

# KLR Algebra 3

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# Card for Quiver Hecke Algebras

$$\text{Denote } q_{ij}(u, v) = \begin{cases} 0 & i = j \\ (v - u)^{\#\{i \rightarrow j\}} (u - v)^{\#\{j \rightarrow i\}} & i \neq j \end{cases}$$

Let  $R(\underline{\mathbf{v}})$  be the algebra generated by  $x_1, \dots, x_n, \tau_1, \dots, \tau_{n-1}, \mathbb{1}_{\mathbf{i}}$ , with  $n = |\underline{\mathbf{v}}|$  and  $\mathbf{i} \vdash \underline{\mathbf{v}}$  with relations

- ▶  $\mathbb{1}_{\mathbf{i}} \mathbb{1}_{\mathbf{j}} = \delta_{\mathbf{i}, \mathbf{j}} \mathbb{1}_{\mathbf{i}}$ .
- ▶  $x_k x_\ell = x_\ell x_k$
- ▶  $\tau_k \tau_\ell = \tau_\ell \tau_k \quad |k - \ell| > 1$ .
- ▶  $(\tau_k x_\ell - x_{s_k(\ell)} \tau_k) \mathbb{1}_{\mathbf{i}} = \delta_{\mathbf{i}_k, \mathbf{i}_{k+1}} (\delta_{k+1, \ell} - \delta_{k, \ell}) \mathbb{1}_{\mathbf{i}}$
- ▶  $\tau_\ell^2 \mathbb{1}_{\mathbf{i}} = q_{\ell, \ell+1}(x_\ell, x_{\ell+1}) \mathbb{1}_{\mathbf{i}}$
- ▶  $(\tau_{\ell+1} \tau_\ell \tau_{\ell+1} - \tau_\ell \tau_{\ell+1} \tau_\ell) \mathbb{1}_{\mathbf{i}} =$   

$$= \delta_{\mathbf{i}_\ell, \mathbf{i}_{\ell+2}} \frac{q_{\mathbf{i}_\ell \mathbf{i}_{\ell+1}}(x_\ell, x_{\ell+1}) - q_{\mathbf{i}_{\ell+2} \mathbf{i}_{\ell+1}}(x_{\ell+2}, x_{\ell+1})}{x_\ell - x_{\ell+2}} \mathbb{1}_{\mathbf{i}}.$$

This standard notation differs from mine (used in last two slides) by  $x_i \leftrightarrow -x_i$ .

# Card for Quantum Groups

For simplicity, we consider only simply-laced case. Let  $U_q(\mathfrak{g})$  be the algebra generated by  $E_i, F_i, K_i, K_i^{-1}$  with  $i \in \mathbb{I}$  with relations

$$\blacktriangleright K_i K_i^{-1} = K_i^{-1} K_i = 1$$

$$\blacktriangleright [E_i, F_j] = \delta_{ij} \frac{K_i - K_i^{-1}}{q - q^{-1}}$$

$$\blacktriangleright \begin{cases} K_i E_j K_j^{-1} = q^{i \cdot j} E_j \\ K_i F_j K_j^{-1} = q^{-i \cdot j} F_j \end{cases}$$

$$\blacktriangleright \begin{cases} E_i E_j = E_j E_i & i \cdot j = 0 \\ E_i^2 E_j - [2] E_i E_j E_i + E_j E_i^2 = 0 & i \cdot j = 1 \end{cases}$$

$$\blacktriangleright \begin{cases} F_i F_j = F_j F_i & i \cdot j = 0 \\ F_i^2 F_j - [2] F_i F_j F_i + F_j F_i^2 = 0 & i \cdot j = 1 \end{cases}$$

# Categorification Theorem

Let

$$[R\text{Proj}] = \bigoplus_{\mathfrak{v}} [R(\mathfrak{v}) \text{Proj}].$$

We proved the following categorification theorem.

**Theorem (Categorification)**

$$U_q^-(\mathfrak{g}) \xrightarrow{\sim} [R\text{Proj}] \quad F_i \mapsto [\Theta_i].$$

Here  $\Theta_i = R(\alpha_i) = \mathbb{Q}[x_1] \mathbb{1}_i$ .

