more characters that are not parentheses. The third part is a conditional subpattern that tests whether the first set of parentheses matched or not. If they did, that is, if subject started with an opening parenthesis, the condition is true, and so the yes-pattern is executed and a closing parenthesis is required. Otherwise, since no-pattern is not present, the subpattern matches nothing. In other words, this pattern matches a sequence of non-parentheses, optionally enclosed in parentheses.

If you were embedding this pattern in a larger one, you could use a relative reference:

...other stuff... (\( )? [^()]+ (?(-1) \)) ...

This makes the fragment independent of the parentheses in the larger pattern.

Checking for a used subpattern by name

Perl uses the syntax (?(<name>)...) or (?('name')...) to test for a used subpattern by name. For compatibility with earlier versions of PCRE, which had this facility before Perl, the syntax (?(name)...) is also recognized. However, there is a possible ambiguity with this syntax, because subpattern names may consist entirely of digits. PCRE looks first for a named subpattern; if it cannot find one and the name consists entirely of digits, PCRE looks for a subpattern of that number, which must be greater than zero. Using subpattern names that consist entirely of digits is not recommended.

Rewriting the above example to use a named subpattern gives this:

(?<OPEN> \( )? [^()]+ (?(<OPEN>) \))

Checking for pattern recursion

If the condition is the string (R), and there is no subpattern with the name R, the condition is true if a recursive call to the whole pattern or any subpattern has been made. If digits or a name preceded by ampersand follow the letter R, for example:

(?(R3)...) or (?(R&name)...)

the condition is true if the most recent recursion is into the subpattern whose number or name is given. This condition does not check the entire recursion stack.

At "top level", all these recursion test conditions are false. Recursive patterns are described below.

Defining subpatterns for use by reference only

If the condition is the string (DEFINE), and there is no subpattern with the name DEFINE, the condition is always false. In this case, there may be only one alternative in the subpattern. It is always skipped if control reaches this point in the pattern; the idea of DEFINE is that it can be used to define "subroutines" that can be referenced from elsewhere. (The use of "subroutines" is described below.) For example, a pattern to match an IPv4 address could be written like this (ignore whitespace and line breaks):

(?(DEFINE) (?<byte> 2[0-4]\d | 25[0-5] | 1\d\d | [1-9]?\d)) \b (?&byte) (\.(?&byte)){3} \b

The first part of the pattern is a DEFINE group inside which a another group named "byte" is defined. This matches an individual component of an IPv4 address (a number less than 256). When matching takes place, this part of the pattern is skipped because DEFINE acts like a false condition.

The rest of the pattern uses references to the named group to match the four dot-separated components of an IPv4 address, insisting on a word boundary at each end.

Assertion conditions

If the condition is not in any of the above formats, it must be an assertion. This may be a positive or negative lookahead or lookbehind assertion. Consider this pattern, again containing non-significant whitespace, and with the two alternatives on the second line:

```
(?(?=[^a-z]*[a-z])
\d{2}-[a-z]{3}-\d{2}  |  \d{2}-\d{2}-\d{2} )
```

The condition is a positive lookahead assertion that matches an optional sequence of non-letters followed by a letter. In other words, it tests for the presence of at least one letter in the subject. If a letter is found, the subject is matched against the first alternative; otherwise it is matched against the second. This pattern matches strings in one of the two forms dd-aaa-dd or dd-dd-dd, where aaa are letters and dd are digits.

Comments

The sequence (?# marks the start of a comment that continues up to the next closing parenthesis. Nested parentheses are not permitted. The characters that make up a comment play no part in the pattern matching at all.

If the `extended` option is set, an unescaped # character outside a character class introduces a comment that continues to immediately after the next newline in the pattern.

Recursive patterns

Consider the problem of matching a string in parentheses, allowing for unlimited nested parentheses. Without the use of recursion, the best that can be done is to use a pattern that matches up to some fixed depth of nesting. It is not possible to handle an arbitrary nesting depth.

For some time, Perl has provided a facility that allows regular expressions to recurse (amongst other things). It does this by interpolating Perl code in the expression at run time, and the code can refer to the expression itself. A Perl pattern using code interpolation to solve the parentheses problem can be created like this:

$re = qr{\((?: (?>[^()]+) | (?p{$re}))* \)}x;

The (?p{...}) item interpolates Perl code at run time, and in this case refers recursively to the pattern in which it appears.

Obviously, PCRE cannot support the interpolation of Perl code. Instead, it supports special syntax for recursion of the entire pattern, and also for individual subpattern recursion. After its introduction in PCRE and Python, this kind of recursion was introduced into Perl at release 5.10.

A special item that consists of (? followed by a number greater than zero and a closing parenthesis is a recursive call of the subpattern of the given number, provided that it occurs inside that subpattern. (If not, it is a "subroutine" call, which is described in the next section.) The special item (?R) or (?0) is a recursive call of the entire regular expression.

In PCRE (like Python, but unlike Perl), a recursive subpattern call is always treated as an atomic group. That is, once it has matched some of the subject string, it is never re-entered, even if it contains untried alternatives and there is a subsequent matching failure.

This PCRE pattern solves the nested parentheses problem (assume the `extended` option is set so that whitespace is ignored):

\(((?>[^()]+) | (?R))* \)

First it matches an opening parenthesis. Then it matches any number of substrings which can either be a sequence of non-parentheses, or a recursive match of the pattern itself (that is, a correctly parenthesized substring). Finally there is a closing parenthesis.

If this were part of a larger pattern, you would not want to recurse the entire pattern, so instead you could use this:

(\(((?>[^()]+) | (?1))* \))

We have put the pattern into parentheses, and caused the recursion to refer to them instead of the whole pattern.

In a larger pattern, keeping track of parenthesis numbers can be tricky. This is made easier by the use of relative references. (A Perl 5.10 feature.) Instead of (?1) in the pattern above you can write (?-2) to refer to the second most recently opened parentheses preceding the recursion. In other words, a negative number counts capturing parentheses leftwards from the point at which it is encountered.

It is also possible to refer to subsequently opened parentheses, by writing references such as (?+2). However, these cannot be recursive because the reference is not inside the parentheses that are referenced. They are always "subroutine" calls, as described in the next section.

An alternative approach is to use named parentheses instead. The Perl syntax for this is (?&name); PCRE's earlier syntax (?P>name) is also supported. We could rewrite the above example as follows:

(?<pn> \(((?>[^()]+) | (?&pn))* \))

If there is more than one subpattern with the same name, the earliest one is used.

This particular example pattern that we have been looking at contains nested unlimited repeats, and so the use of atomic grouping for matching strings of non-parentheses is important when applying the pattern to strings that do not match. For example, when this pattern is applied to

(aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa()

it yields "no match" quickly. However, if atomic grouping is not used, the match runs for a very long time indeed because there are so many different ways the + and * repeats can carve up the subject, and all have to be tested before failure can be reported.

At the end of a match, the values set for any capturing subpatterns are those from the outermost level of the recursion at which the subpattern value is set. If the pattern above is matched against

(ab(cd)ef)

the value for the capturing parentheses is "ef", which is the last value taken on at the top level. If additional parentheses are added, giving

```
\( ( ( (?>[^()]+) | (?R) )* ) \)
   ^                        ^
   ^                        ^
```

the string they capture is "ab(cd)ef", the contents of the top level parentheses.

Do not confuse the (?R) item with the condition (R), which tests for recursion. Consider this pattern, which matches text in angle brackets, allowing for arbitrary nesting. Only digits are allowed in nested brackets (that is, when recursing), whereas any characters are permitted at the outer level.

< (?: (?(R) \d++ | [^<>]*+) | (?R)) * >

In this pattern, (?(R) is the start of a conditional subpattern, with two different alternatives for the recursive and non-recursive cases. The (?R) item is the actual recursive call.

Subpatterns as subroutines

If the syntax for a recursive subpattern reference (either by number or by name) is used outside the parentheses to which it refers, it operates like a subroutine in a programming language. The "called" subpattern may be defined before or after the reference. A numbered reference can be absolute or relative, as in these examples:

- (...(absolute)...)...(?2)...
- (...(relative)...)...(?-1)...
- (...(?+1)...(relative)...

An earlier example pointed out that the pattern

(sens|respons)e and \1ibility

matches "sense and sensibility" and "response and responsibility", but not "sense and responsibility". If instead the pattern

(sens|respons)e and (?1)ibility

is used, it does match "sense and responsibility" as well as the other two strings. Another example is given in the discussion of DEFINE above.

Like recursive subpatterns, a "subroutine" call is always treated as an atomic group. That is, once it has matched some of the subject string, it is never re-entered, even if it contains untried alternatives and there is a subsequent matching failure.

When a subpattern is used as a subroutine, processing options such as case-independence are fixed when the subpattern is defined. They cannot be changed for different calls. For example, consider this pattern:

(abc)(?i:(?-1))

It matches "abcabc". It does not match "abcABC" because the change of processing option does not affect the called subpattern.

Backtracking control

Perl 5.10 introduced a number of "Special Backtracking Control Verbs", which are described in the Perl documentation as "experimental and subject to change or removal in a future version of Perl". It goes on to say: "Their usage in production code should be noted to avoid problems during upgrades." The same remarks apply to the PCRE features described in this section.

The new verbs make use of what was previously invalid syntax: an opening parenthesis followed by an asterisk. In Perl, they are generally of the form (*VERB:ARG) but PCRE does not support the use of arguments, so its general form is just (*VERB). Any number of these verbs may occur in a pattern. There are two kinds:

Verbs that act immediately

The following verbs act as soon as they are encountered:

(*ACCEPT)

This verb causes the match to end successfully, skipping the remainder of the pattern. When inside a recursion, only the innermost pattern is ended immediately. PCRE differs from Perl in what happens if the (*ACCEPT) is inside capturing parentheses. In Perl, the data so far is captured: in PCRE no data is captured. For example:

A(A|B(*ACCEPT)|C)D

This matches "AB", "AAD", or "ACD", but when it matches "AB", no data is captured.

(*FAIL) or (*F)

This verb causes the match to fail, forcing backtracking to occur. It is equivalent to (?!) but easier to read. The Perl documentation notes that it is probably useful only when combined with (?{}) or (??{}). Those are, of course, Perl features that are not present in PCRE. The nearest equivalent is the callout feature, as for example in this pattern:

a+(?C)(*FAIL)

A match with the string "aaaa" always fails, but the callout is taken before each backtrack happens (in this example, 10 times).

Verbs that act after backtracking

The following verbs do nothing when they are encountered. Matching continues with what follows, but if there is no subsequent match, a failure is forced. The verbs differ in exactly what kind of failure occurs.

(*COMMIT)

This verb causes the whole match to fail outright if the rest of the pattern does not match. Even if the pattern is unanchored, no further attempts to find a match by advancing the start point take place. Once (*COMMIT) has been passed, `re:run/3` is committed to finding a match at the current starting point, or not at all. For example:

a+(*COMMIT)b

This matches "xxaab" but not "aacaab". It can be thought of as a kind of dynamic anchor, or "I've started, so I must finish."

(*PRUNE)

This verb causes the match to fail at the current position if the rest of the pattern does not match. If the pattern is unanchored, the normal "bumpalong" advance to the next starting character then happens. Backtracking can occur as usual to the left of (*PRUNE), or when matching to the right of (*PRUNE), but if there is no match to the right,

backtracking cannot cross (*PRUNE). In simple cases, the use of (*PRUNE) is just an alternative to an atomic group or possessive quantifier, but there are some uses of (*PRUNE) that cannot be expressed in any other way.

(*SKIP)

This verb is like (*PRUNE), except that if the pattern is unanchored, the "bumpalong" advance is not to the next character, but to the position in the subject where (*SKIP) was encountered. (*SKIP) signifies that whatever text was matched leading up to it cannot be part of a successful match. Consider:

a+(*SKIP)b

If the subject is "aaaac...", after the first match attempt fails (starting at the first character in the string), the starting point skips on to start the next attempt at "c". Note that a possessive quantifier does not have the same effect in this example; although it would suppress backtracking during the first match attempt, the second attempt would start at the second character instead of skipping on to "c".

(*THEN)

This verb causes a skip to the next alternation if the rest of the pattern does not match. That is, it cancels pending backtracking, but only within the current alternation. Its name comes from the observation that it can be used for a pattern-based if-then-else block:

(COND1 (*THEN) FOO | COND2 (*THEN) BAR | COND3 (*THEN) BAZ) ...

If the COND1 pattern matches, FOO is tried (and possibly further items after the end of the group if FOO succeeds); on failure the matcher skips to the second alternative and tries COND2, without backtracking into COND1. If (*THEN) is used outside of any alternation, it acts exactly like (*PRUNE).

regexp

Erlang module

> ## Note:
>
> This module has been obsoleted by the *re* module and will be removed in a future release.

This module contains functions for regular expression matching and substitution.

Exports

`match(String, RegExp) -> MatchRes`

Types:

String = RegExp = string()

MatchRes = {match,Start,Length} | nomatch | {error,errordesc()}

Start = Length = integer()

Finds the first, longest match of the regular expression `RegExp` in `String`. This function searches for the longest possible match and returns the first one found if there are several expressions of the same length. It returns as follows:

`{match,Start,Length}`

if the match succeeded. `Start` is the starting position of the match, and `Length` is the length of the matching string.

`nomatch`

if there were no matching characters.

`{error,Error}`

if there was an error in `RegExp`.

`first_match(String, RegExp) -> MatchRes`

Types:

String = RegExp = string()

MatchRes = {match,Start,Length} | nomatch | {error,errordesc()}

Start = Length = integer()

Finds the first match of the regular expression `RegExp` in `String`. This call is usually faster than `match` and it is also a useful way to ascertain that a match exists. It returns as follows:

`{match,Start,Length}`

if the match succeeded. `Start` is the starting position of the match and `Length` is the length of the matching string.

`nomatch`

if there were no matching characters.

`{error,Error}`

if there was an error in `RegExp`.

matches(String, RegExp) -> MatchRes

Types:

String = RegExp = string()

MatchRes = {match, Matches} | {error, errordesc()}

Matches = list()

Finds all non-overlapping matches of the expression `RegExp` in `String`. It returns as follows:

`{match, Matches}`

if the regular expression was correct. The list will be empty if there was no match. Each element in the list looks like `{Start, Length}`, where `Start` is the starting position of the match, and `Length` is the length of the matching string.

`{error,Error}`

if there was an error in `RegExp`.

sub(String, RegExp, New) -> SubRes

Types:

String = RegExp = New = string()

SubRes = {ok,NewString,RepCount} | {error,errordesc()}

RepCount = integer()

Substitutes the first occurrence of a substring matching `RegExp` in `String` with the string `New`. A & in the string `New` is replaced by the matched substring of `String`. \& puts a literal & into the replacement string. It returns as follows:

`{ok,NewString,RepCount}`

if `RegExp` is correct. `RepCount` is the number of replacements which have been made (this will be either 0 or 1).

`{error, Error}`

if there is an error in `RegExp`.

gsub(String, RegExp, New) -> SubRes

Types:

String = RegExp = New = string()

SubRes = {ok,NewString,RepCount} | {error,errordesc()}

RepCount = integer()

The same as `sub`, except that all non-overlapping occurrences of a substring matching `RegExp` in `String` are replaced by the string `New`. It returns:

`{ok,NewString,RepCount}`

if `RegExp` is correct. `RepCount` is the number of replacements which have been made.

`{error, Error}`

if there is an error in `RegExp`.

split(String, RegExp) -> SplitRes

Types:

String = RegExp = string()

SubRes = {ok,FieldList} | {error,errordesc()}

Fieldlist = [string()]

`String` is split into fields (sub-strings) by the regular expression `RegExp`.

If the separator expression is `" "` (a single space), then the fields are separated by blanks and/or tabs and leading and trailing blanks and tabs are discarded. For all other values of the separator, leading and trailing blanks and tabs are not discarded. It returns:

`{ok, FieldList}`

to indicate that the string has been split up into the fields of `FieldList`.

`{error, Error}`

if there is an error in `RegExp`.

sh_to_awk(ShRegExp) -> AwkRegExp

Types:

ShRegExp AwkRegExp = string()

SubRes = {ok,NewString,RepCount} | {error,errordesc()}

RepCount = integer()

Converts the `sh` type regular expression `ShRegExp` into a full `AWK` regular expression. Returns the converted regular expression string. `sh` expressions are used in the shell for matching file names and have the following special characters:

`*`

matches any string including the null string.

`?`

matches any single character.

`[...]`

matches any of the enclosed characters. Character ranges are specified by a pair of characters separated by a `-`. If the first character after `[` is a `!`, then any character not enclosed is matched.

It may sometimes be more practical to use `sh` type expansions as they are simpler and easier to use, even though they are not as powerful.

parse(RegExp) -> ParseRes

Types:

RegExp = string()

ParseRes = {ok,RE} | {error,errordesc()}

Parses the regular expression `RegExp` and builds the internal representation used in the other regular expression functions. Such representations can be used in all of the other functions instead of a regular expression string. This is more efficient when the same regular expression is used in many strings. It returns:

`{ok, RE}`if `RegExp`is correct and `RE`is the internal representation.

`{error, Error}`if there is an error in `RegExpString`.

`format_error(ErrorDescriptor) -> Chars`

Types:

ErrorDescriptor = errordesc()

Chars = [char() | Chars]

Returns a string which describes the error `ErrorDescriptor` returned when there is an error in a regular expression.

Regular Expressions

The regular expressions allowed here is a subset of the set found in `egrep` and in the `AWK` programming language, as defined in the book, `The AWK Programming Language, by A. V. Aho, B. W. Kernighan, P. J. Weinberger`. They are composed of the following characters:

c

matches the non-metacharacter `c`.

\c

matches the escape sequence or literal character `c`.

.

matches any character.

^

matches the beginning of a string.

$

matches the end of a string.

[abc...]

character class, which matches any of the characters `abc...` Character ranges are specified by a pair of characters separated by a `-`.

[^abc...]

negated character class, which matches any character except `abc...`.

r1 | r2

alternation. It matches either `r1` or `r2`.

r1r2

concatenation. It matches `r1` and then `r2`.

r+

matches one or more `r`s.

r*

matches zero or more `r`s.

r?

matches zero or one `r`s.

(r)

grouping. It matches `r`.

The escape sequences allowed are the same as for Erlang strings:

`\b`
: backspace

`\f`
: form feed

`\n`
: newline (line feed)

`\r`
: carriage return

`\t`
: tab

`\e`
: escape

`\v`
: vertical tab

`\s`
: space

`\d`
: delete

`\ddd`
: the octal value ddd

`\xhh`
: The hexadecimal value `hh`.

`\x{h...}`
: The hexadecimal value `h...`.

`\c`
: any other character literally, for example `\\` for backslash, `\"` for ")

To make these functions easier to use, in combination with the function `io:get_line` which terminates the input line with a new line, the `$` characters also matches a string ending with `"...\n"`. The following examples define Erlang data types:

```
Atoms     [a-z][0-9a-zA-Z_]*

Variables [A-Z_][0-9a-zA-Z_]*

Floats    (\+|-)?[0-9]+\.[0-9]+((E|e)(\+|-)?[0-9]+)?
```

Regular expressions are written as Erlang strings when used with the functions in this module. This means that any `\` or `"` characters in a regular expression string must be written with `\` as they are also escape characters for the string. For example, the regular expression string for Erlang floats is: `"(\\+|-)?[0-9]+\\.[0-9]+((E|e)(\\+|-)?[0-9]+)?"`.

It is not really necessary to have the escape sequences as part of the regular expression syntax as they can always be generated directly in the string. They are included for completeness and can they can also be useful when generating regular expressions, or when they are entered other than with Erlang strings.

sets =:=

Erlang module

Sets are collections of elements with no duplicate elements. The representation of a set is not defined.

This module provides exactly the same interface as the module `ordsets` but with a defined representation. One difference is that while this module considers two elements as different if they do not match (`=:=`), `ordsets` considers two elements as different if and only if they do not compare equal (`==`).

DATA TYPES

```
set()
  as returned by new/0
```

Exports

`new() -> Set`

Types:

Set = **set()**

Returns a new empty set.

`is_set(Set) -> bool()`

Types:

Set = **term()**

Returns `true` if `Set` is a set of elements, otherwise `false`.

`size(Set) -> int()`

Types:

Set = **term()**

Returns the number of elements in `Set`.

`to_list(Set) -> List`

Types:

Set = **set()**

List = **[term()]**

Returns the elements of `Set` as a list.