## Some Abstract Nonsense

Consider the combinatorial category “*quiver Hecke category*”

$$\mathcal{H}(Q) = \left\{ \begin{array}{ll} \mathbf{Obj} : & \underline{i} \quad \text{some sequence of } \mathbb{I} \\ \mathbf{Mor} : & \downarrow \\ & \underline{j} \quad \text{elements from } \in \mathbb{1}_{\underline{j}} R(\underline{\mathbf{v}}) \mathbb{1}_{\underline{i}} \end{array} \right.$$

It is the same to the full subcategory of  $\{R(\underline{\mathbf{v}})\mathbb{1}_{\underline{i}}\}_{\underline{i}}$ . But this is more combinatorial.

This is a monoidal category over  $\mathbb{Q}$ . Furthermore, the *Karoubi envelope* (=idempotent completion) is equivalent to the direct sum of the category  $\bigoplus_{\underline{\mathbf{v}}} R(\underline{\mathbf{v}}) \text{Proj}$ .

$$\mathcal{H}(Q)^{\oplus} \cong \bigoplus_{\underline{\mathbf{v}}} R(\underline{\mathbf{v}}) \text{Proj}.$$

## Modified Version

- ▶ We use the modified quantum groups

$$\dot{U}_q(\mathfrak{g}) = \bigoplus_{\mu, \lambda \in \Lambda} \mathbb{1}_\mu \cdot U_q(\mathfrak{g}) \cdot \mathbb{1}_\lambda$$

$(\Lambda = \text{weight lattice})$

$$\begin{cases} \mathbb{1}_\lambda \mathbb{1}_\mu = \delta_{\lambda\mu} \mathbb{1}_\lambda \\ K_i \mathbb{1}_\lambda = \mathbb{1}_\lambda K_i = q^{\langle \alpha_i^\vee, \lambda \rangle} \mathbb{1}_\lambda \\ E_i \mathbb{1}_\lambda = \mathbb{1}_{\lambda + \alpha_i} E_i \\ F_i \mathbb{1}_\lambda = \mathbb{1}_{\lambda - \alpha_i} F_i \end{cases}$$

- ▶ For  $\dot{U}_q(\mathfrak{g})$ , there is no coproduct. But there is a global symmetric paring, see [Lusztig Chapter 26].

$$\begin{aligned} (\lambda, \mu) \neq (\lambda', \mu') &\Rightarrow \langle \mathbb{1}_\mu f \mathbb{1}_\lambda, \mathbb{1}_{\mu'} g \mathbb{1}_{\lambda'} \rangle = 0 \\ \langle x^+ \mathbb{1}_\lambda, y^+ \mathbb{1}_\mu \rangle &= \langle x, y \rangle_f & \langle xy, z \rangle &= \langle y, \rho(x)z \rangle & \langle x, y \rangle &= \langle y, x \rangle \end{aligned}$$

Here

$$\rho : \begin{cases} \mathbb{1}_{\lambda + \alpha_i} E_i \mathbb{1}_\lambda \mapsto q^{(\dots)} \mathbb{1}_\lambda F_i \mathbb{1}_{\lambda + \alpha_i} \\ \mathbb{1}_{\lambda + \alpha_i} F_i \mathbb{1}_\lambda \mapsto q^{(\dots)} \mathbb{1}_\lambda E_i \mathbb{1}_{\lambda + \alpha_i} \end{cases}$$

We can think  $U_q(\mathfrak{g})$  a category with only one point

$$U_q(\mathfrak{g}) = \left\{ \begin{array}{c} \bullet \\ \circlearrowleft E_i, F_i, K_i \end{array} \right\}$$

But  $\dot{U}_q(\mathfrak{g})$  a category with objects to be elements of  $\Lambda$

$$\dot{U}_q(\mathfrak{g}) = \left\{ \begin{array}{c} \dots \swarrow F_j \quad \dots \swarrow F_j \\ \dots \xleftrightarrow[E_i]{F_j} \mathbb{1}_{\lambda+\alpha_j} \xleftrightarrow[E_i]{F_j} \mathbb{1}_{\lambda+\alpha_i+\alpha_j} \xleftrightarrow[E_i]{F_i} \dots \\ \dots \swarrow F_j \quad \dots \swarrow F_j \\ \dots \xleftrightarrow[E_i]{F_i} \mathbb{1}_{\lambda} \xleftrightarrow[E_i]{F_j} \mathbb{1}_{\lambda+\alpha_i} \xleftrightarrow[E_i]{F_i} \dots \\ \dots \swarrow E_j \quad \dots \swarrow E_j \end{array} \right\}$$

## Remark

We think  $\dot{\mathbf{U}}_q(\mathfrak{g})$  as an algebra, also as a category

- ▶ If we just want to compute this category  $\dot{\mathbf{U}}_q(\mathfrak{g})$ , there is some diagrammatic approach. (In this “categorification”, relations are equalities, not isomorphisms)
- ▶ For  $\mathfrak{gl}_n$ , known as *ladder diagram*, with  $\mathfrak{sl}_N$ -web, it gives a diagrammatic approach to the Howe duality. It also provides knot invariants. REF REF
- ▶ But today, we will categorify  $\dot{\mathbf{U}}_q(\mathfrak{g})$ . (In this “categorification”, relations are isomorphisms, not equalities)



# Our Main Purpose

We will construct the category  $\mathcal{U}_q(\mathfrak{g})$ .

- ▶ Its object is all the monomials in  $\dot{\mathbf{U}}_q(\mathfrak{g})$

$$\mathbf{Obj}(\mathcal{U}_q(\mathfrak{g})) = \{E_i^a F_j^b \cdots \mathbb{1}_\lambda\}.$$

- ▶ The morphism will be constructed such that

$$\langle X, Y \rangle = \mathbf{dim} \operatorname{Hom}_{\mathcal{U}_q(\mathfrak{g})}(X, Y).$$

and categorify

$$\left\{ \begin{array}{l} \mathbb{1}_\lambda \mathbb{1}_\mu = \delta_{\lambda\mu} \mathbb{1}_\lambda \\ E_i \mathbb{1}_\lambda = \mathbb{1}_{\lambda + \alpha_i} E_i \\ F_i \mathbb{1}_\lambda = \mathbb{1}_{\lambda - \alpha_i} F_i \end{array} \right. \quad \left\{ \begin{array}{l} E_i E_j = E_j E_i \quad i \cdot j = 0 \\ E_i^2 E_j - [2] E_i E_j E_i + E_j E_i^2 = 0 \quad i \cdot j = 1 \end{array} \right.$$
$$\left\{ \begin{array}{l} F_i F_j = F_j F_i \quad i \cdot j = 0 \\ F_i^2 F_j - [2] F_i F_j F_i + F_j F_i^2 = 0 \quad i \cdot j = 1 \end{array} \right.$$
$$\left\{ \begin{array}{l} E_i F_j = F_j E_i \quad i \neq j, \\ E_i F_i \cdot \mathbb{1}_\lambda = F_i E_i \cdot \mathbb{1}_\lambda + [\langle \alpha_i^\vee, \lambda \rangle] \cdot \mathbb{1}_\lambda \quad i \neq j, \end{array} \right.$$

# Objects

We use the following notations

$$\mathbb{1}_\lambda = \lambda \quad E_i = \uparrow^i \quad F_i = \downarrow^i.$$

Since in  $\dot{\mathbf{U}}(\mathfrak{g})$ ,  $E_i$  and  $F_i$  does not exists singly, so more precisely,

$$\lambda' = \lambda + \alpha_i \quad \mathbb{1}_{\lambda'} E_i \mathbb{1}_\lambda = \lambda' \uparrow^i \lambda \quad \mathbb{1}_\lambda F_i \mathbb{1}_{\lambda'} = \lambda \downarrow^i \lambda'.$$