`from_list(List) -> Set`

Types:

List = **[term()]**

Set = **set()**

Returns an set of the elements in `List`.

`is_element(Element, Set) -> bool()`

Types:

Element = term()

Set = set()

Returns `true` if `Element` is an element of `Set`, otherwise `false`.

`add_element(Element, Set1) -> Set2`

Types:

Element = term()

Set1 = Set2 = set()

Returns a new set formed from `Set1` with `Element` inserted.

`del_element(Element, Set1) -> Set2`

Types:

Element = term()

Set1 = Set2 = set()

Returns `Set1`, but with `Element` removed.

`union(Set1, Set2) -> Set3`

Types:

Set1 = Set2 = Set3 = set()

Returns the merged (union) set of `Set1` and `Set2`.

`union(SetList) -> Set`

Types:

SetList = [set()]

Set = set()

Returns the merged (union) set of the list of sets.

`intersection(Set1, Set2) -> Set3`

Types:

Set1 = Set2 = Set3 = set()

Returns the intersection of `Set1` and `Set2`.

`intersection(SetList) -> Set`

Types:

SetList = [set()]

Set = set()

Returns the intersection of the non-empty list of sets.

`is_disjoint(Set1, Set2) -> bool()`

Types:

Set1 = Set2 = set()

Returns `true` if `Set1` and `Set2` are disjoint (have no elements in common), and `false` otherwise.

`subtract(Set1, Set2) -> Set3`

Types:

Set1 = Set2 = Set3 = set()

Returns only the elements of `Set1` which are not also elements of `Set2`.

`is_subset(Set1, Set2) -> bool()`

Types:

Set1 = Set2 = set()

Returns `true` when every element of `Set1` is also a member of `Set2`, otherwise `false`.

`fold(Function, Acc0, Set) -> Acc1`

Types:

Function = fun (E, AccIn) -> AccOut

Acc0 = Acc1 = AccIn = AccOut = term()

Set = set()

Fold `Function` over every element in `Set` returning the final value of the accumulator.

`filter(Pred, Set1) -> Set2`

Types:

Pred = fun (E) -> bool()

Set1 = Set2 = set()

Filter elements in `Set1` with boolean function `Fun`.

See Also

ordsets(3), *gb_sets(3)*

shell

Erlang module

The module `shell` implements an Erlang shell.

The shell is a user interface program for entering expression sequences. The expressions are evaluated and a value is returned. A history mechanism saves previous commands and their values, which can then be incorporated in later commands. How many commands and results to save can be determined by the user, either interactively, by calling `shell:history/1` and `shell:results/1`, or by setting the application configuration parameters `shell_history_length` and `shell_saved_results` for the application STDLIB.

The shell uses a helper process for evaluating commands in order to protect the history mechanism from exceptions. By default the evaluator process is killed when an exception occurs, but by calling `shell:catch_exception/1` or by setting the application configuration parameter `shell_catch_exception` for the application STDLIB this behavior can be changed. See also the example below.

Variable bindings, and local process dictionary changes which are generated in user expressions are preserved, and the variables can be used in later commands to access their values. The bindings can also be forgotten so the variables can be re-used.

The special shell commands all have the syntax of (local) function calls. They are evaluated as normal function calls and many commands can be used in one expression sequence.

If a command (local function call) is not recognized by the shell, an attempt is first made to find the function in the module `user_default`, where customized local commands can be placed. If found, then the function is evaluated. Otherwise, an attempt is made to evaluate the function in the module `shell_default`. The module `user_default` must be explicitly loaded.

The shell also permits the user to start multiple concurrent jobs. A job can be regarded as a set of processes which can communicate with the shell.

There is some support for reading and printing records in the shell. During compilation record expressions are translated to tuple expressions. In runtime it is not known whether a tuple actually represents a record. Nor are the record definitions used by compiler available at runtime. So in order to read the record syntax and print tuples as records when possible, record definitions have to be maintained by the shell itself. The shell commands for reading, defining, forgetting, listing, and printing records are described below. Note that each job has its own set of record definitions. To facilitate matters record definitions in the modules `shell_default` and `user_default` (if loaded) are read each time a new job is started. For instance, adding the line

```
-include_lib("kernel/include/file.hrl").
```

to `user_default` makes the definition of `file_info` readily available in the shell.

The shell runs in two modes:

- `Normal (possibly restricted)` mode, in which commands can be edited and expressions evaluated.
- Job Control Mode `JCL`, in which jobs can be started, killed, detached and connected.

Only the currently connected job can 'talk' to the shell.

Shell Commands

`b()`

Prints the current variable bindings.

`f()`

Removes all variable bindings.

`f(X)`

Removes the binding of variable `X`.

`h()`

Prints the history list.

`history(N)`

Sets the number of previous commands to keep in the history list to `N`. The previous number is returned. The default number is 20.

`results(N)`

Sets the number of results from previous commands to keep in the history list to `N`. The previous number is returned. The default number is 20.

`e(N)`

Repeats the command `N`, if `N` is positive. If it is negative, the `N`th previous command is repeated (i.e., `e(-1)` repeats the previous command).

`v(N)`

Uses the return value of the command `N` in the current command, if `N` is positive. If it is negative, the return value of the `N`th previous command is used (i.e., `v(-1)` uses the value of the previous command).

`help()`

Evaluates `shell_default:help()`.

`c(File)`

Evaluates `shell_default:c(File)`. This compiles and loads code in `File` and purges old versions of code, if necessary. Assumes that the file and module names are the same.

`catch_exception(Bool)`

Sets the exception handling of the evaluator process. The previous exception handling is returned. The default (`false`) is to kill the evaluator process when an exception occurs, which causes the shell to create a new evaluator process. When the exception handling is set to `true` the evaluator process lives on which means that for instance ports and ETS tables as well as processes linked to the evaluator process survive the exception.

`rd(RecordName, RecordDefinition)`

Defines a record in the shell. `RecordName` is an atom and `RecordDefinition` lists the field names and the default values. Usually record definitions are made known to the shell by use of the `rr` commands described below, but sometimes it is handy to define records on the fly.

`rf()`

Removes all record definitions, then reads record definitions from the modules `shell_default` and `user_default` (if loaded). Returns the names of the records defined.

`rf(RecordNames)`

Removes selected record definitions. `RecordNames` is a record name or a list of record names. Use `'_'` to remove all record definitions.

`rl()`

Prints all record definitions.

`rl(RecordNames)`

Prints selected record definitions. `RecordNames` is a record name or a list of record names.

`rp(Term)`

Prints a term using the record definitions known to the shell. All of `Term` is printed; the depth is not limited as is the case when a return value is printed.

`rr(Module)`

Reads record definitions from a module's BEAM file. If there are no record definitions in the BEAM file, the source file is located and read instead. Returns the names of the record definitions read. `Module` is an atom.

`rr(Wildcard)`

Reads record definitions from files. Existing definitions of any of the record names read are replaced. `Wildcard` is a wildcard string as defined in `filelib(3)` but not an atom.

`rr(WildcardOrModule, RecordNames)`

Reads record definitions from files but discards record names not mentioned in `RecordNames` (a record name or a list of record names).

`rr(WildcardOrModule, RecordNames, Options)`

Reads record definitions from files. The compiler options `{i, Dir}`, `{d, Macro}`, and `{d, Macro, Value}` are recognized and used for setting up the include path and macro definitions. Use `'_'` as value of `RecordNames` to read all record definitions.

Example

The following example is a long dialogue with the shell. Commands starting with `>` are inputs to the shell. All other lines are output from the shell. All commands in this example are explained at the end of the dialogue. .

```
strider 1> erl
Erlang (BEAM) emulator version 5.3 [hipe] [threads:0]

Eshell V5.3  (abort with ^G)
1>Str = "abcd".
"abcd"
2> L = length(Str).
4
3> Descriptor = {L, list_to_atom(Str)}.
{4,abcd}
4> L.
4
5> b().
Descriptor = {4,abcd}
L = 4
Str = "abcd"
ok
6> f(L).
ok
7> b().
Descriptor = {4,abcd}
Str = "abcd"
ok
8> f(L).
ok
9> {L, _} = Descriptor.
{4,abcd}
10> L.
```

```
4
11> {P, Q, R} = Descriptor.
** exception error: no match of right hand side value {4,abcd}
12> P.
* 1: variable 'P' is unbound **
13> Descriptor.
{4,abcd}
14>{P, Q} = Descriptor.
{4,abcd}
15> P.
4
16> f().
ok
17> put(aa, hello).
undefined
18> get(aa).
hello
19> Y = test1:demo(1).
11
20> get().
[{aa,worked}]
21> put(aa, hello).
worked
22> Z = test1:demo(2).
** exception error: no match of right hand side value 1
     in function  test1:demo/1
23> Z.
* 1: variable 'Z' is unbound **
24> get(aa).
hello
25> erase(), put(aa, hello).
undefined
26> spawn(test1, demo, [1]).
<0.57.0>
27> get(aa).
hello
28> io:format("hello hello\n").
hello hello
ok
29> e(28).
hello hello
ok
30> v(28).
ok
31> c(ex).
{ok,ex}
32> rr(ex).
[rec]
33> rl(rec).
-record(rec,{a,b = val()}).
ok
34> #rec{}.
** exception error: undefined shell command val/0
35> #rec{b = 3}.
#rec{a = undefined,b = 3}
36> rp(v(-1)).
#rec{a = undefined,b = 3}
ok
37> rd(rec, {f = orddict:new()}).
rec
38> #rec{}.
#rec{f = []}
ok
39> rd(rec, {c}), A.
* 1: variable 'A' is unbound **
```

```
40> #rec{}.
#rec{c = undefined}
ok
41> test1:loop(0).
Hello Number: 0
Hello Number: 1
Hello Number: 2
Hello Number: 3

User switch command
 --> i
 --> c
.
.
.
Hello Number: 3374
Hello Number: 3375
Hello Number: 3376
Hello Number: 3377
Hello Number: 3378
** exception exit: killed
42> E = ets:new(t, []).
17
43> ets:insert({d,1,2}).
** exception error: undefined function ets:insert/1
44> ets:insert(E, {d,1,2}).
** exception error: argument is of wrong type
     in function  ets:insert/2
        called as ets:insert(16,{d,1,2})
45> f(E).
ok
46> catch_exception(true).
false
47> E = ets:new(t, []).
18
48> ets:insert({d,1,2}).
* exception error: undefined function ets:insert/1
49> ets:insert(E, {d,1,2}).
true
50> halt().
strider 2>
```

Comments

Command 1 sets the variable `Str` to the string `"abcd"`.

Command 2 sets `L` to the length of the string evaluating the BIF `atom_to_list`.

Command 3 builds the tuple `Descriptor`.

Command 4 prints the value of the variable `L`.

Command 5 evaluates the internal shell command `b()`, which is an abbreviation of "bindings". This prints the current shell variables and their bindings. The `ok` at the end is the return value of the `b()` function.

Command 6 `f(L)` evaluates the internal shell command `f(L)` (abbreviation of "forget"). The value of the variable `L` is removed.

Command 7 prints the new bindings.

Command 8 has no effect since `L` has no value.

Command 9 performs a pattern matching operation on `Descriptor`, binding a new value to `L`.

Command 10 prints the current value of `L`.

Command 11 tries to match `{P, Q, R}` against `Descriptor` which is `{4, abc}`. The match fails and none of the new variables become bound. The printout starting with "`** exception error:`" is not the value of the expression (the expression had no value because its evaluation failed), but rather a warning printed by the system to inform the user that an error has occurred. The values of the other variables (`L`, `Str`, etc.) are unchanged.

Commands 12 and 13 show that `P` is unbound because the previous command failed, and that `Descriptor` has not changed.

Commands 14 and 15 show a correct match where `P` and `Q` are bound.

Command 16 clears all bindings.

The next few commands assume that `test1:demo(X)` is defined in the following way:

```
demo(X) ->
    put(aa, worked),
    X = 1,
    X + 10.
```

Commands 17 and 18 set and inspect the value of the item `aa` in the process dictionary.

Command 19 evaluates `test1:demo(1)`. The evaluation succeeds and the changes made in the process dictionary become visible to the shell. The new value of the dictionary item `aa` can be seen in command 20.

Commands 21 and 22 change the value of the dictionary item `aa` to `hello` and call `test1:demo(2)`. Evaluation fails and the changes made to the dictionary in `test1:demo(2)`, before the error occurred, are discarded.

Commands 23 and 24 show that `Z` was not bound and that the dictionary item `aa` has retained its original value.

Commands 25, 26 and 27 show the effect of evaluating `test1:demo(1)` in the background. In this case, the expression is evaluated in a newly spawned process. Any changes made in the process dictionary are local to the newly spawned process and therefore not visible to the shell.

Commands 28, 29 and 30 use the history facilities of the shell.

Command 29 is `e(28)`. This re-evaluates command 28. Command 30 is `v(28)`. This uses the value (result) of command 28. In the cases of a pure function (a function with no side effects), the result is the same. For a function with side effects, the result can be different.

The next few commands show some record manipulation. It is assumed that `ex.erl` defines a record like this:

```
-record(rec, {a, b = val()}).

val() ->
    3.
```

Commands 31 and 32 compiles the file `ex.erl` and reads the record definitions in `ex.beam`. If the compiler did not output any record definitions on the BEAM file, `rr(ex)` tries to read record definitions from the source file instead.

Command 33 prints the definition of the record named `rec`.

Command 34 tries to create a `rec` record, but fails since the function `val/0` is undefined. Command 35 shows the workaround: explicitly assign values to record fields that cannot otherwise be initialized.

Command 36 prints the newly created record using record definitions maintained by the shell.

Command 37 defines a record directly in the shell. The definition replaces the one read from the file `ex.beam`.

Command 38 creates a record using the new definition, and prints the result.

Command 39 and 40 show that record definitions are updated as side effects. The evaluation of the command fails but the definition of `rec` has been carried out.

For the next command, it is assumed that `test1:loop(N)` is defined in the following way:

```
loop(N) ->
    io:format("Hello Number: ~w~n", [N]),
    loop(N+1).
```

Command 41 evaluates `test1:loop(0)`, which puts the system into an infinite loop. At this point the user types `Control G`, which suspends output from the current process, which is stuck in a loop, and activates `JCL` mode. In `JCL` mode the user can start and stop jobs.

In this particular case, the `i` command ("interrupt") is used to terminate the looping program, and the `c` command is used to connect to the shell again. Since the process was running in the background before we killed it, there will be more printouts before the "`** exception exit: killed`" message is shown.

Command 42 creates an ETS table.

Command 43 tries to insert a tuple into the ETS table but the first argument (the table) is missing. The exception kills the evaluator process.

Command 44 corrects the mistake, but the ETS table has been destroyed since it was owned by the killed evaluator process.

Command 46 sets the exception handling of the evaluator process to `true`. The exception handling can also be set when starting Erlang, like this: `erl -stdlib shell_catch_exception true`.

Command 48 makes the same mistake as in command 43, but this time the evaluator process lives on. The single star at the beginning of the printout signals that the exception has been caught.

Command 49 successfully inserts the tuple into the ETS table.

The `halt()` command exits the Erlang runtime system.

JCL Mode

When the shell starts, it starts a single evaluator process. This process, together with any local processes which it spawns, is referred to as a `job`. Only the current job, which is said to be `connected`, can perform operations with standard IO. All other jobs, which are said to be `detached`, are `blocked` if they attempt to use standard IO.

All jobs which do not use standard IO run in the normal way.

The shell escape key `^G` (Control G) detaches the current job and activates `JCL` mode. The `JCL` mode prompt is "`-->`". If "`?`" is entered at the prompt, the following help message is displayed:

```
        --> ?
        c [nn]            - connect to job
        i [nn]            - interrupt job
        k [nn]            - kill job
        j                 - list all jobs
        s [shell]         - start local shell
        r [node [shell]]  - start remote shell
        q        - quit erlang
        ? | h             - this message
```

The `JCL` commands have the following meaning:

`c [nn]`

Connects to job number <nn> or the current job. The standard shell is resumed. Operations which use standard IO by the current job will be interleaved with user inputs to the shell.

`i [nn]`

Stops the current evaluator process for job number `nn` or the current job, but does not kill the shell process. Accordingly, any variable bindings and the process dictionary will be preserved and the job can be connected again. This command can be used to interrupt an endless loop.

`k [nn]`

Kills job number `nn` or the current job. All spawned processes in the job are killed, provided they have not evaluated the `group_leader/1` BIF and are located on the local machine. Processes spawned on remote nodes will not be killed.

`j`

Lists all jobs. A list of all known jobs is printed. The current job name is prefixed with '*'.

`s`

Starts a new job. This will be assigned the new index `[nn]` which can be used in references.

`s [shell]`

Starts a new job. This will be assigned the new index `[nn]` which can be used in references. If the optional argument `shell` is given, it is assumed to be a module that implements an alternative shell.

`r [node]`

Starts a remote job on `node`. This is used in distributed Erlang to allow a shell running on one node to control a number of applications running on a network of nodes. If the optional argument `shell` is given, it is assumed to be a module that implements an alternative shell.

`q`

Quits Erlang. Note that this option is disabled if Erlang is started with the ignore break, `+Bi`, system flag (which may be useful e.g. when running a restricted shell, see below).

`?`

Displays this message.

It is possible to alter the behavior of shell escape by means of the STDLIB application variable `shell_esc`. The value of the variable can be either `jcl` (`erl -stdlib shell_esc jcl`) or `abort` (`erl -stdlib shell_esc abort`). The first option sets ^G to activate `JCL` mode (which is also default behavior). The latter sets ^G to terminate the current shell and start a new one. `JCL` mode cannot be invoked when `shell_esc` is set to `abort`.

If you want an Erlang node to have a remote job active from the start (rather than the default local job), you start Erlang with the `-remsh` flag. Example: `erl -sname this_node -remsh other_node@other_host`

Restricted Shell

The shell may be started in a restricted mode. In this mode, the shell evaluates a function call only if allowed. This feature makes it possible to, for example, prevent a user from accidentally calling a function from the prompt that could harm a running system (useful in combination with the the system flag *`+Bi`*).

When the restricted shell evaluates an expression and encounters a function call or an operator application, it calls a callback function (with information about the function call in question). This callback function returns `true` to let the shell go ahead with the evaluation, or `false` to abort it. There are two possible callback functions for the user to implement:

```
local_allowed(Func, ArgList, State) -> {true,NewState} | {false,NewState}
```

to determine if the call to the local function `Func` with arguments `ArgList` should be allowed.

`non_local_allowed(FuncSpec, ArgList, State) -> {true,NewState} | {false,NewState} | {{redirect,NewFuncSpec,NewArgList},NewState}`

to determine if the call to non-local function `FuncSpec` (`{Module,Func}` or a fun) with arguments `ArgList` should be allowed. The return value `{redirect,NewFuncSpec,NewArgList}` can be used to let the shell evaluate some other function than the one specified by `FuncSpec` and `ArgList`.

These callback functions are in fact called from local and non-local evaluation function handlers, described in the *erl_eval* manual page. (Arguments in `ArgList` are evaluated before the callback functions are called.)

The `State` argument is a tuple `{ShellState,ExprState}`. The return value `NewState` has the same form. This may be used to carry a state between calls to the callback functions. Data saved in `ShellState` lives through an entire shell session. Data saved in `ExprState` lives only through the evaluation of the current expression.

There are two ways to start a restricted shell session:

- Use the STDLIB application variable `restricted_shell` and specify, as its value, the name of the callback module. Example (with callback functions implemented in callback_mod.erl): `$ erl -stdlib restricted_shell callback_mod`
- From a normal shell session, call function `shell:start_restricted/1`. This exits the current evaluator and starts a new one in restricted mode.

Notes:

- When restricted shell mode is activated or deactivated, new jobs started on the node will run in restricted or normal mode respectively.
- If restricted mode has been enabled on a particular node, remote shells connecting to this node will also run in restricted mode.
- The callback functions cannot be used to allow or disallow execution of functions called from compiled code (only functions called from expressions entered at the shell prompt).

Errors when loading the callback module is handled in different ways depending on how the restricted shell is activated:

- If the restricted shell is activated by setting the kernel variable during emulator startup and the callback module cannot be loaded, a default restricted shell allowing only the commands `q()` and `init:stop()` is used as fallback.
- If the restricted shell is activated using `shell:start_restricted/1` and the callback module cannot be loaded, an error report is sent to the error logger and the call returns `{error,Reason}`.