For example,

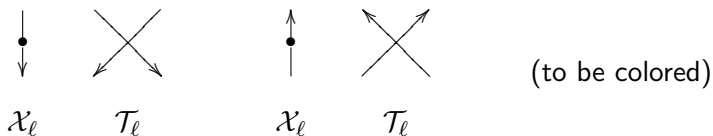
$$F_1 F_2 E_1 \mathbb{1}_\lambda = \lambda''' \downarrow^{\alpha_1} \lambda'' \downarrow^{\alpha_2} \lambda' \uparrow^{\alpha_1} \lambda$$
$$\begin{aligned} \lambda' &= \lambda + \alpha_1 \\ \lambda'' &= \lambda' - \alpha_2 \\ \lambda''' &= \lambda'' - \alpha_1 \end{aligned}$$

# Serre Relations

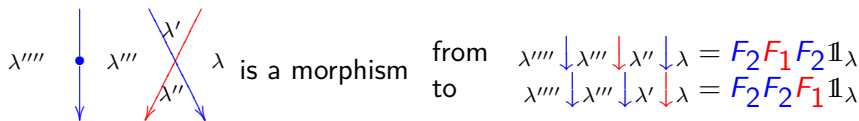
We impose

The relations in  $\mathcal{H}(Q)$  holds for  $\uparrow$  and  $\downarrow$  respectively

That is, we add the following into  $\text{Mor}(\mathcal{U}(\mathfrak{g}))$

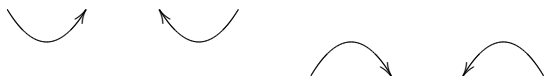


For example,



# Adjointness

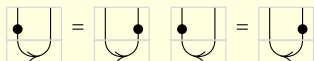
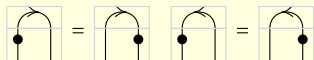
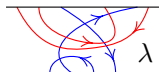
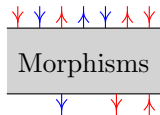
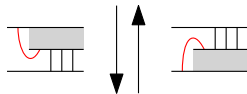
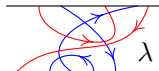
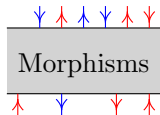
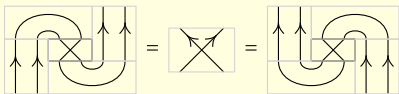
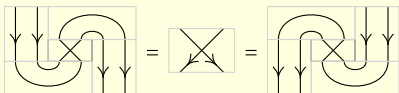
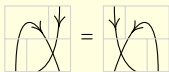
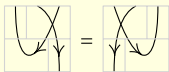
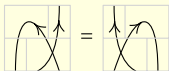
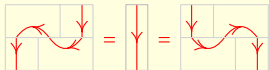
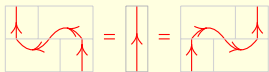
Since under the pairing  $\langle -, - \rangle$ ,  $E_i$  and  $F_i$  are in principle adjoint. To categorify them, we consider the unit and counit



For example,

$$\lambda'' \uparrow \quad \begin{array}{c} \lambda' \uparrow \\ \cup \\ \lambda \end{array} \text{ is a morphism from } \lambda'' \uparrow \lambda \downarrow \lambda' \uparrow \lambda \text{ to } \lambda'' \uparrow \lambda = E_2 \mathbb{1}_\lambda$$

$= E_2 F_1 E_1 \mathbb{1}_\lambda$



Well, there are more relations...

## The rest relation

Then we “force” the following morphisms to be invertible.

- ▶ When  $i \neq j$ , or  $i = j$  with  $\langle \alpha_i^\vee, \lambda \rangle = 0$

$$\begin{array}{c} j \\ \searrow \\ i \end{array} \begin{array}{c} i \\ \nearrow \\ j \end{array} \begin{array}{c} \lambda \\ \times \end{array} : E_i F_j \mathbb{1}_\lambda \longrightarrow F_j E_i \mathbb{1}_\lambda$$