Prompting

The default shell prompt function displays the name of the node (if the node can be part of a distributed system) and the current command number. The user can customize the prompt function by calling `shell:prompt_func/1` or by setting the application configuration parameter `shell_prompt_func` for the application STDLIB.

A customized prompt function is stated as a tuple `{Mod, Func}`. The function is called as `Mod:Func(L)`, where `L` is a list of key-value pairs created by the shell. Currently there is only one pair: `{history, N}`, where N is the current command number. The function should return a list of characters or an atom. This constraint is due to the Erlang I/O-protocol. Unicode characters beyond codepoint 255 are allowed in the list. Note that in restricted mode the call `Mod:Func(L)` must be allowed or the default shell prompt function will be called.

Exports

`history(N) -> integer()`

Types:

N = integer()

Sets the number of previous commands to keep in the history list to `N`. The previous number is returned. The default number is 20.

`results(N) -> integer()`

Types:

N = integer()

Sets the number of results from previous commands to keep in the history list to `N`. The previous number is returned. The default number is 20.

`catch_exception(Bool) -> Bool`

Types:

Bool = bool()

Sets the exception handling of the evaluator process. The previous exception handling is returned. The default (`false`) is to kill the evaluator process when an exception occurs, which causes the shell to create a new evaluator process. When the exception handling is set to `true` the evaluator process lives on which means that for instance ports and ETS tables as well as processes linked to the evaluator process survive the exception.

`prompt_func(PromptFunc) -> prompt_func()`

Types:

PromptFunc = prompt_func()

prompt_func() = default | {Mod, Func}

Mod = Func = atom()

Sets the shell prompt function to `PromptFunc`. The previous prompt function is returned.

`start_restricted(Module) -> ok | {error, Reason}`

Types:

Module = atom()

Reason = atom()

Exits a normal shell and starts a restricted shell. `Module` specifies the callback module for the functions `local_allowed/3` and `non_local_allowed/3`. The function is meant to be called from the shell.

If the callback module cannot be loaded, an error tuple is returned. The `Reason` in the error tuple is the one returned by the code loader when trying to load the code of the callback module.

`stop_restricted() -> ok`

Exits a restricted shell and starts a normal shell. The function is meant to be called from the shell.

shell_default

Erlang module

The functions in `shell_default` are called when no module name is given in a shell command.

Consider the following shell dialogue:

```
1 > lists:reverse("abc").
"cba"
2 > c(foo).
{ok, foo}
```

In command one, the module `lists` is called. In command two, no module name is specified. The shell searches the modules `user_default` followed by `shell_default` for the function `foo/1`.

`shell_default` is intended for "system wide" customizations to the shell. `user_default` is intended for "local" or individual user customizations.

Hint

To add your own commands to the shell, create a module called `user_default` and add the commands you want. Then add the following line as the *first* line in your `.erlang` file in your home directory.

```
code:load_abs("$PATH/user_default").
```

`$PATH` is the directory where your `user_default` module can be found.

slave

Erlang module

This module provides functions for starting Erlang slave nodes. All slave nodes which are started by a master will terminate automatically when the master terminates. All TTY output produced at the slave will be sent back to the master node. File I/O is done via the master.

Slave nodes on other hosts than the current one are started with the program `rsh`. The user must be allowed to `rsh` to the remote hosts without being prompted for a password. This can be arranged in a number of ways (refer to the `rsh` documentation for details). A slave node started on the same host as the master inherits certain environment values from the master, such as the current directory and the environment variables. For what can be assumed about the environment when a slave is started on another host, read the documentation for the `rsh` program.

An alternative to the `rsh` program can be specified on the command line to `erl` as follows: `-rsh Program`.

The slave node should use the same file system at the master. At least, Erlang/OTP should be installed in the same place on both computers and the same version of Erlang should be used.

Currently, a node running on Windows NT can only start slave nodes on the host on which it is running.

The master node must be alive.

Exports

`start(Host) ->`

`start(Host, Name) ->`

`start(Host, Name, Args) -> {ok, Node} | {error, Reason}`

Types:

Host = Name = atom()

Args = string()

Node = node()

Reason = timeout | no_rsh | {already_running, Node}

Starts a slave node on the host `Host`. Host names need not necessarily be specified as fully qualified names; short names can also be used. This is the same condition that applies to names of distributed Erlang nodes.

The name of the started node will be `Name@Host`. If no name is provided, the name will be the same as the node which executes the call (with the exception of the host name part of the node name).

The slave node resets its `user` process so that all terminal I/O which is produced at the slave is automatically relayed to the master. Also, the file process will be relayed to the master.

The `Args` argument is used to set `erl` command line arguments. If provided, it is passed to the new node and can be used for a variety of purposes. See *erl(1)*

As an example, suppose that we want to start a slave node at host `H` with the node name `Name@H`, and we also want the slave node to have the following properties:

- directory `Dir` should be added to the code path;
- the Mnesia directory should be set to `M`;
- the unix `DISPLAY` environment variable should be set to the display of the master node.

The following code is executed to achieve this:

```
E = " -env DISPLAY " ++ net_adm:localhost() ++ ":0 ",
Arg = "-mnesia_dir " ++ M ++ " -pa " ++ Dir ++ E,
slave:start(H, Name, Arg).
```

If successful, the function returns `{ok, Node}`, where `Node` is the name of the new node. Otherwise it returns `{error, Reason}`, where `Reason` can be one of:

`timeout`

The master node failed to get in contact with the slave node. This can happen in a number of circumstances:

- Erlang/OTP is not installed on the remote host
- the file system on the other host has a different structure to the the master
- the Erlang nodes have different cookies.

`no_rsh`

There is no `rsh` program on the computer.

`{already_running, Node}`

A node with the name `Name@Host` already exists.

`start_link(Host) ->`

`start_link(Host, Name) ->`

`start_link(Host, Name, Args) -> {ok, Node} | {error, Reason}`

Types:

Host = Name = atom()

Args = string()

Node = node()

Reason = timeout | no_rsh | {already_running, Node}

Starts a slave node in the same way as `start/1,2,3`, except that the slave node is linked to the currently executing process. If that process terminates, the slave node also terminates.

See `start/1,2,3` for a description of arguments and return values.

`stop(Node) -> ok`

Types:

Node = node()

Stops (kills) a node.

`pseudo([Master | ServerList]) -> ok`

Types:

Master = node()

ServerList = [atom()]

Calls `pseudo(Master, ServerList)`. If we want to start a node from the command line and set up a number of pseudo servers, an Erlang runtime system can be started as follows:

```
% erl -name abc -s slave pseudo klacke@super x --
```

pseudo(Master, ServerList) -> ok

Types:

Master = node()

ServerList = [atom()]

Starts a number of pseudo servers. A pseudo server is a server with a registered name which does absolutely nothing but pass on all message to the real server which executes at a master node. A pseudo server is an intermediary which only has the same registered name as the real server.

For example, if we have started a slave node `N` and want to execute `pxw` graphics code on this node, we can start the server `pxw_server` as a pseudo server at the slave node. The following code illustrates:

```
rpc:call(N, slave, pseudo, [node(), [pxw_server]]).
```

relay(Pid)

Types:

Pid = pid()

Runs a pseudo server. This function never returns any value and the process which executes the function will receive messages. All messages received will simply be passed on to `Pid`.

sofs

Erlang module

The `sofs` module implements operations on finite sets and relations represented as sets. Intuitively, a set is a collection of elements; every element belongs to the set, and the set contains every element.

Given a set A and a sentence S(x), where x is a free variable, a new set B whose elements are exactly those elements of A for which S(x) holds can be formed, this is denoted B = {x in A : S(x)}. Sentences are expressed using the logical operators "for some" (or "there exists"), "for all", "and", "or", "not". If the existence of a set containing all the specified elements is known (as will always be the case in this module), we write B = {x : S(x)}.

The *unordered set* containing the elements a, b and c is denoted {a, b, c}. This notation is not to be confused with tuples. The *ordered pair* of a and b, with first *coordinate* a and second coordinate b, is denoted (a, b). An ordered pair is an *ordered set* of two elements. In this module ordered sets can contain one, two or more elements, and parentheses are used to enclose the elements. Unordered sets and ordered sets are orthogonal, again in this module; there is no unordered set equal to any ordered set.

The set that contains no elements is called the *empty set*. If two sets A and B contain the same elements, then A is *equal* to B, denoted A = B. Two ordered sets are equal if they contain the same number of elements and have equal elements at each coordinate. If a set A contains all elements that B contains, then B is a *subset* of A. The *union* of two sets A and B is the smallest set that contains all elements of A and all elements of B. The *intersection* of two sets A and B is the set that contains all elements of A that belong to B. Two sets are *disjoint* if their intersection is the empty set. The *difference* of two sets A and B is the set that contains all elements of A that do not belong to B. The *symmetric difference* of two sets is the set that contains those element that belong to either of the two sets, but not both. The *union* of a collection of sets is the smallest set that contains all the elements that belong to at least one set of the collection. The *intersection* of a non-empty collection of sets is the set that contains all elements that belong to every set of the collection.

The *Cartesian product* of two sets X and Y, denoted X × Y, is the set {a : a = (x, y) for some x in X and for some y in Y}. A *relation* is a subset of X × Y. Let R be a relation. The fact that (x, y) belongs to R is written as x R y. Since relations are sets, the definitions of the last paragraph (subset, union, and so on) apply to relations as well. The *domain* of R is the set {x : x R y for some y in Y}. The *range* of R is the set {y : x R y for some x in X}. The *converse* of R is the set {a : a = (y, x) for some (x, y) in R}. If A is a subset of X, then the *image* of A under R is the set {y : x R y for some x in A}, and if B is a subset of Y, then the *inverse image* of B is the set {x : x R y for some y in B}. If R is a relation from X to Y and S is a relation from Y to Z, then the *relative product* of R and S is the relation T from X to Z defined so that x T z if and only if there exists an element y in Y such that x R y and y S z. The *restriction* of R to A is the set S defined so that x S y if and only if there exists an element x in A such that x R y. If S is a restriction of R to A, then R is an *extension* of S to X. If X = Y then we call R a relation *in* X. The *field* of a relation R in X is the union of the domain of R and the range of R. If R is a relation in X, and if S is defined so that x S y if x R y and not x = y, then S is the *strict* relation corresponding to R, and vice versa, if S is a relation in X, and if R is defined so that x R y if x S y or x = y, then R is the *weak* relation corresponding to S. A relation R in X is *reflexive* if x R x for every element x of X; it is *symmetric* if x R y implies that y R x; and it is *transitive* if x R y and y R z imply that x R z.

A *function* F is a relation, a subset of X × Y, such that the domain of F is equal to X and such that for every x in X there is a unique element y in Y with (x, y) in F. The latter condition can be formulated as follows: if x F y and x F z then y = z. In this module, it will not be required that the domain of F be equal to X for a relation to be considered a function. Instead of writing (x, y) in F or x F y, we write F(x) = y when F is a function, and say that F maps x onto y, or that the value of F at x is y. Since functions are relations, the definitions of the last paragraph (domain, range, and so on) apply to functions as well. If the converse of a function F is a function F', then F' is called the *inverse* of F. The relative product of two functions F1 and F2 is called the *composite* of F1 and F2 if the range of F1 is a subset of the domain of F2.

Sometimes, when the range of a function is more important than the function itself, the function is called a *family*. The domain of a family is called the *index set*, and the range is called the *indexed set*. If x is a family from I to X, then x[i] denotes the value of the function at index i. The notation "a family in X" is used for such a family. When the indexed set is a set of subsets of a set X, then we call x a *family of subsets* of X. If x is a family of subsets of X, then the union of the range of x is called the *union of the family* x. If x is non-empty (the index set is non-empty), the *intersection of the family* x is the intersection of the range of x. In this module, the only families that will be considered are families of subsets of some set X; in the following the word "family" will be used for such families of subsets.

A *partition* of a set X is a collection S of non-empty subsets of X whose union is X and whose elements are pairwise disjoint. A relation in a set is an *equivalence relation* if it is reflexive, symmetric and transitive. If R is an equivalence relation in X, and x is an element of X, the *equivalence class* of x with respect to R is the set of all those elements y of X for which x R y holds. The equivalence classes constitute a partitioning of X. Conversely, if C is a partition of X, then the relation that holds for any two elements of X if they belong to the same equivalence class, is an equivalence relation induced by the partition C. If R is an equivalence relation in X, then the *canonical map* is the function that maps every element of X onto its equivalence class.

Relations as defined above (as sets of ordered pairs) will from now on be referred to as *binary relations*. We call a set of ordered sets (x[1], ..., x[n]) an *(n-ary) relation*, and say that the relation is a subset of the Cartesian product X[1] × ... × X[n] where x[i] is an element of X[i], 1 <= i <= n. The *projection* of an n-ary relation R onto coordinate i is the set {x[i] : (x[1], ..., x[i], ..., x[n]) in R for some x[j] in X[j], 1 <= j <= n and not i = j}. The projections of a binary relation R onto the first and second coordinates are the domain and the range of R respectively. The relative product of binary relations can be generalized to n-ary relations as follows. Let TR be an ordered set (R[1], ..., R[n]) of binary relations from X to Y[i] and S a binary relation from (Y[1] × ... × Y[n]) to Z. The *relative product* of TR and S is the binary relation T from X to Z defined so that x T z if and only if there exists an element y[i] in Y[i] for each 1 <= i <= n such that x R[i] y[i] and (y[1], ..., y[n]) S z. Now let TR be a an ordered set (R[1], ..., R[n]) of binary relations from X[i] to Y[i] and S a subset of X[1] × ... × X[n]. The *multiple relative product* of TR and S is defined to be the set {z : z = ((x[1], ..., x[n]), (y[1],...,y[n])) for some (x[1], ..., x[n]) in S and for some (x[i], y[i]) in R[i], 1 <= i <= n}. The *natural join* of an n-ary relation R and an m-ary relation S on coordinate i and j is defined to be the set {z : z = (x[1], ..., x[n], y[1], ..., y[j-1], y[j+1], ..., y[m]) for some (x[1], ..., x[n]) in R and for some (y[1], ..., y[m]) in S such that x[i] = y[j]}.

The sets recognized by this module will be represented by elements of the relation Sets, defined as the smallest set such that:

- for every atom T except '_' and for every term X, (T, X) belongs to Sets (*atomic sets*);
- (['_'], []) belongs to Sets (the *untyped empty set*);
- for every tuple T = {T[1], ..., T[n]} and for every tuple X = {X[1], ..., X[n]}, if (T[i], X[i]) belongs to Sets for every 1 <= i <= n then (T, X) belongs to Sets (*ordered sets*);
- for every term T, if X is the empty list or a non-empty sorted list [X[1], ..., X[n]] without duplicates such that (T, X[i]) belongs to Sets for every 1 <= i <= n, then ([T], X) belongs to Sets (*typed unordered sets*).

An *external set* is an element of the range of Sets. A *type* is an element of the domain of Sets. If S is an element (T, X) of Sets, then T is a *valid type* of X, T is the type of S, and X is the external set of S. *from_term/2* creates a set from a type and an Erlang term turned into an external set.

The actual sets represented by Sets are the elements of the range of the function Set from Sets to Erlang terms and sets of Erlang terms:

- Set(T,Term) = Term, where T is an atom;
- Set({T[1], ..., T[n]}, {X[1], ..., X[n]}) = (Set(T[1], X[1]), ..., Set(T[n], X[n]));
- Set([T], [X[1], ..., X[n]]) = {Set(T, X[1]), ..., Set(T, X[n])};
- Set([T], []) = {}.

When there is no risk of confusion, elements of Sets will be identified with the sets they represent. For instance, if U is the result of calling `union/2` with S1 and S2 as arguments, then U is said to be the union of S1 and S2. A more precise formulation would be that Set(U) is the union of Set(S1) and Set(S2).

The types are used to implement the various conditions that sets need to fulfill. As an example, consider the relative product of two sets R and S, and recall that the relative product of R and S is defined if R is a binary relation to Y and S is a binary relation from Y. The function that implements the relative product, *relative_product/2*, checks that the arguments represent binary relations by matching [{A,B}] against the type of the first argument (Arg1 say), and [{C,D}] against the type of the second argument (Arg2 say). The fact that [{A,B}] matches the type of Arg1 is to be interpreted as Arg1 representing a binary relation from X to Y, where X is defined as all sets Set(x) for some element x in Sets the type of which is A, and similarly for Y. In the same way Arg2 is interpreted as representing a binary relation from W to Z. Finally it is checked that B matches C, which is sufficient to ensure that W is equal to Y. The untyped empty set is handled separately: its type, ['_'], matches the type of any unordered set.

A few functions of this module (`drestriction/3`, `family_projection/2`, `partition/2`, `partition_family/2`, `projection/2`, `restriction/3`, `substitution/2`) accept an Erlang function as a means to modify each element of a given unordered set. Such a function, called SetFun in the following, can be specified as a functional object (fun), a tuple `{external, Fun}`, or an integer. If SetFun is specified as a fun, the fun is applied to each element of the given set and the return value is assumed to be a set. If SetFun is specified as a tuple `{external, Fun}`, Fun is applied to the external set of each element of the given set and the return value is assumed to be an external set. Selecting the elements of an unordered set as external sets and assembling a new unordered set from a list of external sets is in the present implementation more efficient than modifying each element as a set. However, this optimization can only be utilized when the elements of the unordered set are atomic or ordered sets. It must also be the case that the type of the elements matches some clause of Fun (the type of the created set is the result of applying Fun to the type of the given set), and that Fun does nothing but selecting, duplicating or rearranging parts of the elements. Specifying a SetFun as an integer I is equivalent to specifying `{external, fun(X) -> element(I, X) end}`, but is to be preferred since it makes it possible to handle this case even more efficiently. Examples of SetFuns:

```
{sofs, union}
fun(S) -> sofs:partition(1, S) end
{external, fun(A) -> A end}
{external, fun({A,_,C}) -> {C,A} end}
{external, fun({_,{_,C}}) -> C end}
{external, fun({_,{_,{_,E}=C}}) -> {E,{E,C}} end}
2
```

The order in which a SetFun is applied to the elements of an unordered set is not specified, and may change in future versions of sofs.

The execution time of the functions of this module is dominated by the time it takes to sort lists. When no sorting is needed, the execution time is in the worst case proportional to the sum of the sizes of the input arguments and the returned value. A few functions execute in constant time: `from_external`, `is_empty_set`, `is_set`, `is_sofs_set`, `to_external`, `type`.

The functions of this module exit the process with a `badarg`, `bad_function`, or `type_mismatch` message when given badly formed arguments or sets the types of which are not compatible.

When comparing external sets the operator `==/2` is used.

Types

```
anyset() = - an unordered, ordered or atomic set -
binary_relation() = - a binary relation -
bool() = true | false
external_set() = - an external set -
family() = - a family (of subsets) -
function() = - a function -
ordset() = - an ordered set -
```

```
relation() = - an n-ary relation -
set() = - an unordered set -
set_of_sets() = - an unordered set of set() -
set_fun() = integer() >= 1
          | {external, fun(external_set()) -> external_set()}
          | fun(anyset()) -> anyset()
spec_fun() = {external, fun(external_set()) -> bool()}
           | fun(anyset()) -> bool()
type() = - a type -
```

Exports

a_function(Tuples [, Type]) -> Function

Types:

Function = function()

Tuples = [tuple()]

Type = type()

Creates a *function*. `a_function(F, T)` is equivalent to `from_term(F, T)`, if the result is a function. If no *type* is explicitly given, `[{atom, atom}]` is used as type of the function.

canonical_relation(SetOfSets) -> BinRel

Types:

BinRel = binary_relation()

SetOfSets = set_of_sets()

Returns the binary relation containing the elements (E, Set) such that Set belongs to SetOfSets and E belongs to Set. If SetOfSets is a *partition* of a set X and R is the equivalence relation in X induced by SetOfSets, then the returned relation is the *canonical map* from X onto the equivalence classes with respect to R.

```
1> Ss = sofs:from_term([[a,b],[b,c]]),
CR = sofs:canonical_relation(Ss),
sofs:to_external(CR).
[{a,[a,b]},{b,[a,b]},{b,[b,c]},{c,[b,c]}]
```

composite(Function1, Function2) -> Function3

Types:

Function1 = Function2 = Function3 = function()

Returns the *composite* of the functions Function1 and Function2.

```
1> F1 = sofs:a_function([{a,1},{b,2},{c,2}]),
F2 = sofs:a_function([{1,x},{2,y},{3,z}]),
F = sofs:composite(F1, F2),
sofs:to_external(F).
[{a,x},{b,y},{c,y}]
```

constant_function(Set, AnySet) -> Function

Types:

AnySet = anyset()

Function = function()

Set = set()

Creates the *function* that maps each element of the set Set onto AnySet.

```
1> S = sofs:set([a,b]),
E = sofs:from_term(1),
R = sofs:constant_function(S, E),
sofs:to_external(R).
[{a,1},{b,1}]
```

converse(BinRel1) -> BinRel2

Types:

BinRel1 = BinRel2 = binary_relation()

Returns the *converse* of the binary relation BinRel1.

```
1> R1 = sofs:relation([{1,a},{2,b},{3,a}]),
R2 = sofs:converse(R1),
sofs:to_external(R2).
[{a,1},{a,3},{b,2}]
```

difference(Set1, Set2) -> Set3

Types:

Set1 = Set2 = Set3 = set()

Returns the *difference* of the sets Set1 and Set2.

digraph_to_family(Graph [, Type]) -> Family

Types:

Graph = digraph() - see digraph(3) -

Family = family()

Type = type()

Creates a *family* from the directed graph Graph. Each vertex a of Graph is represented by a pair (a, {b[1], ..., b[n]}) where the b[i]'s are the out-neighbours of a. If no type is explicitly given, [{atom, [atom]}] is used as type of the family. It is assumed that Type is a *valid type* of the external set of the family.

If G is a directed graph, it holds that the vertices and edges of G are the same as the vertices and edges of `family_to_digraph(digraph_to_family(G))`.

domain(BinRel) -> Set

Types:

BinRel = binary_relation()

Set = set()

Returns the *domain* of the binary relation BinRel.

```
1> R = sofs:relation([{1,a},{1,b},{2,b},{2,c}]),
S = sofs:domain(R),
sofs:to_external(S).
[1,2]
```

drestriction(BinRel1, Set) -> BinRel2

Types:

BinRel1 = **BinRel2** = **binary_relation()**

Set = **set()**

Returns the difference between the binary relation BinRel1 and the *restriction* of BinRel1 to Set.

```
1> R1 = sofs:relation([{1,a},{2,b},{3,c}]),
S = sofs:set([2,4,6]),
R2 = sofs:drestriction(R1, S),
sofs:to_external(R2).
[{1,a},{3,c}]
```

`drestriction(R, S)` is equivalent to `difference(R, restriction(R, S))`.

drestriction(SetFun, Set1, Set2) -> Set3

Types:

SetFun = **set_fun()**

Set1 = **Set2** = **Set3** = **set()**

Returns a subset of Set1 containing those elements that do not yield an element in Set2 as the result of applying SetFun.

```
1> SetFun = {external, fun({_A,B,C}) -> {B,C} end},
R1 = sofs:relation([{a,aa,1},{b,bb,2},{c,cc,3}]),
R2 = sofs:relation([{bb,2},{cc,3},{dd,4}]),
R3 = sofs:drestriction(SetFun, R1, R2),
sofs:to_external(R3).
[{a,aa,1}]
```

`drestriction(F, S1, S2)` is equivalent to `difference(S1, restriction(F, S1, S2))`.

empty_set() -> Set

Types:

Set = **set()**

Returns the *untyped empty set*. `empty_set()` is equivalent to `from_term([], ['_'])`.

extension(BinRel1, Set, AnySet) -> BinRel2

Types:

AnySet = **anyset()**

BinRel1 = **BinRel2** = **binary_relation()**

Set = **set()**

Returns the *extension* of BinRel1 such that for each element E in Set that does not belong to the *domain* of BinRel1, BinRel2 contains the pair (E, AnySet).

```
1> S = sofs:set([b,c]),
A = sofs:empty_set(),
R = sofs:family([{a,[1,2]},{b,[3]}]),
X = sofs:extension(R, S, A),
sofs:to_external(X).
[{a,[1,2]},{b,[3]},{c,[]}]
```

family(Tuples [, Type]) -> Family

Types:

Family = **family()**

Tuples = **[tuple()]**

Type = **type()**

Creates a *family of subsets*. `family(F, T)` is equivalent to `from_term(F, T)`, if the result is a family. If no *type* is explicitly given, `[{atom, [atom]}]` is used as type of the family.

family_difference(Family1, Family2) -> Family3

Types:

Family1 = **Family2** = **Family3** = **family()**

If Family1 and Family2 are *families*, then Family3 is the family such that the index set is equal to the index set of Family1, and Family3[i] is the difference between Family1[i] and Family2[i] if Family2 maps i, Family1[i] otherwise.

```
1> F1 = sofs:family([{a,[1,2]},{b,[3,4]}]),
F2 = sofs:family([{b,[4,5]},{c,[6,7]}]),
F3 = sofs:family_difference(F1, F2),
sofs:to_external(F3).
[{a,[1,2]},{b,[3]}]
```

family_domain(Family1) -> Family2

Types:

Family1 = **Family2** = **family()**

If Family1 is a *family* and Family1[i] is a binary relation for every i in the index set of Family1, then Family2 is the family with the same index set as Family1 such that Family2[i] is the *domain* of Family1[i].

```
1> FR = sofs:from_term([{a,[{1,a},{2,b},{3,c}]},{b,[]},{c,[{4,d},{5,e}]}]),
F = sofs:family_domain(FR),
sofs:to_external(F).
[{a,[1,2,3]},{b,[]},{c,[4,5]}]
```

family_field(Family1) -> Family2

Types:

Family1 = **Family2** = **family()**

If Family1 is a *family* and Family1[i] is a binary relation for every i in the index set of Family1, then Family2 is the family with the same index set as Family1 such that Family2[i] is the *field* of Family1[i].

```
1> FR = sofs:from_term([{a,[{1,a},{2,b},{3,c}]},{b,[]},{c,[{4,d},{5,e}]}]),
F = sofs:family_field(FR),
sofs:to_external(F).
[{a,[1,2,3,a,b,c]},{b,[]},{c,[4,5,d,e]}]
```

`family_field(Family1)` is equivalent to `family_union(family_domain(Family1), family_range(Family1))`.

`family_intersection(Family1) -> Family2`

Types:

Family1 = Family2 = family()

If Family1 is a *family* and Family1[i] is a set of sets for every i in the index set of Family1, then Family2 is the family with the same index set as Family1 such that Family2[i] is the *intersection* of Family1[i].

If Family1[i] is an empty set for some i, then the process exits with a `badarg` message.

```
1> F1 = sofs:from_term([{a,[[1,2,3],[2,3,4]]},{b,[[x,y,z],[x,y]]}]),
F2 = sofs:family_intersection(F1),
sofs:to_external(F2).
[{a,[2,3]},{b,[x,y]}]
```

`family_intersection(Family1, Family2) -> Family3`

Types:

Family1 = Family2 = Family3 = family()

If Family1 and Family2 are *families*, then Family3 is the family such that the index set is the intersection of Family1's and Family2's index sets, and Family3[i] is the intersection of Family1[i] and Family2[i].

```
1> F1 = sofs:family([{a,[1,2]},{b,[3,4]},{c,[5,6]}]),
F2 = sofs:family([{b,[4,5]},{c,[7,8]},{d,[9,10]}]),
F3 = sofs:family_intersection(F1, F2),
sofs:to_external(F3).
[{b,[4]},{c,[]}]
```

`family_projection(SetFun, Family1) -> Family2`

Types:

SetFun = set_fun()

Family1 = Family2 = family()

Set = set()

If Family1 is a *family* then Family2 is the family with the same index set as Family1 such that Family2[i] is the result of calling SetFun with Family1[i] as argument.

```
1> F1 = sofs:from_term([{a,[[1,2],[2,3]]},{b,[[]]}]),
F2 = sofs:family_projection({sofs, union}, F1),
sofs:to_external(F2).
[{a,[1,2,3]},{b,[]}]
```

family_range(Family1) -> Family2

Types:

Family1 = Family2 = family()

If Family1 is a *family* and Family1[i] is a binary relation for every i in the index set of Family1, then Family2 is the family with the same index set as Family1 such that Family2[i] is the *range* of Family1[i].

```
1> FR = sofs:from_term([{a,[{1,a},{2,b},{3,c}]},{b,[]},{c,[{4,d},{5,e}]}]),
F = sofs:family_range(FR),
sofs:to_external(F).
[{a,[a,b,c]},{b,[]},{c,[d,e]}]
```

family_specification(Fun, Family1) -> Family2

Types:

Fun = spec_fun()

Family1 = Family2 = family()

If Family1 is a *family*, then Family2 is the *restriction* of Family1 to those elements i of the index set for which Fun applied to Family1[i] returns `true`. If Fun is a tuple `{external, Fun2}`, Fun2 is applied to the *external set* of Family1[i], otherwise Fun is applied to Family1[i].

```
1> F1 = sofs:family([{a,[1,2,3]},{b,[1,2]},{c,[1]}]),
SpecFun = fun(S) -> sofs:no_elements(S) =:= 2 end,
F2 = sofs:family_specification(SpecFun, F1),
sofs:to_external(F2).
[{b,[1,2]}]
```

family_to_digraph(Family [, GraphType]) -> Graph

Types:

Graph = digraph()

Family = family()

GraphType = - see digraph(3) -

Creates a directed graph from the *family* Family. For each pair (a, {b[1], ..., b[n]}) of Family, the vertex a as well the edges (a, b[i]) for 1 <= i <= n are added to a newly created directed graph.

If no graph type is given, `digraph:new/1` is used for creating the directed graph, otherwise the GraphType argument is passed on as second argument to `digraph:new/2`.

It F is a family, it holds that F is a subset of `digraph_to_family(family_to_digraph(F), type(F))`. Equality holds if `union_of_family(F)` is a subset of `domain(F)`.

Creating a cycle in an acyclic graph exits the process with a `cyclic` message.

family_to_relation(Family) -> BinRel

Types:

Family = family()

BinRel = binary_relation()

If Family is a *family*, then BinRel is the binary relation containing all pairs (i, x) such that i belongs to the index set of Family and x belongs to Family[i].

```
1> F = sofs:family([{a,[]}, {b,[1]}, {c,[2,3]}]),
R = sofs:family_to_relation(F),
sofs:to_external(R).
[{b,1},{c,2},{c,3}]
```

family_union(Family1) -> Family2

Types:

Family1 = Family2 = family()

If Family1 is a *family* and Family1[i] is a set of sets for each i in the index set of Family1, then Family2 is the family with the same index set as Family1 such that Family2[i] is the *union* of Family1[i].

```
1> F1 = sofs:from_term([{a,[[1,2],[2,3]]},{b,[[]]}]),
F2 = sofs:family_union(F1),
sofs:to_external(F2).
[{a,[1,2,3]},{b,[]}]
```

`family_union(F)` is equivalent to `family_projection({sofs,union}, F)`.

family_union(Family1, Family2) -> Family3

Types:

Family1 = Family2 = Family3 = family()

If Family1 and Family2 are *families*, then Family3 is the family such that the index set is the union of Family1's and Family2's index sets, and Family3[i] is the union of Family1[i] and Family2[i] if both maps i, Family1[i] or Family2[i] otherwise.

```
1> F1 = sofs:family([{a,[1,2]},{b,[3,4]},{c,[5,6]}]),
F2 = sofs:family([{b,[4,5]},{c,[7,8]},{d,[9,10]}]),
F3 = sofs:family_union(F1, F2),
sofs:to_external(F3).
[{a,[1,2]},{b,[3,4,5]},{c,[5,6,7,8]},{d,[9,10]}]
```

field(BinRel) -> Set

Types:

BinRel = binary_relation()

Set = set()

Returns the *field* of the binary relation BinRel.

```
1> R = sofs:relation([{1,a},{1,b},{2,b},{2,c}]),
S = sofs:field(R),
sofs:to_external(S).
[1,2,a,b,c]
```

`field(R)` is equivalent to `union(domain(R), range(R))`.

from_external(ExternalSet, Type) -> AnySet

Types:

ExternalSet = external_set()

AnySet = anyset()

Type = type()

Creates a set from the *external set* ExternalSet and the *type* Type. It is assumed that Type is a *valid type* of ExternalSet.

from_sets(ListOfSets) -> Set

Types:

Set = set()

ListOfSets = [anyset()]

Returns the *unordered set* containing the sets of the list ListOfSets.

```
1> S1 = sofs:relation([{a,1},{b,2}]),
S2 = sofs:relation([{x,3},{y,4}]),
S = sofs:from_sets([S1,S2]),
sofs:to_external(S).
[[{a,1},{b,2}],[{x,3},{y,4}]]
```

from_sets(TupleOfSets) -> Ordset

Types:

Ordset = ordset()

TupleOfSets = tuple-of(anyset())

Returns the *ordered set* containing the sets of the non-empty tuple TupleOfSets.

from_term(Term [, Type]) -> AnySet

Types:

AnySet = anyset()

Term = term()

Type = type()

Creates an element of *Sets* by traversing the term Term, sorting lists, removing duplicates and deriving or verifying a *valid type* for the so obtained external set. An explicitly given *type* Type can be used to limit the depth of the traversal; an atomic type stops the traversal, as demonstrated by this example where "foo" and {"foo"} are left unmodified:

```
1> S = sofs:from_term([{{"foo"},[1,1]},{"foo",[2,2]}], [{atom,[atom]}]),
sofs:to_external(S).
[{{"foo"},[1]},{"foo",[2]}]
```

`from_term` can be used for creating atomic or ordered sets. The only purpose of such a set is that of later building unordered sets since all functions in this module that *do* anything operate on unordered sets. Creating unordered sets from a collection of ordered sets may be the way to go if the ordered sets are big and one does not want to waste heap by rebuilding the elements of the unordered set. An example showing that a set can be built "layer by layer":

```
1> A = sofs:from_term(a),
```

```
S = sofs:set([1,2,3]),
P1 = sofs:from_sets({A,S}),
P2 = sofs:from_term({b,[6,5,4]}),
Ss = sofs:from_sets([P1,P2]),
sofs:to_external(Ss).
[{a,[1,2,3]},{b,[4,5,6]}]
```

Other functions that create sets are `from_external/2` and `from_sets/1`. Special cases of `from_term/2` are `a_function/1,2`, `empty_set/0`, `family/1,2`, `relation/1,2`, and `set/1,2`.

image(BinRel, Set1) -> Set2

Types:

BinRel = binary_relation()

Set1 = Set2 = set()

Returns the *image* of the set Set1 under the binary relation BinRel.

```
1> R = sofs:relation([{1,a},{2,b},{2,c},{3,d}]),
S1 = sofs:set([1,2]),
S2 = sofs:image(R, S1),
sofs:to_external(S2).
[a,b,c]
```

intersection(SetOfSets) -> Set

Types:

Set = set()

SetOfSets = set_of_sets()

Returns the *intersection* of the set of sets SetOfSets.

Intersecting an empty set of sets exits the process with a `badarg` message.

intersection(Set1, Set2) -> Set3

Types:

Set1 = Set2 = Set3 = set()

Returns the *intersection* of Set1 and Set2.

intersection_of_family(Family) -> Set

Types:

Family = family()

Set = set()

Returns the intersection of the *family* Family.

Intersecting an empty family exits the process with a `badarg` message.

```
1> F = sofs:family([{a,[0,2,4]},{b,[0,1,2]},{c,[2,3]}]),
S = sofs:intersection_of_family(F),
sofs:to_external(S).
[2]
```

inverse(Function1) -> Function2

Types:

Function1 = Function2 = function()

Returns the *inverse* of the function Function1.

```
1> R1 = sofs:relation([{1,a},{2,b},{3,c}]),
R2 = sofs:inverse(R1),
sofs:to_external(R2).
[{a,1},{b,2},{c,3}]
```

inverse_image(BinRel, Set1) -> Set2

Types:

BinRel = binary_relation()

Set1 = Set2 = set()

Returns the *inverse image* of Set1 under the binary relation BinRel.

```
1> R = sofs:relation([{1,a},{2,b},{2,c},{3,d}]),
S1 = sofs:set([c,d,e]),
S2 = sofs:inverse_image(R, S1),
sofs:to_external(S2).
[2,3]
```

is_a_function(BinRel) -> Bool

Types:

Bool = bool()

BinRel = binary_relation()

Returns `true` if the binary relation BinRel is a *function* or the untyped empty set, `false` otherwise.

is_disjoint(Set1, Set2) -> Bool

Types:

Bool = bool()

Set1 = Set2 = set()

Returns `true` if Set1 and Set2 are *disjoint*, `false` otherwise.

is_empty_set(AnySet) -> Bool

Types:

AnySet = anyset()

Bool = bool()

Returns `true` if Set is an empty unordered set, `false` otherwise.

is_equal(AnySet1, AnySet2) -> Bool

Types:

AnySet1 = AnySet2 = anyset()

Bool = bool()

Returns `true` if the AnySet1 and AnySet2 are *equal*, `false` otherwise. This example shows that `==/2` is used when comparing sets for equality:

```
1> S1 = sofs:set([1.0]),
S2 = sofs:set([1]),
sofs:is_equal(S1, S2).
true
```

is_set(AnySet) -> Bool

Types:

AnySet = anyset()

Bool = bool()

Returns `true` if AnySet is an *unordered set*, and `false` if AnySet is an ordered set or an atomic set.

is_sofs_set(Term) -> Bool

Types:

Bool = bool()

Term = term()

Returns `true` if Term is an *unordered set*, an ordered set or an atomic set, `false` otherwise.

is_subset(Set1, Set2) -> Bool

Types:

Bool = bool()

Set1 = Set2 = set()

Returns `true` if Set1 is a *subset* of Set2, `false` otherwise.

is_type(Term) -> Bool

Types:

Bool = bool()

Term = term()

Returns `true` if the term Term is a *type*.

join(Relation1, I, Relation2, J) -> Relation3

Types:

Relation1 = Relation2 = Relation3 = relation()

I = J = integer() > 0

Returns the *natural join* of the relations Relation1 and Relation2 on coordinates I and J.

```
1> R1 = sofs:relation([{a,x,1},{b,y,2}]),
R2 = sofs:relation([{1,f,g},{1,h,i},{2,3,4}]),
J = sofs:join(R1, 3, R2, 1),
sofs:to_external(J).
```

```
[{a,x,1,f,g},{a,x,1,h,i},{b,y,2,3,4}]
```

multiple_relative_product(TupleOfBinRels, BinRel1) -> BinRel2

Types:

TupleOfBinRels = tuple-of(BinRel)

BinRel = BinRel1 = BinRel2 = binary_relation()

If TupleOfBinRels is a non-empty tuple {R[1], ..., R[n]} of binary relations and BinRel1 is a binary relation, then BinRel2 is the *multiple relative product* of the ordered set (R[i], ..., R[n]) and BinRel1.

```
1> Ri = sofs:relation([{a,1},{b,2},{c,3}]),
R = sofs:relation([{a,b},{b,c},{c,a}]),
MP = sofs:multiple_relative_product({Ri, Ri}, R),
sofs:to_external(sofs:range(MP)).
[{1,2},{2,3},{3,1}]
```

no_elements(ASet) -> NoElements

Types:

ASet = set() | ordset()

NoElements = integer() >= 0

Returns the number of elements of the ordered or unordered set ASet.

partition(SetOfSets) -> Partition

Types:

SetOfSets = set_of_sets()

Partition = set()

Returns the *partition* of the union of the set of sets SetOfSets such that two elements are considered equal if they belong to the same elements of SetOfSets.

```
1> Sets1 = sofs:from_term([[a,b,c],[d,e,f],[g,h,i]]),
Sets2 = sofs:from_term([[b,c,d],[e,f,g],[h,i,j]]),
P = sofs:partition(sofs:union(Sets1, Sets2)),
sofs:to_external(P).
[[a],[b,c],[d],[e,f],[g],[h,i],[j]]
```

partition(SetFun, Set) -> Partition

Types:

SetFun = set_fun()

Partition = set()

Set = set()

Returns the *partition* of Set such that two elements are considered equal if the results of applying SetFun are equal.

```
1> Ss = sofs:from_term([[a],[b],[c,d],[e,f]]),
SetFun = fun(S) -> sofs:from_term(sofs:no_elements(S)) end,
P = sofs:partition(SetFun, Ss),
```

```
sofs:to_external(P).
[[[a],[b]],[[c,d],[e,f]]]
```

partition(SetFun, Set1, Set2) -> {Set3, Set4}

Types:

SetFun = set_fun()

Set1 = Set2 = Set3 = Set4 = set()

Returns a pair of sets that, regarded as constituting a set, forms a *partition* of Set1. If the result of applying SetFun to an element of Set1 yields an element in Set2, the element belongs to Set3, otherwise the element belongs to Set4.

```
1> R1 = sofs:relation([{1,a},{2,b},{3,c}]),
S = sofs:set([2,4,6]),
{R2,R3} = sofs:partition(1, R1, S),
{sofs:to_external(R2),sofs:to_external(R3)}.
{[{2,b}],[{1,a},{3,c}]}
```

`partition(F, S1, S2)` is equivalent to `{restriction(F, S1, S2), drestriction(F, S1, S2)}`.

partition_family(SetFun, Set) -> Family

Types:

Family = family()

SetFun = set_fun()

Set = set()

Returns the *family* Family where the indexed set is a *partition* of Set such that two elements are considered equal if the results of applying SetFun are the same value i. This i is the index that Family maps onto the *equivalence class*.

```
1> S = sofs:relation([{a,a,a,a},{a,a,b,b},{a,b,b,b}]),
SetFun = {external, fun({A,_,C,_}) -> {A,C} end},
F = sofs:partition_family(SetFun, S),
sofs:to_external(F).
[{{a,a},[{a,a,a,a}]},{{a,b},[{a,a,b,b},{a,b,b,b}]}]
```

product(TupleOfSets) -> Relation

Types:

Relation = relation()

TupleOfSets = tuple-of(set())

Returns the *Cartesian product* of the non-empty tuple of sets TupleOfSets. If (x[1], ..., x[n]) is an element of the n-ary relation Relation, then x[i] is drawn from element i of TupleOfSets.

```
1> S1 = sofs:set([a,b]),
S2 = sofs:set([1,2]),
S3 = sofs:set([x,y]),
P3 = sofs:product({S1,S2,S3}),
sofs:to_external(P3).
[{a,1,x},{a,1,y},{a,2,x},{a,2,y},{b,1,x},{b,1,y},{b,2,x},{b,2,y}]
```

product(Set1, Set2) -> BinRel

Types:

BinRel = binary_relation()

Set1 = Set2 = set()

Returns the *Cartesian product* of Set1 and Set2.

```
1> S1 = sofs:set([1,2]),
S2 = sofs:set([a,b]),
R = sofs:product(S1, S2),
sofs:to_external(R).
[{1,a},{1,b},{2,a},{2,b}]
```

`product(S1, S2)` is equivalent to `product({S1, S2})`.

projection(SetFun, Set1) -> Set2

Types:

SetFun = set_fun()

Set1 = Set2 = set()

Returns the set created by substituting each element of Set1 by the result of applying SetFun to the element.