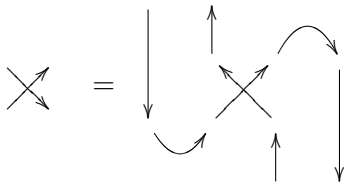
- ▶ When  $d := \langle \alpha_i^\vee, \lambda \rangle \geq 0$ ,

$$\begin{array}{c} i \\ \searrow \\ i \end{array} \begin{array}{c} i \\ \nearrow \\ i \end{array} \begin{array}{c} \lambda \\ \times \end{array} \oplus \bigoplus_{i=0}^{d-1} \begin{array}{c} \lambda \\ \curvearrowright \\ i \end{array} \begin{array}{c} \times \\ i \end{array} : E_i F_i \longrightarrow F_i E_i \mathbb{1}_\lambda \oplus \mathbb{1}_\lambda \oplus \cdots \oplus \mathbb{1}_\lambda$$

- ▶ When  $d := \langle \alpha_i^\vee, \lambda \rangle \leq 0$ ,

$$\begin{array}{c} i \\ \searrow \\ i \end{array} \begin{array}{c} i \\ \nearrow \\ i \end{array} \begin{array}{c} \lambda \\ \times \end{array} \oplus \bigoplus_{i=0}^{|d|-1} \begin{array}{c} \times \\ i \end{array} \begin{array}{c} \cup \\ i \end{array} : E_i F_i \mathbb{1}_\lambda \oplus \cdots \oplus \mathbb{1}_\lambda \longrightarrow F_i E_i \mathbb{1}_\lambda \oplus \cdots \oplus \mathbb{1}_\lambda$$

Where



Actually, it is proved that if we assume the morphisms to be invertible, then the inverse can also be expressed by a long list of diagrams and relations.

REF

Be careful, it even has imaginary “minus •” on bubbles.

# Degrees

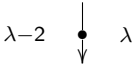
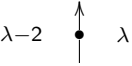
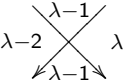

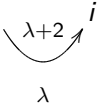



Let us determine the degrees

deg	2	2	$-i \cdot j$	$-i \cdot j$
deg	$1 + \langle \alpha_i, \lambda \rangle$	$1 + \langle \alpha_i, \lambda \rangle$	$1 - \langle \alpha_i, \lambda \rangle$	$1 - \langle \alpha_i, \lambda \rangle$



## Example of $\mathfrak{sl}_2$

In this case, let us assume  $\Lambda = \mathbb{Z}$ , and  $\alpha_1 = 2 \in \Lambda$ . Then  $\langle \alpha_1^\vee, b \rangle = b$ . Then

				
deg	2	2	-2	-2
				
deg	$1 + \lambda$	$1 + \lambda$	$1 - \lambda$	$1 - \lambda$

# Summary

▶ 
$$\begin{cases} \mathbb{1}_\lambda \mathbb{1}_\mu = \delta_{\lambda\mu} \mathbb{1}_\lambda \\ E_i \mathbb{1}_\lambda = \mathbb{1}_{\lambda+\alpha_i} E_i & F_i \mathbb{1}_\lambda = \mathbb{1}_{\lambda-\alpha_i} F_i \end{cases} \quad \text{By notation}$$

▶ 
$$\begin{cases} E_i E_j = E_j E_i & i \cdot j = 0 \\ E_i^2 E_j - [2] E_i E_j E_i + E_j E_i^2 = 0 & i \cdot j = 1 \\ F_i F_j = F_j F_i & i \cdot j = 0 \\ F_i^2 F_j - [2] F_i F_j F_i + F_j F_i^2 = 0 & i \cdot j = 1 \end{cases} \quad \text{By quiver-Hecke}$$

▶ 
$$\begin{cases} E_i F_j = F_j E_i & i \neq j, \\ E_i F_i \cdot \mathbb{1}_\lambda = F_i E_i \cdot \mathbb{1}_\lambda + [\langle \alpha_i^\vee, \lambda \rangle] \cdot \mathbb{1}_\lambda & i = j, \end{cases} \quad \text{By forcing}$$

▶ Actually, there is a combinatorial formula for monomial, summing over all homotopy class, see [KL, section 2.2].

From the construction above,  $[\mathcal{U}_q(\mathfrak{g})] = \mathbf{U}_q(\mathfrak{g})$ .

## The 2-category $\mathcal{U}_q(\mathfrak{g})$

But actually, this construction here is too “artificial”. Actually, the purpose of  $\mathcal{U}_q(\mathfrak{g})$  here is different from  $R(\underline{\mathbf{v}})$ .