If SetFun is a number i >= 1 and Set1 is a relation, then the returned set is the *projection* of Set1 onto coordinate i.

```
1> S1 = sofs:from_term([{1,a},{2,b},{3,a}]),
S2 = sofs:projection(2, S1),
sofs:to_external(S2).
[a,b]
```

range(BinRel) -> Set

Types:

BinRel = binary_relation()

Set = set()

Returns the *range* of the binary relation BinRel.

```
1> R = sofs:relation([{1,a},{1,b},{2,b},{2,c}]),
S = sofs:range(R),
sofs:to_external(S).
[a,b,c]
```

relation(Tuples [, Type]) -> Relation

Types:

N = integer()

Type = N | type()

Relation = relation()

Tuples = [tuple()]

Creates a *relation*. `relation(R, T)` is equivalent to `from_term(R, T)`, if T is a *type* and the result is a relation. If Type is an integer N, then `[{atom, ..., atom}])`, where the size of the tuple is N, is used as type of the relation. If no type is explicitly given, the size of the first tuple of Tuples is used if there is such a tuple. `relation([])` is equivalent to `relation([], 2)`.

`relation_to_family(BinRel) -> Family`

Types:

Family = family()

BinRel = binary_relation()

Returns the *family* Family such that the index set is equal to the *domain* of the binary relation BinRel, and Family[i] is the *image* of the set of i under BinRel.

```
1> R = sofs:relation([{b,1},{c,2},{c,3}]),
F = sofs:relation_to_family(R),
sofs:to_external(F).
[{b,[1]},{c,[2,3]}]
```

`relative_product(TupleOfBinRels [, BinRel1]) -> BinRel2`

Types:

TupleOfBinRels = tuple-of(BinRel)

BinRel = BinRel1 = BinRel2 = binary_relation()

If TupleOfBinRels is a non-empty tuple {R[1], ..., R[n]} of binary relations and BinRel1 is a binary relation, then BinRel2 is the *relative product* of the ordered set (R[i], ..., R[n]) and BinRel1.

If BinRel1 is omitted, the relation of equality between the elements of the *Cartesian product* of the ranges of R[i], range R[1] × ... × range R[n], is used instead (intuitively, nothing is "lost").

```
1> TR = sofs:relation([{1,a},{1,aa},{2,b}]),
R1 = sofs:relation([{1,u},{2,v},{3,c}]),
R2 = sofs:relative_product({TR, R1}),
sofs:to_external(R2).
[{1,{a,u}},{1,{aa,u}},{2,{b,v}}]
```

Note that `relative_product({R1}, R2)` is different from `relative_product(R1, R2)`; the tuple of one element is not identified with the element itself.

`relative_product(BinRel1, BinRel2) -> BinRel3`

Types:

BinRel1 = BinRel2 = BinRel3 = binary_relation()

Returns the *relative product* of the binary relations BinRel1 and BinRel2.

`relative_product1(BinRel1, BinRel2) -> BinRel3`

Types:

BinRel1 = BinRel2 = BinRel3 = binary_relation()

Returns the *relative product* of the *converse* of the binary relation BinRel1 and the binary relation BinRel2.

```
1> R1 = sofs:relation([{1,a},{1,aa},{2,b}]),
R2 = sofs:relation([{1,u},{2,v},{3,c}]),
R3 = sofs:relative_product1(R1, R2),
sofs:to_external(R3).
[{a,u},{aa,u},{b,v}]
```

`relative_product1(R1, R2)` is equivalent to `relative_product(converse(R1), R2)`.

`restriction(BinRel1, Set) -> BinRel2`

Types:

BinRel1 = **BinRel2** = **binary_relation()**

Set = **set()**

Returns the *restriction* of the binary relation BinRel1 to Set.

```
1> R1 = sofs:relation([{1,a},{2,b},{3,c}]),
S = sofs:set([1,2,4]),
R2 = sofs:restriction(R1, S),
sofs:to_external(R2).
[{1,a},{2,b}]
```

`restriction(SetFun, Set1, Set2) -> Set3`

Types:

SetFun = **set_fun()**

Set1 = **Set2** = **Set3** = **set()**

Returns a subset of Set1 containing those elements that yield an element in Set2 as the result of applying SetFun.

```
1> S1 = sofs:relation([{1,a},{2,b},{3,c}]),
S2 = sofs:set([b,c,d]),
S3 = sofs:restriction(2, S1, S2),
sofs:to_external(S3).
[{2,b},{3,c}]
```

`set(Terms [, Type]) -> Set`

Types:

Set = **set()**

Terms = **[term()]**

Type = **type()**

Creates an *unordered set*. `set(L, T)` is equivalent to `from_term(L, T)`, if the result is an unordered set. If no *type* is explicitly given, `[atom]` is used as type of the set.

`specification(Fun, Set1) -> Set2`

Types:

Fun = **spec_fun()**

Set1 = **Set2** = **set()**

Returns the set containing every element of Set1 for which Fun returns `true`. If Fun is a tuple `{external, Fun2}`, Fun2 is applied to the *external set* of each element, otherwise Fun is applied to each element.

```
1> R1 = sofs:relation([{a,1},{b,2}]),
R2 = sofs:relation([{x,1},{x,2},{y,3}]),
S1 = sofs:from_sets([R1,R2]),
S2 = sofs:specification({sofs,is_a_function}, S1),
sofs:to_external(S2).
[[{a,1},{b,2}]]
```

strict_relation(BinRel1) -> BinRel2

Types:

BinRel1 = BinRel2 = binary_relation()

Returns the *strict relation* corresponding to the binary relation BinRel1.

```
1> R1 = sofs:relation([{1,1},{1,2},{2,1},{2,2}]),
R2 = sofs:strict_relation(R1),
sofs:to_external(R2).
[{1,2},{2,1}]
```

substitution(SetFun, Set1) -> Set2

Types:

SetFun = set_fun()

Set1 = Set2 = set()

Returns a function, the domain of which is Set1. The value of an element of the domain is the result of applying SetFun to the element.

```
1> L = [{a,1},{b,2}].
[{a,1},{b,2}]
2> sofs:to_external(sofs:projection(1,sofs:relation(L))).
[a,b]
3> sofs:to_external(sofs:substitution(1,sofs:relation(L))).
[{{a,1},a},{{b,2},b}]
4> SetFun = {external, fun({A,_}=E) -> {E,A} end},
sofs:to_external(sofs:projection(SetFun,sofs:relation(L))).
[{{a,1},a},{{b,2},b}]
```

The relation of equality between the elements of {a,b,c}:

```
1> I = sofs:substitution(fun(A) -> A end, sofs:set([a,b,c])),
sofs:to_external(I).
[{a,a},{b,b},{c,c}]
```

Let SetOfSets be a set of sets and BinRel a binary relation. The function that maps each element Set of SetOfSets onto the *image* of Set under BinRel is returned by this function:

```
```

```
images(SetOfSets, BinRel) ->
   Fun = fun(Set) -> sofs:image(BinRel, Set) end,
   sofs:substitution(Fun, SetOfSets).
```

Here might be the place to reveal something that was more or less stated before, namely that external unordered sets are represented as sorted lists. As a consequence, creating the image of a set under a relation R may traverse all elements of R (to that comes the sorting of results, the image). In `images/2`, BinRel will be traversed once for each element of SetOfSets, which may take too long. The following efficient function could be used instead under the assumption that the image of each element of SetOfSets under BinRel is non-empty:

```
images2(SetOfSets, BinRel) ->
   CR = sofs:canonical_relation(SetOfSets),
   R = sofs:relative_product1(CR, BinRel),
   sofs:relation_to_family(R).
```

symdiff(Set1, Set2) -> Set3

Types:

Set1 = Set2 = Set3 = set()

Returns the *symmetric difference* (or the Boolean sum) of Set1 and Set2.

```
1> S1 = sofs:set([1,2,3]),
S2 = sofs:set([2,3,4]),
P = sofs:symdiff(S1, S2),
sofs:to_external(P).
[1,4]
```

symmetric_partition(Set1, Set2) -> {Set3, Set4, Set5}

Types:

Set1 = Set2 = Set3 = Set4 = Set5 = set()

Returns a triple of sets: Set3 contains the elements of Set1 that do not belong to Set2; Set4 contains the elements of Set1 that belong to Set2; Set5 contains the elements of Set2 that do not belong to Set1.

to_external(AnySet) -> ExternalSet

Types:

ExternalSet = external_set()

AnySet = anyset()

Returns the *external set* of an atomic, ordered or unordered set.

to_sets(ASet) -> Sets

Types:

ASet = set() | ordset()

Sets = tuple_of(AnySet) | [AnySet]

Returns the elements of the ordered set ASet as a tuple of sets, and the elements of the unordered set ASet as a sorted list of sets without duplicates.

type(AnySet) -> Type

Types:

AnySet = anyset()

Type = type()

Returns the *type* of an atomic, ordered or unordered set.

union(SetOfSets) -> Set

Types:

Set = set()

SetOfSets = set_of_sets()

Returns the *union* of the set of sets SetOfSets.

union(Set1, Set2) -> Set3

Types:

Set1 = Set2 = Set3 = set()

Returns the *union* of Set1 and Set2.

union_of_family(Family) -> Set

Types:

Family = family()

Set = set()

Returns the union of the *family* Family.

```
1> F = sofs:family([{a,[0,2,4]},{b,[0,1,2]},{c,[2,3]}]),
S = sofs:union_of_family(F),
sofs:to_external(S).
[0,1,2,3,4]
```

weak_relation(BinRel1) -> BinRel2

Types:

BinRel1 = BinRel2 = binary_relation()

Returns a subset S of the *weak relation* W corresponding to the binary relation BinRel1. Let F be the *field* of BinRel1. The subset S is defined so that x S y if x W y for some x in F and for some y in F.

```
1> R1 = sofs:relation([{1,1},{1,2},{3,1}]),
R2 = sofs:weak_relation(R1),
sofs:to_external(R2).
[{1,1},{1,2},{2,2},{3,1},{3,3}]
```

See Also

dict(3), *digraph(3)*, *orddict(3)*, *ordsets(3)*, *sets(3)*

string

Erlang module

This module contains functions for string processing.

Exports

`len(String) -> Length`

Types:

String = **string()**

Length = **integer()**

Returns the number of characters in the string.

`equal(String1, String2) -> bool()`

Types:

String1 = **String2** = **string()**

Tests whether two strings are equal. Returns `true` if they are, otherwise `false`.

`concat(String1, String2) -> String3`

Types:

String1 = **String2** = **String3** = **string()**

Concatenates two strings to form a new string. Returns the new string.

`chr(String, Character) -> Index`

`rchr(String, Character) -> Index`

Types:

String = **string()**

Character = **char()**

Index = **integer()**

Returns the index of the first/last occurrence of `Character` in `String`. `0` is returned if `Character` does not occur.

`str(String, SubString) -> Index`

`rstr(String, SubString) -> Index`

Types:

String = **SubString** = **string()**

Index = **integer()**

Returns the position where the first/last occurrence of `SubString` begins in `String`. `0` is returned if `SubString` does not exist in `String`. For example:

```
> string:str(" Hello Hello World World ", "Hello World").
```

```
8
```

span(String, Chars) -> Length

cspan(String, Chars) -> Length

Types:

String = Chars = string()

Length = integer()

Returns the length of the maximum initial segment of String, which consists entirely of characters from (not from) Chars.

For example:

```
> string:span("\t    abcdef", " \t").
5
> string:cspan("\t    abcdef", " \t").
0
```

substr(String, Start) -> SubString

substr(String, Start, Length) -> Substring

Types:

String = SubString = string()

Start = Length = integer()

Returns a substring of `String`, starting at the position `Start`, and ending at the end of the string or at length `Length`.

For example:

```
> substr("Hello World", 4, 5).
"lo Wo"
```

tokens(String, SeparatorList) -> Tokens

Types:

String = SeparatorList = string()

Tokens = [string()]

Returns a list of tokens in `String`, separated by the characters in `SeparatorList`.

For example:

```
> tokens("abc defxxghix jkl", "x ").
["abc", "def", "ghi", "jkl"]
```

join(StringList, Separator) -> String

[Host,Port]=string:tokens(“ 127.0.0.1:8080",":").
{_,[Host,Port]}=regexp:split("127.0.0.1:8080",":")

Types:

regexp:split("127.0.0.1:8080",":").
{_,[Host,Port]}=regexp:split("127.0.0.1:8080",":")

12> string:join(["a","b","c","d"],"/").
"a/b/c/d"

StringList = [string()]

Separator = string()

Returns a string with the elements of `StringList` separated by the string in `Seperator`.

For example:

```
> join(["one", "two", "three"], ", ").
"one, two, three"
```

chars(Character, Number) -> String

chars(Character, Number, Tail) -> String

Types:

Character = char()

Number = integer()

String = string()

Returns a string consisting of `Number` of characters `Character`. Optionally, the string can end with the string `Tail`.

copies(String, Number) -> Copies

Types:

String = Copies = string()

Number = integer()

Returns a string containing `String` repeated `Number` times.

words(String) -> Count

words(String, Character) -> Count

Types:

String = string()

Character = char()

Count = integer()

Returns the number of words in `String`, separated by blanks or `Character`.

For example:

```
> words(" Hello old boy!", $o).
4
```

sub_word(String, Number) -> Word

sub_word(String, Number, Character) -> Word

Types:

String = Word = string()

Character = char()

Number = integer()

Returns the word in position `Number` of `String`. Words are separated by blanks or `Character`s.

For example:

```
> string:sub_word(" Hello old boy !",3,$o).
"ld b"
```

strip(String) -> Stripped

strip(String, Direction) -> Stripped

strip(String, Direction, Character) -> Stripped

```
20> string:strip("   12        ").
"12"
21> string:strip("1   12        1").
"1   12        1"
22> string:strip("   12        1").
"12        1"
23>
23> string:strip("   12        ",left).
"12        "
24> string:strip("   12        ",right).
"   12"
25> string:strip("   12        ",both).
"12"
26> string:strip("======12=======",both,$=).
"12"
```

Types:

String = Stripped = string()

Direction = left | right | both

Character = char()

Returns a string, where leading and/or trailing blanks or a number of `Character` have been removed. `Direction` can be `left`, `right`, or `both` and indicates from which direction blanks are to be removed. The function `strip/1` is equivalent to `strip(String, both)`.

For example:

```
> string:strip("...Hello.....", both, $.).
"Hello"
```

left(String, Number) -> Left

left(String, Number, Character) -> Left

Types:

String = Left = string()

Character = char

Number = integer()

Returns the `String` with the length adjusted in accordance with `Number`. The left margin is fixed. If the `length(String)` < `Number`, `String` is padded with blanks or `Character`s.

For example:

```
> string:left("Hello",10,$.).
"Hello....."
```

right(String, Number) -> Right

right(String, Number, Character) -> Right

Types:

String = Right = string()

Character = char

Number = integer()

Returns the `String` with the length adjusted in accordance with `Number`. The right margin is fixed. If the length of `(String)` < `Number`, `String` is padded with blanks or `Character`s.

For example:

```
> string:right("Hello", 10, $.).
".....Hello"
```

centre(String, Number) -> Centered

centre(String, Number, Character) -> Centered

Types:

String = Centered = string()

Character = char

Number = integer()

Returns a string, where `String` is centred in the string and surrounded by blanks or characters. The resulting string will have the length `Number`.

sub_string(String, Start) -> SubString

sub_string(String, Start, Stop) -> SubString

Types:

String = SubString = string()

Start = Stop = integer()

Returns a substring of `String`, starting at the position `Start` to the end of the string, or to and including the `Stop` position.

For example:

```
sub_string("Hello World", 4, 8).
"lo Wo"
```

to_float(String) -> {Float,Rest} | {error,Reason}

Types:

String = string()

Float = float()

Rest = string()

Reason = no_float | not_a_list

Argument `String` is expected to start with a valid text represented float (the digits being ASCII values). Remaining characters in the string after the float are returned in `Rest`.

Example:

```
> {F1,Fs} = string:to_float("1.0-1.0e-1"),
> {F2,[]} = string:to_float(Fs),
> F1+F2.
0.9
> string:to_float("3/2=1.5").
{error,no_float}
> string:to_float("-1.5eX").
{-1.5,"eX"}
```

to_integer(String) -> {Int,Rest} | {error,Reason}

Types:

String = string()

Int = integer()

Rest = string()

Reason = no_integer | not_a_list

Argument `String` is expected to start with a valid text represented integer (the digits being ASCII values). Remaining characters in the string after the integer are returned in `Rest`.

Example:

```
> {I1,Is} = string:to_integer("33+22"),
> {I2,[]} = string:to_integer(Is),
> I1-I2.
11
> string:to_integer("0.5").
{0,".5"}
> string:to_integer("x=2").
{error,no_integer}
```

to_lower(String) -> Result

to_lower(Char) -> CharResult

to_upper(String) -> Result

to_upper(Char) -> CharResult

Types:

String = Result = string()

Char = CharResult = integer()

The given string or character is case-converted. Note that the supported character set is ISO/IEC 8859-1 (a.k.a. Latin 1), all values outside this set is unchanged

Notes

Some of the general string functions may seem to overlap each other. The reason for this is that this string package is the combination of two earlier packages and all the functions of both packages have been retained.

> ### Note:
>
> Any undocumented functions in `string` should not be used.

supervisor

Erlang module
supervisor本身就是一个gen_server

A behaviour module for implementing a supervisor, a process which supervises other processes called child processes. A child process can either be another supervisor or a worker process. Worker processes are normally implemented using one of the `gen_event`, `gen_fsm`, or `gen_server` behaviours. A supervisor implemented using this module will have a standard set of interface functions and include functionality for tracing and error reporting. Supervisors are used to build an hierarchical process structure called a supervision tree, a nice way to structure a fault tolerant application. Refer to *OTP Design Principles* for more information.

A supervisor assumes the definition of which child processes to supervise to be located in a callback module exporting a pre-defined set of functions.

Unless otherwise stated, all functions in this module will fail if the specified supervisor does not exist or if bad arguments are given.

Supervision Principles

worker进程

The supervisor is responsible for starting, stopping and monitoring its child processes. The basic idea of a supervisor is that it should keep its child processes alive by restarting them when necessary.

The children of a supervisor is defined as a list of *child specifications*. When the supervisor is started, the child processes are started in order from left to right according to this list. When the supervisor terminates, it first terminates its child processes in reversed start order, from right to left.

A supervisor can have one of the following *restart strategies*:

- `one_for_one` - if one child process terminates and should be restarted, only that child process is affected.
- `one_for_all` - if one child process terminates and should be restarted, all other child processes are terminated and then all child processes are restarted.
- `rest_for_one` - if one child process terminates and should be restarted, the 'rest' of the child processes -- i.e. the child processes after the terminated child process in the start order -- are terminated. Then the terminated child process and all child processes after it are restarted.
- `simple_one_for_one` - a simplified `one_for_one` supervisor, where all child processes are dynamically added instances of the same process type, i.e. running the same code.

  The functions `terminate_child/2`, `delete_child/2` and `restart_child/2` are invalid for `simple_one_for_one` supervisors and will return `{error,simple_one_for_one}` if the specified supervisor uses this restart strategy.

To prevent a supervisor from getting into an infinite loop of child process terminations and restarts, a *maximum restart frequency* is defined using two integer values `MaxR` and `MaxT`. If more than `MaxR` restarts occur within `MaxT` seconds, the supervisor terminates all child processes and then itself. 先终止所有的子进程 然后自杀

This is the type definition of a child specification:

避免进入一个死掉 重启的死循环,所以会有一个重启策略的参数

```
child_spec() = {Id,StartFunc,Restart,Shutdown,Type,Modules}
 Id = term()
 StartFunc = {M,F,A}
  M = F = atom()
  A = [term()]
 Restart = permanent | transient | temporary
 Shutdown = brutal_kill | int()>=0 | infinity
 Type = worker | supervisor
 Modules = [Module] | dynamic
```

重启 关闭 类型 依赖什么模块

```
Module = atom()
```

- `Id` is a name that is used to identify the child specification internally by the supervisor.
- `StartFunc` defines the function call used to start the child process. It should be a module-function-arguments tuple `{M,F,A}` used as `apply(M,F,A)`.

  The start function *must create and link to* the child process, and should return `{ok,Child}` or `{ok,Child,Info}` where `Child` is the pid of the child process and `Info` an arbitrary term which is ignored by the supervisor.

  The start function can also return `ignore` if the child process for some reason cannot be started, in which case the child specification will be kept by the supervisor but the non-existing child process will be ignored.

  If something goes wrong, the function may also return an error tuple `{error,Error}`.

  Note that the `start_link` functions of the different behaviour modules fulfill the above requirements.
- `Restart` defines when a terminated child process should be restarted. A `permanent` child process should always be restarted, a `temporary` child process should never be restarted and a `transient` child process should be restarted only if it terminates abnormally, i.e. with another exit reason than `normal`.
- `Shutdown` defines how a child process should be terminated. `brutal_kill` means the child process will be unconditionally terminated using `exit(Child,kill)`. An integer timeout value means that the supervisor will tell the child process to terminate by calling `exit(Child,shutdown)` and then wait for an exit signal with reason `shutdown` back from the child process. If no exit signal is received within the specified number of milliseconds, the child process is unconditionally terminated using `exit(Child,kill)`.

  If the child process is another supervisor, `Shutdown` should be set to `infinity` to give the subtree ample time to shutdown.

  Important note on simple-one-for-one supervisors: The dynamically created child processes of a simple-one-for-one supervisor are not explicitly killed, regardless of shutdown strategy, but are expected to terminate when the supervisor does (that is, when an exit signal from the parent process is received).

  Note that all child processes implemented using the standard OTP behavior modules automatically adhere to the shutdown protocol.
- `Type` specifies if the child process is a supervisor or a worker.
- `Modules` is used by the release handler during code replacement to determine which processes are using a certain module. As a rule of thumb `Modules` should be a list with one element `[Module]`, where `Module` is the callback module, if the child process is a supervisor, gen_server or gen_fsm. If the child process is an event manager (gen_event) with a dynamic set of callback modules, `Modules` should be `dynamic`. See *OTP Design Principles* for more information about release handling.
- Internally, the supervisor also keeps track of the pid `Child` of the child process, or `undefined` if no pid exists.

Exports

`start_link(Module, Args) -> Result`

`start_link(SupName, Module, Args) -> Result`

Types:

SupName = {local,Name} | {global,Name}
Name = atom()
Module = atom()

Args = term()
Result = {ok,Pid} | ignore | {error,Error}
Pid = pid()
Error = {already_started,Pid}} | shutdown | term()

Creates a supervisor process as part of a supervision tree. The function will, among other things, ensure that the supervisor is linked to the calling process (its supervisor).

The created supervisor process calls `Module:init/1` to find out about restart strategy, maximum restart frequency and child processes. To ensure a synchronized start-up procedure, `start_link/2,3` does not return until `Module:init/1` has returned and all child processes have been started.

If `SupName={local,Name}` the supervisor is registered locally as `Name` using `register/2`. If `SupName={global,Name}` the supervisor is registered globally as `Name` using `global:register_name/2`. If no name is provided, the supervisor is not registered.

`Module` is the name of the callback module.

`Args` is an arbitrary term which is passed as the argument to `Module:init/1`.

If the supervisor and its child processes are successfully created (i.e. if all child process start functions return `{ok,Child}`, `{ok,Child,Info}`, or `ignore`) the function returns `{ok,Pid}`, where `Pid` is the pid of the supervisor. If there already exists a process with the specified `SupName` the function returns `{error, {already_started,Pid}}`, where `Pid` is the pid of that process.

If `Module:init/1` returns `ignore`, this function returns `ignore` as well and the supervisor terminates with reason `normal`. If `Module:init/1` fails or returns an incorrect value, this function returns `{error,Term}` where `Term` is a term with information about the error, and the supervisor terminates with reason `Term`.

If any child process start function fails or returns an error tuple or an erroneous value, the function returns `{error,shutdown}` and the supervisor terminates all started child processes and then itself with reason `shutdown`.

`start_child(SupRef, ChildSpec) -> Result`

Types:

SupRef = Name | {Name,Node} | {global,Name} | pid()
Name = Node = atom()
ChildSpec = child_spec() | [term()]
Result = {ok,Child} | {ok,Child,Info} | {error,Error}
Child = pid() | undefined
Info = term()
Error = already_present | {already_started,Child} | term()

Dynamically adds a child specification to the supervisor `SupRef` which starts the corresponding child process.

`SupRef` can be:

- the pid,
- `Name`, if the supervisor is locally registered,
- `{Name,Node}`, if the supervisor is locally registered at another node, or
- `{global,Name}`, if the supervisor is globally registered.

`ChildSpec` should be a valid child specification (unless the supervisor is a `simple_one_for_one` supervisor, see below). The child process will be started by using the start function as defined in the child specification.

```
start_client(Sock) ->
    {ok, Child} = supervisor:start_child(tcp_client_sup, []),
    ok = gen_tcp:controlling_process(Sock, Child),
    Child ! {go, Sock}.
```

注意:simple_one_for_one子进程的启动方式

If the case of a `simple_one_for_one` supervisor, the child specification defined in `Module:init/1` will be used and `ChildSpec` should instead be an arbitrary list of terms `List`. The child process will then be started by appending `List` to the existing start function arguments, i.e. by calling `apply(M, F, A++List)` where `{M,F,A}` is the start function defined in the child specification.

If there already exists a child specification with the specified `Id`, `ChildSpec` is discarded and the function returns `{error,already_present}` or `{error,{already_started,Child}}`, depending on if the corresponding child process is running or not.

If the child process start function returns `{ok,Child}` or `{ok,Child,Info}`, the child specification and pid is added to the supervisor and the function returns the same value.

If the child process start function returns `ignore`, the child specification is added to the supervisor, the pid is set to `undefined` and the function returns `{ok,undefined}`.

If the child process start function returns an error tuple or an erroneous value, or if it fails, the child specification is discarded and the function returns `{error,Error}` where `Error` is a term containing information about the error and child specification.

`terminate_child(SupRef, Id) -> Result`

Types:

SupRef = Name | {Name,Node} | {global,Name} | pid()

Name = Node = atom()

Id = term()

Result = ok | {error,Error}

Error = not_found | simple_one_for_one

Tells the supervisor `SupRef` to terminate the child process corresponding to the child specification identified by `Id`. The process, if there is one, is terminated but the child specification is kept by the supervisor. This means that the child process may be later be restarted by the supervisor. The child process can also be restarted explicitly by calling `restart_child/2`. Use `delete_child/2` to remove the child specification.

See `start_child/2` for a description of `SupRef`.

If successful, the function returns `ok`. If there is no child specification with the specified `Id`, the function returns `{error,not_found}`.

`delete_child(SupRef, Id) -> Result`

Types:

SupRef = Name | {Name,Node} | {global,Name} | pid()

Name = Node = atom()

Id = term()

Result = ok | {error,Error}

Error = running | not_found | simple_one_for_one

Tells the supervisor `SupRef` to delete the child specification identified by `Id`. The corresponding child process must not be running, use `terminate_child/2` to terminate it.

See `start_child/2` for a description of `SupRef`.

If successful, the function returns `ok`. If the child specification identified by `Id` exists but the corresponding child process is running, the function returns `{error,running}`. If the child specification identified by `Id` does not exist, the function returns `{error,not_found}`.

`restart_child(SupRef, Id) -> Result`

Types:

SupRef = Name | {Name,Node} | {global,Name} | pid()

Name = Node = atom()

Id = term()

Result = {ok,Child} | {ok,Child,Info} | {error,Error}

Child = pid() | undefined

Error = running | not_found | simple_one_for_one | term()

Tells the supervisor `SupRef` to restart a child process corresponding to the child specification identified by `Id`. The child specification must exist and the corresponding child process must not be running.

See `start_child/2` for a description of `SupRef`.

If the child specification identified by `Id` does not exist, the function returns `{error,not_found}`. If the child specification exists but the corresponding process is already running, the function returns `{error,running}`.

If the child process start function returns `{ok,Child}` or `{ok,Child,Info}`, the pid is added to the supervisor and the function returns the same value.

If the child process start function returns `ignore`, the pid remains set to `undefined` and the function returns `{ok,undefined}`.

If the child process start function returns an error tuple or an erroneous value, or if it fails, the function returns `{error,Error}` where `Error` is a term containing information about the error.

`which_children(SupRef) -> [{Id,Child,Type,Modules}]`

Types:

SupRef = Name | {Name,Node} | {global,Name} | pid()

Name = Node = atom()

Id = term() | undefined

Child = pid() | undefined

Type = worker | supervisor

Modules = [Module] | dynamic

Module = atom()

Returns a newly created list with information about all child specifications and child processes belonging to the supervisor `SupRef`.

Note that calling this function when supervising a large number of children under low memory conditions can cause an out of memory exception.

See `start_child/2` for a description of `SupRef`.

The information given for each child specification/process is:

- `Id` - as defined in the child specification or `undefined` in the case of a `simple_one_for_one` supervisor.
- `Child` - the pid of the corresponding child process, or `undefined` if there is no such process.
- `Type` - as defined in the child specification.
- `Modules` - as defined in the child specification.

`count_children(SupRef) -> PropListOfCounts`

Types:

SupRef = Name | {Name,Node} | {global,Name} | pid()

Name = Node = atom()

PropListOfCounts = [{specs, ChildSpecCount}, {active, ActiveProcessCount}, {supervisors, ChildSupervisorCount}, {workers, ChildWorkerCount}]

Returns a property list (see `proplists`) containing the counts for each of the following elements of the supervisor's child specifications and managed processes:

- `specs` - the total count of children, dead or alive.
- `active` - the count of all actively running child processes managed by this supervisor.
- `supervisors` - the count of all children marked as child_type = supervisor in the spec list, whether or not the child process is still alive.
- `workers` - the count of all children marked as child_type = worker in the spec list, whether or not the child process is still alive.

`check_childspecs([ChildSpec]) -> Result`

Types:

ChildSpec = child_spec()

Result = ok | {error,Error}

Error = term()

This function takes a list of child specification as argument and returns `ok` if all of them are syntactically correct, or `{error,Error}` otherwise.

CALLBACK FUNCTIONS

The following functions should be exported from a `supervisor` callback module.

Exports

`Module:init(Args) -> Result`

Types:

Args = term()

Result = {ok,{{RestartStrategy,MaxR,MaxT},[ChildSpec]}} | ignore

RestartStrategy = one_for_all | one_for_one | rest_for_one | simple_one_for_one

MaxR = MaxT = int()>=0

ChildSpec = child_spec()

Whenever a supervisor is started using `supervisor:start_link/2,3`, this function is called by the new process to find out about restart strategy, maximum restart frequency and child specifications.

`Args` is the `Args` argument provided to the start function.

`RestartStrategy` is the restart strategy and `MaxR` and `MaxT` defines the maximum restart frequency of the supervisor. `[ChildSpec]` is a list of valid child specifications defining which child processes the supervisor should start and monitor. See the discussion about Supervision Principles above.

Note that when the restart strategy is `simple_one_for_one`, the list of child specifications must be a list with one child specification only. (The `Id` is ignored). No child process is then started during the initialization phase, but all children are assumed to be started dynamically using `supervisor:start_child/2`.

The function may also return `ignore`.

SEE ALSO

gen_event(3), *gen_fsm(3)*, *gen_server(3)*, *sys(3)*

supervisor_bridge

Erlang module

A behaviour module for implementing a supervisor_bridge, a process which connects a subsystem not designed according to the OTP design principles to a supervision tree. The supervisor_bridge sits between a supervisor and the subsystem. It behaves like a real supervisor to its own supervisor, but has a different interface than a real supervisor to the subsystem. Refer to *OTP Design Principles* for more information.

A supervisor_bridge assumes the functions for starting and stopping the subsystem to be located in a callback module exporting a pre-defined set of functions.

The `sys` module can be used for debugging a supervisor_bridge.

Unless otherwise stated, all functions in this module will fail if the specified supervisor_bridge does not exist or if bad arguments are given.

Exports

`start_link(Module, Args) -> Result`

`start_link(SupBridgeName, Module, Args) -> Result`

Types:

SupBridgeName = {local,Name} | {global,Name}

Name = atom()

Module = atom()

Args = term()

Result = {ok,Pid} | ignore | {error,Error}

Pid = pid()

Error = {already_started,Pid} | term()

Creates a supervisor_bridge process, linked to the calling process, which calls `Module:init/1` to start the subsystem. To ensure a synchronized start-up procedure, this function does not return until `Module:init/1` has returned.

If `SupBridgeName={local,Name}` the supervisor_bridge is registered locally as `Name` using `register/2`. If `SupBridgeName={global,Name}` the supervisor_bridge is registered globally as `Name` using `global:register_name/2`. If no name is provided, the supervisor_bridge is not registered. If there already exists a process with the specified `SupBridgeName` the function returns `{error,{already_started,Pid}}`, where `Pid` is the pid of that process.

`Module` is the name of the callback module.

`Args` is an arbitrary term which is passed as the argument to `Module:init/1`.

If the supervisor_bridge and the subsystem are successfully started the function returns `{ok,Pid}`, where `Pid` is is the pid of the supervisor_bridge.

If `Module:init/1` returns `ignore`, this function returns `ignore` as well and the supervisor_bridge terminates with reason `normal`. If `Module:init/1` fails or returns an error tuple or an incorrect value, this function returns `{error,Term}` where `Term` is a term with information about the error, and the supervisor_bridge terminates with reason `Term`.

CALLBACK FUNCTIONS

The following functions should be exported from a `supervisor_bridge` callback module.

Exports

`Module:init(Args) -> Result`

Types:

Args = term()

Result = {ok,Pid,State} | ignore | {error,Error}

Pid = pid()

State = term()

Error = term()

Whenever a supervisor_bridge is started using `supervisor_bridge:start_link/2,3`, this function is called by the new process to start the subsystem and initialize.

`Args` is the `Args` argument provided to the start function.

The function should return `{ok,Pid,State}` where `Pid` is the pid of the main process in the subsystem and `State` is any term.

If later `Pid` terminates with a reason `Reason`, the supervisor bridge will terminate with reason `Reason` as well. If later the supervisor_bridge is stopped by its supervisor with reason `Reason`, it will call `Module:terminate(Reason,State)` to terminate.

If something goes wrong during the initialization the function should return `{error,Error}` where `Error` is any term, or `ignore`.

`Module:terminate(Reason, State)`

Types:

Reason = shutdown | term()

State = term()

This function is called by the supervisor_bridge when it is about to terminate. It should be the opposite of `Module:init/1` and stop the subsystem and do any necessary cleaning up. The return value is ignored.

`Reason` is `shutdown` if the supervisor_bridge is terminated by its supervisor. If the supervisor_bridge terminates because a a linked process (apart from the main process of the subsystem) has terminated with reason `Term`, `Reason` will be `Term`.

`State` is taken from the return value of `Module:init/1`.

SEE ALSO

supervisor(3), sys(3)

sys

Erlang module

模块包含了发送系统消息的方法和调试使用的消息.

This module contains functions for sending system messages used by programs, and messages used for debugging purposes.

Functions used for implementation of processes should also understand system messages such as debugging messages and code change. These functions must be used to implement the use of system messages for a process; either directly, or through standard behaviours, such as `gen_server`.

The following types are used in the functions defined below:

- `Name = pid() | atom() | {global, atom()}`
- `Timeout = int() >= 0 | infinity`
- `system_event() = {in, Msg} | {in, Msg, From} | {out, Msg, To} | term()`

The default timeout is 5000 ms, unless otherwise specified. The `timeout` defines the time period to wait for the process to respond to a request. If the process does not respond, the function evaluates `exit({timeout, {M, F, A}})`.

The functions make reference to a debug structure. The debug structure is a list of `dbg_opt()`. `dbg_opt()` is an internal data type used by the `handle_system_msg/6` function. No debugging is performed if it is an empty list.

System Messages

Processes which are not implemented as one of the standard behaviours must still understand system messages. There are three different messages which must be understood:

- Plain system messages. These are received as `{system, From, Msg}`. The content and meaning of this message are not interpreted by the receiving process module. When a system message has been received, the function `sys:handle_system_msg/6` is called in order to handle the request.
- Shutdown messages. If the process traps exits, it must be able to handle an shut-down request from its parent, the supervisor. The message `{'EXIT', Parent, Reason}` from the parent is an order to terminate. The process must terminate when this message is received, normally with the same `Reason` as `Parent`.
- There is one more message which the process must understand if the modules used to implement the process change dynamically during runtime. An example of such a process is the `gen_event` processes. This message is `{get_modules, From}`. The reply to this message is `From ! {modules, Modules}`, where `Modules` is a list of the currently active modules in the process.

  This message is used by the release handler to find which processes execute a certain module. The process may at a later time be suspended and ordered to perform a code change for one of its modules.

System Events

When debugging a process with the functions of this module, the process generates *system_events* which are then treated in the debug function. For example, `trace` formats the system events to the tty.

There are three predefined system events which are used when a process receives or sends a message. The process can also define its own system events. It is always up to the process itself to format these events.

Exports

log(Name,Flag)

log(Name,Flag,Timeout) -> ok | {ok, [system_event()]}

Types:

Flag = true | {true, N} | false | get | print

N = integer() > 0

Turns the logging of system events On or Off. If On, a maximum of `N` events are kept in the debug structure (the default is 10). If `Flag` is `get`, a list of all logged events is returned. If `Flag` is `print`, the logged events are printed to `standard_io`. The events are formatted with a function that is defined by the process that generated the event (with a call to `sys:handle_debug/4`).

log_to_file(Name,Flag)

log_to_file(Name,Flag,Timeout) -> ok | {error, open_file}

Types:

Flag = FileName | false

FileName = string()

Enables or disables the logging of all system events in textual format to the file. The events are formatted with a function that is defined by the process that generated the event (with a call to `sys:handle_debug/4`).

statistics(Name,Flag)

statistics(Name,Flag,Timeout) -> ok | {ok, Statistics}

Types:

Flag = true | false | get

Statistics = [{start_time, {Date1, Time1}}, {current_time, {Date, Time2}}, {reductions, integer()}, {messages_in, integer()}, {messages_out, integer()}]

Date1 = Date2 = {Year, Month, Day}

Time1 = Time2 = {Hour, Min, Sec}

Enables or disables the collection of statistics. If `Flag` is `get`, the statistical collection is returned.

trace(Name,Flag)

trace(Name,Flag,Timeout) -> void()

Types:

Flag = boolean()

Prints all system events on `standard_io`. The events are formatted with a function that is defined by the process that generated the event (with a call to `sys:handle_debug/4`).

no_debug(Name)

no_debug(Name,Timeout) -> void()

Turns off all debugging for the process. This includes functions that have been installed explicitly with the `install` function, for example triggers.

suspend(Name)

suspend(Name,Timeout) -> void()

Suspends the process. When the process is suspended, it will only respond to other system messages, but not other messages.

resume(Name)

resume(Name,Timeout) -> void()

Resumes a suspended process.

change_code(Name, Module, OldVsn, Extra)

change_code(Name, Module, OldVsn, Extra, Timeout) -> ok | {error, Reason}

Types:

OldVsn = undefined | term()

Module = atom()

Extra = term()

Tells the process to change code. The process must be suspended to handle this message. The `Extra` argument is reserved for each process to use as its own. The function `Mod:system_code_change/4` is called. `OldVsn` is the old version of the `Module`.

get_status(Name)

get_status(Name,Timeout) -> {status, Pid, {module, Mod}, [PDict, SysState, Parent, Dbg, Misc]}

我们在windows系统中调用的时候发现可能会卡死Shell 这个比较奇怪
是不是竞技场不能进入的原因?????

Types:

PDict = [{Key, Value}]

SysState = running | suspended

Parent = pid()

Dbg = [dbg_opt()]

Misc = term()

Gets the status of the process.

The value of `Misc` varies for different types of processes. For example, a `gen_server` process returns the callback module's state, and a `gen_fsm` process returns information such as its current state name. Callback modules for `gen_server` and `gen_fsm` can also customise the value of `Misc` by exporting a `format_status/2` function that contributes module-specific information; see *gen_server:format_status/2* and *gen_fsm:format_status/2* for more details.

install(Name,{Func,FuncState})

install(Name,{Func,FuncState},Timeout)

Types:

Func = dbg_fun()

dbg_fun() = fun(FuncState, Event, ProcState) -> done | NewFuncState

FuncState = term()

Event = system_event()

ProcState = term()

NewFuncState = term()

This function makes it possible to install other debug functions than the ones defined above. An example of such a function is a trigger, a function that waits for some special event and performs some action when the event is generated. This could, for example, be turning on low level tracing.

`Func` is called whenever a system event is generated. This function should return `done`, or a new func state. In the first case, the function is removed. It is removed if the function fails.

remove(Name,Func)

remove(Name,Func,Timeout) -> void()

Types:

Func = dbg_fun()

Removes a previously installed debug function from the process. `Func` must be the same as previously installed.

Process Implementation Functions

The following functions are used when implementing a special process. This is an ordinary process which does not use a standard behaviour, but a process which understands the standard system messages.

Exports

debug_options(Options) -> [dbg_opt()]

Types:

Options = [Opt]

Opt = trace | log | statistics | {log_to_file, FileName} | {install, {Func, FuncState}}

Func = dbg_fun()

FuncState = term()

This function can be used by a process that initiates a debug structure from a list of options. The values of the `Opt` argument are the same as the corresponding functions.

get_debug(Item,Debug,Default) -> term()

Types:

Item = log | statistics

Debug = [dbg_opt()]

Default = term()

This function gets the data associated with a debug option. `Default` is returned if the `Item` is not found. Can be used by the process to retrieve debug data for printing before it terminates.

`handle_debug([dbg_opt()],FormFunc,Extra,Event) -> [dbg_opt()]`

Types:

FormFunc = dbg_fun()

Extra = term()

Event = system_event()

This function is called by a process when it generates a system event. `FormFunc` is a formatting function which is called as `FormFunc(Device, Event, Extra)` in order to print the events, which is necessary if tracing is activated. `Extra` is any extra information which the process needs in the format function, for example the name of the process.

`handle_system_msg(Msg,From,Parent,Module,Debug,Misc)`

Types:

Msg = term()

From = pid()

Parent = pid()

Module = atom()

Debug = [dbg_opt()]

Misc = term()

This function is used by a process module that wishes to take care of system messages. The process receives a `{system, From, Msg}` message and passes the `Msg` and `From` to this function.

This function *never* returns. It calls the function `Module:system_continue(Parent, NDebug, Misc)` where the process continues the execution, or `Module:system_terminate(Reason, Parent, Debug, Misc)` if the process should terminate. The `Module` must export `system_continue/3`, `system_terminate/4`, and `system_code_change/4` (see below).

The `Misc` argument can be used to save internal data in a process, for example its state. It is sent to `Module:system_continue/3` or `Module:system_terminate/4`

`print_log(Debug) -> void()`

Types:

Debug = [dbg_opt()]

Prints the logged system events in the debug structure using `FormFunc` as defined when the event was generated by a call to `handle_debug/4`.

`Mod:system_continue(Parent, Debug, Misc)`

Types:

Parent = pid()

Debug = [dbg_opt()]

Misc = term()

This function is called from `sys:handle_system_msg/6` when the process should continue its execution (for example after it has been suspended). This function never returns.

Mod:system_terminate(Reason, Parent, Debug, Misc)

Types:

Reason = term()

Parent = pid()

Debug = [dbg_opt()]

Misc = term()

This function is called from `sys:handle_system_msg/6` when the process should terminate. For example, this function is called when the process is suspended and its parent orders shut-down. It gives the process a chance to do a clean-up. This function never returns.

Mod:system_code_change(Misc, Module, OldVsn, Extra) -> {ok, NMisc}

Types:

Misc = term()

OldVsn = undefined | term()

Module = atom()

Extra = term()

NMisc = term()

Called from `sys:handle_system_msg/6` when the process should perform a code change. The code change is used when the internal data structure has changed. This function converts the `Misc` argument to the new data structure. `OldVsn` is the *vsn* attribute of the old version of the `Module`. If no such attribute was defined, the atom `undefined` is sent.

timer

Erlang module

This module provides useful functions related to time. Unless otherwise stated, time is always measured in `milliseconds`. All timer functions return immediately, regardless of work carried out by another process.

Successful evaluations of the timer functions yield return values containing a timer reference, denoted `TRef` below. By using `cancel/1`, the returned reference can be used to cancel any requested action. A `TRef` is an Erlang term, the contents of which must not be altered.

The timeouts are not exact, but should be `at least` as long as requested.

Exports

`start() -> ok`

Starts the timer server. Normally, the server does not need to be started explicitly. It is started dynamically if it is needed. This is useful during development, but in a target system the server should be started explicitly. Use configuration parameters for `kernel` for this.

`apply_after(Time, Module, Function, Arguments) -> {ok, Tref} | {error, Reason}`

Types:

Time = integer() in Milliseconds

Module = Function = atom()

Arguments = [term()]

Evaluates `apply(M, F, A)` after `Time` amount of time has elapsed. Returns `{ok, TRef}`, or `{error, Reason}`.

`send_after(Time, Pid, Message) -> {ok, TRef} | {error,Reason}`

`send_after(Time, Message) -> {ok, TRef} | {error,Reason}`

Types:

Time = integer() in Milliseconds

Pid = pid() | atom()

Message = term()

Result = {ok, TRef} | {error, Reason}

`send_after/3`

Evaluates `Pid ! Message` after `Time` amount of time has elapsed. (`Pid` can also be an atom of a registered name.) Returns `{ok, TRef}`, or `{error, Reason}`.

`send_after/2`

Same as `send_after(Time, self(), Message)`.

exit_after(Time, Pid, Reason1) -> {ok, TRef} | {error,Reason2}

可否使用它来做玩家进程生命周期的延长?

exit_after(Time, Reason1) -> {ok, TRef} | {error,Reason2}

kill_after(Time, Pid)-> {ok, TRef} | {error,Reason2}

kill_after(Time) -> {ok, TRef} | {error,Reason2}

Types:

Time = integer() in milliseconds

Pid = pid() | atom()

Reason1 = Reason2 = term()

`exit_after/3`

Send an exit signal with reason `Reason1` to Pid `Pid`. Returns `{ok, TRef}`, or `{error, Reason2}`.

`exit_after/2`

Same as `exit_after(Time, self(), Reason1)`.

`kill_after/2`

Same as `exit_after(Time, Pid, kill)`.

`kill_after/1`

Same as `exit_after(Time, self(), kill)`.

apply_interval(Time, Module, Function, Arguments) -> {ok, TRef} | {error, Reason}

Types:

Time = integer() in milliseconds

Module = Function = atom()

Arguments = [term()]

Evaluates `apply(Module, Function, Arguments)` repeatedly at intervals of `Time`. Returns `{ok, TRef}`, or `{error, Reason}`.

send_interval(Time, Pid, Message) -> {ok, TRef} | {error, Reason}

send_interval(Time, Message) -> {ok, TRef} | {error, Reason}

Types:

Time = integer() in milliseconds

Pid = pid() | atom()

Message = term()

Reason = term()

{ok, TRef} = timer:send_interval(timer:seconds(1), doit)
在mochiweb的reloader里面使用了这个方法

进程终止的时候:
terminate(_Reason, State) ->
{ok, cancel} = timer:cancel(State#state.tref),
ok.

`send_interval/3`

Evaluates `Pid ! Message` repeatedly after `Time` amount of time has elapsed. (`Pid` can also be an atom of a registered name.) Returns `{ok, TRef}` or `{error, Reason}`.

`send_interval/2`

Same as `send_interval(Time, self(), Message)`.

`cancel(TRef) -> {ok, cancel} | {error, Reason}`

Cancels a previously requested timeout. `TRef` is a unique timer reference returned by the timer function in question. Returns `{ok, cancel}`, or `{error, Reason}` when `TRef` is not a timer reference.

`sleep(Time) -> ok`

Types:

Time = integer() in milliseconds or the atom infinity

Suspends the process calling this function for `Time` amount of milliseconds and then returns `ok`, or suspend the process forever if `Time` is the atom `infinity`. Naturally, this function does *not* return immediately.

`tc(Module, Function, Arguments) -> {Time, Value}`

`tc(Fun, Arguments) -> {Time, Value}`

Types:

Module = Function = atom()

Fun = fun()

Arguments = [term()]

Time = integer() in microseconds

Value = term()

`tc/3`

Evaluates `apply(Module, Function, Arguments)` and measures the elapsed real time as reported by `now/0`. Returns `{Time, Value}`, where `Time` is the elapsed real time in *microseconds*, and `Value` is what is returned from the apply.

`tc/2`

Evaluates `apply(Fun, Arguments)`. Otherwise works like `tc/3`.

`now_diff(T2, T1) -> Tdiff`

```
%%在mochiweb项目中这个方法被频繁使用
3> erlang:now().
{1337,668563,379435}
4> timer:now_diff(erlang:now(),v(3)).
21192126
5> timer:now_diff(erlang:now(),v(3)).
27216184
```

Types:

T1 = T2 = {MegaSecs, Secs, MicroSecs}

Tdiff = MegaSecs = Secs = MicroSecs = integer()

Calculates the time difference `Tdiff = T2 - T1` in *microseconds*, where `T1` and `T2` probably are timestamp tuples returned from `erlang:now/0`.

`seconds(Seconds) -> Milliseconds`

Returns the number of milliseconds in `Seconds`.

`minutes(Minutes) -> Milliseconds`

Return the number of milliseconds in `Minutes`.

hours(Hours) -> Milliseconds

Returns the number of milliseconds in `Hours`.

hms(Hours, Minutes, Seconds) -> Milliseconds

Returns the number of milliseconds in `Hours + Minutes + Seconds`.

Examples

This example illustrates how to print out "Hello World!" in 5 seconds:

```
1> timer:apply_after(5000, io, format, ["~nHello World!~n", []]).
{ok,TRef}
Hello World!
```

The following coding example illustrates a process which performs a certain action and if this action is not completed within a certain limit, then the process is killed.

```
Pid = spawn(mod, fun, [foo, bar]),
%% If pid is not finished in 10 seconds, kill him
{ok, R} = timer:kill_after(timer:seconds(10), Pid),
...
%% We change our mind...
timer:cancel(R),
...
```

WARNING

A timer can always be removed by calling `cancel/1`.

An interval timer, i.e. a timer created by evaluating any of the functions `apply_interval/4`, `send_interval/3`, and `send_interval/2`, is linked to the process towards which the timer performs its task.

A one-shot timer, i.e. a timer created by evaluating any of the functions `apply_after/4`, `send_after/3`, `send_after/2`, `exit_after/3`, `exit_after/2`, `kill_after/2`, and `kill_after/1` is not linked to any process. Hence, such a timer is removed only when it reaches its timeout, or if it is explicitly removed by a call to `cancel/1`.

unicode

Erlang module

This module contains functions for converting between different character representations. Basically it converts between iso-latin-1 characters and Unicode ditto, but it can also convert between different Unicode encodings (like UTF-8, UTF-16 and UTF-32).

The default Unicode encoding in Erlang is in binaries UTF-8, which is also the format in which built in functions and libraries in OTP expect to find binary Unicode data. In lists, Unicode data is encoded as integers, each integer representing one character and encoded simply as the Unicode codepoint for the character.

Other Unicode encodings than integers representing codepoints or UTF-8 in binaries are referred to as "external encodings". The iso-latin-1 encoding is in binaries and lists referred to as latin1-encoding.

It is recommended to only use external encodings for communication with external entities where this is required. When working inside the Erlang/OTP environment, it is recommended to keep binaries in UTF-8 when representing Unicode characters. Latin1 encoding is supported both for backward compatibility and for communication with external entities not supporting Unicode character sets.

DATA TYPES

```
unicode_binary() = binary() with characters encoded in UTF-8 coding standard
unicode_char() = integer() representing valid unicode codepoint

chardata() = charlist() | unicode_binary()

charlist() = [unicode_char() | unicode_binary() | charlist()]
  a unicode_binary is allowed as the tail of the list
```

```
external_unicode_binary() = binary()
  with characters coded in a user specified Unicode encoding other
  than UTF-8 (UTF-16 or UTF-32)

external_chardata() = external_charlist() | external_unicode_binary()

external_charlist() = [unicode_char() | external_unicode_binary() | external_charlist()]
  an external_unicode_binary is allowed as the tail of the list
```

```
latin1_binary() = binary() with characters coded in iso-latin-1
latin1_char() = integer() representing valid latin1 character (0-255)

latin1_chardata() = latin1_charlist() | latin1_binary()

latin1_charlist() = [latin1_char() | latin1_binary() | latin1_charlist()]
  a latin1_binary is allowed as the tail of the list
```

Exports

`bom_to_encoding(Bin) -> {Encoding,Length}`

Types:

Bin = binary() of byte_size 4 or more

Encoding = latin1 | utf8 | {utf16,little} | {utf16,big} | {utf32,little} | {utf32,big}

Length = int()

Check for a UTF byte order mark (BOM) in the beginning of a binary. If the supplied binary `Bin` begins with a valid byte order mark for either UTF-8, UTF-16 or UTF-32, the function returns the encoding identified along with the length of the BOM in bytes.

If no BOM is found, the function returns `{latin1,0}`

`characters_to_list(Data) -> list() | {error, list(), RestData} | {incomplete, list(), binary()}`

Types:

Data = latin1_chardata() | chardata() | external_chardata()

RestData = latin1_chardata() | chardata() | external_chardata()

Same as characters_to_list(Data,unicode).

`characters_to_list(Data, InEncoding) -> list() | {error, list(), RestData} | {incomplete, list(), binary()}`

Types:

Data = latin1_chardata() | chardata() | external_chardata()

RestData = latin1_chardata() | chardata() | external_chardata()

InEncoding = latin1 | unicode | utf8 | utf16 | utf32 | {utf16,little} | {utf16,big} | {utf32,little} | {utf32,big}

This function converts a possibly deep list of integers and binaries into a list of integers representing unicode characters. The binaries in the input may have characters encoded as latin1 (0 - 255, one character per byte), in which case the `InEncoding` parameter should be given as `latin1`, or have characters encoded as one of the UTF-encodings, which is given as the `InEncoding` parameter. Only when the `InEncoding` is one of the UTF encodings, integers in the list are allowed to be grater than 255.

If `InEncoding` is `latin1`, the `Data` parameter corresponds to the `iodata()` type, but for `unicode`, the `Data` parameter can contain integers greater than 255 (unicode characters beyond the iso-latin-1 range), which would make it invalid as `iodata()`.

The purpose of the function is mainly to be able to convert combinations of unicode characters into a pure unicode string in list representation for further processing. For writing the data to an external entity, the reverse function *characters_to_binary/3* comes in handy.

The option `unicode` is an alias for `utf8`, as this is the preferred encoding for Unicode characters in binaries. `utf16` is an alias for `{utf16,big}` and `utf32` is an alias for `{utf32,big}`. The `big` and `little` atoms denote big or little endian encoding.

If for some reason, the data cannot be converted, either because of illegal unicode/latin1 characters in the list, or because of invalid UTF encoding in any binaries, an error tuple is returned. The error tuple contains the tag `error`, a list representing the characters that could be converted before the error occurred and a representation of the characters including and after the offending integer/bytes. The last part is mostly for debugging as it still constitutes a possibly deep and/or mixed list, not necessarily of the same depth as the original data. The error occurs when traversing the list and whatever's left to decode is simply returned as is.

However, if the input `Data` is a pure binary, the third part of the error tuple is guaranteed to be a binary as well.

Errors occur for the following reasons:

- Integers out of range - If `InEncoding` is `latin1`, an error occurs whenever an integer greater than 255 is found in the lists. If `InEncoding` is of a Unicode type, error occurs whenever an integer greater than

16#10FFFF (the maximum unicode character) or in the range 16#D800 to 16#DFFF (invalid unicode range) is found.

- UTF encoding incorrect - If InEncoding is one of the UTF types, the bytes in any binaries have to be valid in that encoding. Errors can occur for various reasons, including "pure" decoding errors (like the upper bits of the bytes being wrong), the bytes are decoded to a too large number, the bytes are decoded to a code-point in the invalid unicode range or encoding is "overlong", meaning that a number should have been encoded in fewer bytes. The case of a truncated UTF is handled specially, see the paragraph about incomplete binaries below. If InEncoding is latin1, binaries are always valid as long as they contain whole bytes, as each byte falls into the valid iso-latin-1 range.

A special type of error is when no actual invalid integers or bytes are found, but a trailing binary() consists of too few bytes to decode the last character. This error might occur if bytes are read from a file in chunks or binaries in other ways are split on non UTF character boundaries. In this case an incomplete tuple is returned instead of the error tuple. It consists of the same parts as the error tuple, but the tag is incomplete instead of error and the last element is always guaranteed to be a binary consisting of the first part of a (so far) valid UTF character.

If one UTF characters is split over two consecutive binaries in the Data, the conversion succeeds. This means that a character can be decoded from a range of binaries as long as the whole range is given as input without errors occurring. Example:

```
decode_data(Data) ->
    case unicode:characters_to_list(Data,unicode) of
        {incomplete,Encoded, Rest} ->
       More = get_some_more_data(),
Encoded ++ decode_data([Rest, More]);
 {error,Encoded,Rest} ->
       handle_error(Encoded,Rest);
        List ->
       List
    end.
```

Bit-strings that are not whole bytes are however not allowed, so a UTF character has to be split along 8-bit boundaries to ever be decoded.

If any parameters are of the wrong type, the list structure is invalid (a number as tail) or the binaries does not contain whole bytes (bit-strings), a badarg exception is thrown.

characters_to_binary(Data) -> binary() | {error, binary(), RestData} | {incomplete, binary(), binary()}

Types:

Data = latin1_chardata() | chardata() | external_chardata()

RestData = latin1_chardata() | chardata() | external_chardata()

Same as characters_to_binary(Data, unicode, unicode).

characters_to_binary(Data,InEncoding) -> binary() | {error, binary(), RestData} | {incomplete, binary(), binary()}

Types:

Data = latin1_chardata() | chardata() | external_chardata()

RestData = latin1_chardata() | chardata() | external_chardata()

InEncoding = latin1 | unicode | utf8 | utf16 | utf32 | {utf16,little} | {utf16,big} | {utf32,little} | {utf32,big}

Same as characters_to_binary(Data, InEncoding, unicode).

characters_to_binary(Data, InEncoding, OutEncoding) -> binary() | {error, binary(), RestData} | {incomplete, binary(), binary()}

Types:

Data = latin1_chardata() | chardata() | external_chardata()

RestData = latin1_chardata() | chardata() | external_chardata()

InEncoding = latin1 | unicode | utf8 | utf16 | utf32 | {utf16,little} | {utf16,big} | {utf32,little} | {utf32,big}

OutEncoding = latin1 | unicode | utf8 | utf16 | utf32| {utf16,little} | {utf16,big} | {utf32,little} | {utf32,big}

This function behaves as *characters_to_list/2*, but produces an binary instead of a unicode list. The `InEncoding` defines how input is to be interpreted if binaries are present in the `Data`, while `OutEncoding` defines in what format output is to be generated.

The option `unicode` is an alias for `utf8`, as this is the preferred encoding for Unicode characters in binaries. `utf16` is an alias for `{utf16,big}` and `utf32` is an alias for `{utf32,big}`. The `big` and `little` atoms denote big or little endian encoding.

Errors and exceptions occur as in *characters_to_list/2*, but the second element in the `error` or `incomplete` tuple will be a `binary()` and not a `list()`.

encoding_to_bom(InEncoding) -> Bin

Types:

Bin = binary() of byte_size 4 or less

InEncoding = latin1 | unicode | utf8 | utf16 | utf32 | {utf16,little} | {utf16,big} | {utf32,little} | {utf32,big}

Length = int()

Create an UTF byte order mark (BOM) as a binary from the supplied `InEncoding`. The BOM is, if supported at all, expected to be placed first in UTF encoded files or messages.

The function returns `<<>>` for the `latin1` encoding, there is no BOM for ISO-latin-1.

It can be noted that the BOM for UTF-8 is seldom used, and it is really not a *byte order* mark. There are obviously no byte order issues with UTF-8, so the BOM is only there to differentiate UTF-8 encoding from other UTF formats.

win32reg

Erlang module

`win32reg` provides read and write access to the registry on Windows. It is essentially a port driver wrapped around the Win32 API calls for accessing the registry.

The registry is a hierarchical database, used to store various system and software information in Windows. It is available in Windows 95 and Windows NT. It contains installation data, and is updated by installers and system programs. The Erlang installer updates the registry by adding data that Erlang needs.

The registry contains keys and values. Keys are like the directories in a file system, they form a hierarchy. Values are like files, they have a name and a value, and also a type.

Paths to keys are left to right, with sub-keys to the right and backslash between keys. (Remember that backslashes must be doubled in Erlang strings.) Case is preserved but not significant. Example: `"\\hkey_local_machine\\software\\Ericsson\\Erlang\\5.0"` is the key for the installation data for the latest Erlang release.

There are six entry points in the Windows registry, top level keys. They can be abbreviated in the `win32reg` module as:

```
Abbrev.          Registry key
=======          ============
hkcr             HKEY_CLASSES_ROOT
current_user     HKEY_CURRENT_USER
hkcu             HKEY_CURRENT_USER
local_machine    HKEY_LOCAL_MACHINE
hklm             HKEY_LOCAL_MACHINE
users            HKEY_USERS
hku              HKEY_USERS
current_config   HKEY_CURRENT_CONFIG
hkcc             HKEY_CURRENT_CONFIG
dyn_data         HKEY_DYN_DATA
hkdd             HKEY_DYN_DATA
```

The key above could be written as `"\\hklm\\software\\ericsson\\erlang\\5.0"`.

The `win32reg` module uses a current key. It works much like the current directory. From the current key, values can be fetched, sub-keys can be listed, and so on.

Under a key, any number of named values can be stored. They have name, and types, and data.

Currently, the `win32reg` module supports storing only the following types: REG_DWORD, which is an integer, REG_SZ which is a string and REG_BINARY which is a binary. Other types can be read, and will be returned as binaries.

There is also a "default" value, which has the empty string as name. It is read and written with the atom `default` instead of the name.

Some registry values are stored as strings with references to environment variables, e.g. `"%SystemRoot%Windows"`. `SystemRoot` is an environment variable, and should be replaced with its value. A function `expand/1` is provided, so that environment variables surrounded in % can be expanded to their values.

For additional information on the Windows registry consult the Win32 Programmer's Reference.

Exports

change_key(RegHandle, Key) -> ReturnValue

Types:

RegHandle = term()

Key = string()

Changes the current key to another key. Works like cd. The key can be specified as a relative path or as an absolute path, starting with \.

change_key_create(RegHandle, Key) -> ReturnValue

Types:

RegHandle = term()

Key = string()

Creates a key, or just changes to it, if it is already there. Works like a combination of `mkdir` and `cd`. Calls the Win32 API function `RegCreateKeyEx()`.

The registry must have been opened in write-mode.

close(RegHandle)-> ReturnValue

Types:

RegHandle = term()

Closes the registry. After that, the `RegHandle` cannot be used.

current_key(RegHandle) -> ReturnValue

Types:

RegHandle = term()

ReturnValue = {ok, string()}

Returns the path to the current key. This is the equivalent of `pwd`.

Note that the current key is stored in the driver, and might be invalid (e.g. if the key has been removed).

delete_key(RegHandle) -> ReturnValue

Types:

RegHandle = term()

ReturnValue = ok | {error, ErrorId}

Deletes the current key, if it is valid. Calls the Win32 API function `RegDeleteKey()`. Note that this call does not change the current key, (unlike `change_key_create/2`.) This means that after the call, the current key is invalid.

delete_value(RegHandle, Name) -> ReturnValue

Types:

RegHandle = term()

ReturnValue = ok | {error, ErrorId}

Deletes a named value on the current key. The atom `default` is used for the the default value.

The registry must have been opened in write-mode.

`expand(String) -> ExpandedString`

Types:

String = string()

ExpandedString = string()

Expands a string containing environment variables between percent characters. Anything between two % is taken for a environment variable, and is replaced by the value. Two consecutive % is replaced by one %.

A variable name that is not in the environment, will result in an error.

`format_error(ErrorId) -> ErrorString`

Types:

ErrorId = atom()

ErrorString = string()

Convert an POSIX errorcode to a string (by calling `erl_posix_msg:message`).

`open(OpenModeList)-> ReturnValue`

Types:

OpenModeList = [OpenMode]

OpenMode = read | write

Opens the registry for reading or writing. The current key will be the root (`HKEY_CLASSES_ROOT`). The `read` flag in the mode list can be omitted.

Use `change_key/2` with an absolute path after `open`.

`set_value(RegHandle, Name, Value) -> ReturnValue`

Types:

Name = string() | default

Value = string() | integer() | binary()

Sets the named (or default) value to value. Calls the Win32 API function `RegSetValueEx()`. The value can be of three types, and the corresponding registry type will be used. Currently the types supported are: `REG_DWORD` for integers, `REG_SZ` for strings and `REG_BINARY` for binaries. Other types cannot currently be added or changed.

The registry must have been opened in write-mode.

`sub_keys(RegHandle) -> ReturnValue`

Types:

ReturnValue = {ok, SubKeys} | {error, ErrorId}

SubKeys = [SubKey]

SubKey = string()

Returns a list of subkeys to the current key. Calls the Win32 API function `EnumRegKeysEx()`.

Avoid calling this on the root keys, it can be slow.

`value(RegHandle, Name) -> ReturnValue`

Types:

Name = string() | default

ReturnValue = {ok, Value}

Value = string() | integer() | binary()

Retrieves the named value (or default) on the current key. Registry values of type `REG_SZ`, are returned as strings. Type `REG_DWORD` values are returned as integers. All other types are returned as binaries.

`values(RegHandle) -> ReturnValue`

Types:

ReturnValue = {ok, ValuePairs} | {ok, ErrorId}

ValuePairs = [ValuePair]

ValuePair = {Name, Value}

Name = string | default

Value = string() | integer() | binary()

Retrieves a list of all values on the current key. The values have types corresponding to the registry types, see `value`. Calls the Win32 API function `EnumRegValuesEx()`.

SEE ALSO

Win32 Programmer's Reference (from Microsoft)

`erl_posix_msg`

The Windows 95 Registry (book from O'Reilly)

zip

Erlang module

The `zip` module archives and extract files to and from a zip archive. The zip format is specified by the "ZIP Appnote.txt" file available on PKWare's website www.pkware.com.

The zip module supports zip archive versions up to 6.1. However, password-protection and Zip64 is not supported.

By convention, the name of a zip file should end in "`.zip`". To abide to the convention, you'll need to add "`.zip`" yourself to the name.

Zip archives are created with the *zip/2* or the *zip/3* function. (They are also available as `create`, to resemble the `erl_tar` module.)

To extract files from a zip archive, use the *unzip/1* or the *unzip/2* function. (They are also available as `extract`.)

To fold a function over all files in a zip archive, use the *foldl_3*.

To return a list of the files in a zip archive, use the *list_dir/1* or the *list_dir/2* function. (They are also available as `table`.)

To print a list of files to the Erlang shell, use either the *t/1* or *tt/1* function.

In some cases, it is desirable to open a zip archive, and to unzip files from it file by file, without having to reopen the archive. The functions *zip_open*, *zip_get*, *zip_list_dir* and *zip_close* do this.

LIMITATIONS

Zip64 archives are not currently supported.

Password-protected and encrypted archives are not currently supported

Only the DEFLATE (zlib-compression) and the STORE (uncompressed data) zip methods are supported.

The size of the archive is limited to 2 G-byte (32 bits).

Comments for individual files is not supported when creating zip archives. The zip archive comment for the whole zip archive is supported.

There is currently no support for altering an existing zip archive. To add or remove a file from an archive, the whole archive must be recreated.

DATA TYPES

```
zip_file()
```

The record `zip_file` contains the following fields.

`name = string()`

the name of the file

`info = file_info()`

file info as in *file:read_file_info/1*

`comment = string()`

the comment for the file in the zip archive

`offset = integer()`

the offset of the file in the zip archive (used internally)

`comp_size = integer()`

the compressed size of the file (the uncompressed size is found in `info`)

```
zip_comment
```

The record `zip_comment` just contains the archive comment for a zip archive

`comment = string()`

the comment for the zip archive

Exports

`zip(Name, FileList) -> RetValue`

`zip(Name, FileList, Options) -> RetValue`

`create(Name, FileList) -> RetValue`

`create(Name, FileList, Options) -> RetValue`

Types:

Name = filename()

FileList = [FileSpec]

FileSpec = filename() | {filename(), binary()} | {filename(), binary(), #file_info{}}

Options = [Option]

Option = memory | cooked | verbose | {comment, Comment} | {cwd, CWD} | {compress, What} | {uncompress, What}

What = all | [Extension] | {add, [Extension]} | {del, [Extension]}

Extension = string()

Comment = CWD = string()

RetValue = {ok, Name} | {ok, {Name, binary()}} | {error, Reason}

Reason = term()

The `zip` function creates a zip archive containing the files specified in `FileList`.

As synonyms, the functions `create/2` and `create/3` are provided, to make it resemble the `erl_tar` module.

The file-list is a list of files, with paths relative to the current directory, they will be stored with this path in the archive. Files may also be specified with data in binaries, to create an archive directly from data.

Files will be compressed using the DEFLATE compression, as described in the Appnote.txt file. However, files will be stored without compression if they already are compressed. The `zip/2` and `zip/3` checks the file extension to see whether the file should be stored without compression. Files with the following extensions are not compressed: `.Z`, `.zip`, `.zoo`, `.arc`, `.lzh`, `.arj`.

It is possible to override the default behavior and explicitly control what types of files that should be compressed by using the `{compress, What}` and `{uncompress, What}` options. It is possible to have several `compress` and `uncompress` options. In order to trigger compression of a file, its extension must match with the `compress` condition and must not match the `uncompress` condition. For example if `compress` is set to `["gif", "jpg"]`

and `uncompress` is set to `["jpg"]`, only files with `"gif"` as extension will be compressed. No other files will be compressed.

The following options are available:

`cooked`

By default, the `open/2` function will open the zip file in `raw` mode, which is faster but does not allow a remote (erlang) file server to be used. Adding `cooked` to the mode list will override the default and open the zip file without the `raw` option. The same goes for the files added.

`verbose`

Print an informational message about each file being added.

`memory`

The output will not be to a file, but instead as a tuple `{FileName, binary()}`. The binary will be a full zip archive with header, and can be extracted with for instance `unzip/2`.

`{comment, Comment}`

Add a comment to the zip-archive.

`{cwd, CWD}`

Use the given directory as current directory, it will be prepended to file names when adding them, although it will not be in the zip-archive. (Acting like a file:set_cwd/1, but without changing the global cwd property.)

`{compress, What}`

Controls what types of files that will be compressed. It is by default set to `all`. The following values of `What` are allowed:

`all`

means that all files will be compressed (as long as they pass the `uncompress` condition).

`[Extension]`

means that only files with exactly these extensions will be compressed.

`{add,[Extension]}`

adds these extensions to the list of compress extensions.

`{del,[Extension]}`

deletes these extensions from the list of compress extensions.

`{uncompress, What}`

Controls what types of files that will be uncompressed. It is by default set to `[".Z",".zip",".zoo",".arc",".lzh",".arj"]`. The following values of `What` are allowed:

`all`

means that no files will be compressed.

`[Extension]`

means that files with these extensions will be uncompressed.

`{add,[Extension]}`

adds these extensions to the list of uncompress extensions.

`{del,[Extension]}`

deletes these extensions from the list of uncompress extensions.

unzip(Archive) -> RetValue

unzip(Archive, Options) -> RetValue

extract(Archive) -> RetValue

extract(Archive, Options) -> RetValue

Types:

Archive = filename() | binary()

Options = [Option]

Option = {file_list, FileList} | keep_old_files | verbose | memory | {file_filter, FileFilter} | {cwd, CWD}

FileList = [filename()]

FileBinList = [{filename(),binary()}]

FileFilter = fun(ZipFile) -> true | false

CWD = string()

ZipFile = zip_file()

RetValue = {ok,FileList} | {ok,FileBinList} | {error, Reason} | {error, {Name, Reason}}

Reason = term()

The `unzip/1` function extracts all files from a zip archive. The `unzip/2` function provides options to extract some files, and more.

If the `Archive` argument is given as a binary, the contents of the binary is assumed to be a zip archive, otherwise it should be a filename.

The following options are available:

`{file_list, FileList}`

By default, all files will be extracted from the zip archive. With the `{file_list,FileList}` option, the `unzip/2` function will only extract the files whose names are included in `FileList`. The full paths, including the names of all sub directories within the zip archive, must be specified.

`cooked`

By default, the `open/2` function will open the zip file in `raw` mode, which is faster but does not allow a remote (erlang) file server to be used. Adding `cooked` to the mode list will override the default and open zip file without the `raw` option. The same goes for the files extracted.

`keep_old_files`

By default, all existing files with the same name as file in the zip archive will be overwritten. With the `keep_old_files` option, the `unzip/2` function will not overwrite any existing files. Not that even with the `memory` option given, which means that no files will be overwritten, files existing will be excluded from the result.

`verbose`

Print an informational message as each file is being extracted.

`memory`

Instead of extracting to the current directory, the `memory` option will give the result as a list of tuples `{Filename, Binary}`, where `Binary` is a binary containing the extracted data of the file named `Filename` in the zip archive.

`{cwd, CWD}`

Use the given directory as current directory, it will be prepended to file names when extracting them from the zip-archive. (Acting like a file:set_cwd/1, but without changing the global cwd property.)

`foldl(Fun, Acc0, Archive) -> {ok, Acc1} | {error, Reason}`

Types:

Fun = fun(FileInArchive, GetInfo, GetBin, AccIn) -> AccOut

FileInArchive = filename()

GetInfo = fun() -> #file_info{}

GetBin = fun() -> binary()

Acc0 = Acc1 = AccIn = AccOut = term()

Archive = filename() | {filename(), binary()}

The `foldl/3` function calls `Fun(FileInArchive, GetInfo, GetBin, AccIn)` on successive files in the `Archive`, starting with `AccIn == Acc0`. `FileInArchive` is the name that the file has in the archive. `GetInfo` is a fun that returns info about the the file. `GetBin` returns the contents of the file. Both `GetInfo` and `GetBin` must be called within the `Fun`. Their behavior is undefined if they are called outside the context of the `Fun`. The `Fun` must return a new accumulator which is passed to the next call. `foldl/3` returns the final value of the accumulator. `Acc0` is returned if the archive is empty. It is not necessary to iterate over all files in the archive. The iteration may be ended prematurely in a controlled manner by throwing an exception.

For example:

```
> Name = "dummy.zip".
"dummy.zip"
> {ok, {Name, Bin}} = zip:create(Name, [{"foo", <<"FOO">>}, {"bar", <<"BAR">>}], [memory]).
{ok,{"dummy.zip",
     <<80,75,3,4,20,0,0,0,0,0,74,152,97,60,171,39,212,26,3,0,
       0,0,3,0,0,...>>}}
> {ok, FileSpec} = zip:foldl(fun(N, I, B, Acc) -> [{N, B(), I()} | Acc] end, [], {Name, Bin}).
{ok,[{"bar",<<"BAR">>,
      {file_info,3,regular,read_write,
                 {{2010,3,1},{19,2,10}},
                 {{2010,3,1},{19,2,10}},
                 {{2010,3,1},{19,2,10}},
                 54,1,0,0,0,0,0}},
     {"foo",<<"FOO">>,
      {file_info,3,regular,read_write,
                 {{2010,3,1},{19,2,10}},
                 {{2010,3,1},{19,2,10}},
                 {{2010,3,1},{19,2,10}},
                 54,1,0,0,0,0,0}}]}
> {ok, {Name, Bin}} = zip:create(Name, lists:reverse(FileSpec), [memory]).
{ok,{"dummy.zip",
     <<80,75,3,4,20,0,0,0,0,0,74,152,97,60,171,39,212,26,3,0,
       0,0,3,0,0,...>>}}
> catch zip:foldl(fun("foo", _, B, _) -> throw(B()); (_, _, _, Acc) -> Acc end, [], {Name, Bin}).
<<"FOO">>
```

list_dir(Archive) -> RetValue

list_dir(Archive, Options)

table(Archive) -> RetValue

table(Archive, Options)

Types:

Archive = filename() | binary()

RetValue = {ok, [Comment, Files]} | {error, Reason}

Comment = zip_comment()

Files = [zip_file()]

Options = [Option]

Option = cooked

Reason = term()

The `list_dir/1` function retrieves the names of all files in the zip archive `Archive`. The `list_dir/2` function provides options.

As synonyms, the functions `table/2` and `table/3` are provided, to make it resemble the `erl_tar` module.

The result value is the tuple `{ok, List}`, where `List` contains the zip archive comment as the first element.

The following options are available:

`cooked`

By default, the `open/2` function will open the zip file in `raw` mode, which is faster but does not allow a remote (erlang) file server to be used. Adding `cooked` to the mode list will override the default and open zip file without the `raw` option.

t(Archive)

Types:

Archive = filename() | binary() | ZipHandle

ZipHandle = pid()

The `t/1` function prints the names of all files in the zip archive `Archive` to the Erlang shell. (Similar to "`tar t`".)

tt(Archive)

Types:

Archive = filename() | binary()

The `tt/1` function prints names and information about all files in the zip archive `Archive` to the Erlang shell. (Similar to "`tar tv`".)

zip_open(Archive) -> {ok, ZipHandle} | {error, Reason}

zip_open(Archive, Options) -> {ok, ZipHandle} | {error, Reason}

Types:

Archive = filename() | binary()

Options = [Option]

Options = cooked | memory | {cwd, CWD}

CWD = string()

ZipHandle = pid()

The `zip_open` function opens a zip archive, and reads and saves its directory. This means that subsequently reading files from the archive will be faster than unzipping files one at a time with `unzip`.

The archive must be closed with `zip_close/1`.

`zip_list_dir(ZipHandle) -> Result | {error, Reason}`

Types:

Result = [ZipComment, ZipFile...]

ZipComment = #zip_comment{}

ZipFile = #zip_file{}

ZipHandle = pid()

The `zip_list_dir/1` function returns the file list of an open zip archive.

`zip_get(ZipHandle) -> {ok, [Result]} | {error, Reason}`

`zip_get(FileName, ZipHandle) -> {ok, Result} | {error, Reason}`

Types:

FileName = filename()

ZipHandle = pid()

Result = filename() | {filename(), binary()}

The `zip_get` function extracts one or all files from an open archive.

The files will be unzipped to memory or to file, depending on the options given to the `zip_open` function when the archive was opened.

`zip_close(ZipHandle) -> ok | {error, einval}`

Types:

ZipHandle = pid()

The `zip_close/1` function closes a zip archive, previously opened with `zip_open`. All resources are closed, and the handle should not be used after closing.