The significance here is

think  $\dot{\mathbf{U}}_q(\mathfrak{g})$  as an algebra  $\implies$  Category  $\mathcal{U}_q(\mathfrak{g})$  categorify it.

think  $\dot{\mathbf{U}}_q(\mathfrak{g})$  as a category  $\implies$  2-Category  $\mathcal{U}_q(\mathfrak{g})$  categorify it.

(When mathematicians tried to categorify, Rouquier had started to categorify categories)

$$\mathcal{U}_q(\mathfrak{g}) = \left\{ \begin{array}{l} \mathbf{Obj} : \lambda \in \Lambda \\ \mathbf{Mor} : \downarrow_{\mu} \end{array} \right. = \left\{ \begin{array}{l} \mathbf{Obj} : \text{monomials in } E_i, F_i \\ \mathbf{Mor} : \text{spaces generated by diagrams} \end{array} \right.$$

# Representations

- ▶ If we think  $U_q(\mathfrak{g})$  as a one-point category with endomorphism ring  $U_q(\mathfrak{g})$ , then

$$U_q(\mathfrak{g})\text{-Rep} = \text{Fun}(U_q(\mathfrak{g}), \mathbb{k}\text{-Vec}).$$

Equivalently, we find one vector space with an action  $E_i, F_i, K_i$  satisfying the relations.

- ▶ Similarly,

$$\dot{U}_q(\mathfrak{g})\text{-Rep} = \text{Fun}(\dot{U}_q(\mathfrak{g}), \mathbb{k}\text{-Vec}).$$

Equivalently, we find a family of space parameterized by  $\Lambda$ , (i.e. weight space) with maps  $E_i, F_i$  between them.

## 2-Representations

- ▶ Then, we can define 2-representation

$$\mathcal{U}_q(\mathfrak{g})\text{-Rep} = 2\text{-Fun}(\dot{\mathbf{U}}_q(\mathfrak{g}), \mathbb{k}\text{-Cat}).$$

Equivalently, we find a family of categories parameterized by  $\Lambda$ , (i.e. weight space) with functor  $E_i, F_i$  between them; natural transformation (diagrams) between functors.

- ▶ But, can we really find 2-representations?

Yes! From geometry. Springer theory, Nakajima quiver variety, perverse sheaves, .... This is known as “categorification of representations”. REF REF

To do this, may be only do for a half of quantum group (quiver-Hecke algebra) is easier. This is known as “cyclotomic categorification conjecture”. REF

## References and Preview

- ▶ Brundan. On the definition of Kac-Moody 2-category. [arXiv]
- ▶ Kamnitzer. Categorification of Lie algebras [d'apres Rouquier, Khovanov-Lauda]. [arXiv]
- ▶ Rouquier. 2-Kac-Moody algebras. [arXiv]
- ▶ Khovanov, Lauda. A diagrammatic approach to categorification of quantum groups III [arXiv].

# General Remarks on Categorification

As far as I see, there are two kinds of *categorification*

Realization algebra as a  
“Grothendieck group” of  
category

$$\mathcal{H}eck_q(W) \cong \mathbf{SSP}erv_B(G/B)$$

$$U_q(\mathfrak{g})_{\deg=\underline{\mathbf{v}}}^+ \cong \mathbf{SSP}erv_{G(\underline{\mathbf{v}})}(E(\underline{\mathbf{v}}))$$

$$\mathcal{H}all(Q) \cong U_q(\mathfrak{g})^+ \Big|_{q=\mathbf{q}^{1/2}}$$

.....

Computation of a known category combinatorially (invariant theory)

$$\mathbf{SSB}im \cong \mathbf{SSP}erv_B(G/B)$$

$$R(\underline{\mathbf{v}})\text{-Proj} \cong \mathbf{SSP}erv_{G(\underline{\mathbf{v}})}(E(\underline{\mathbf{v}}))$$

$$\mathcal{T}\mathcal{L}^\oplus \cong U_q(\mathfrak{sl}_2)\text{-Rep}$$

.....

Of course, we do both of them —“diagrammatic algebra”.

## When realizing algebras, usually

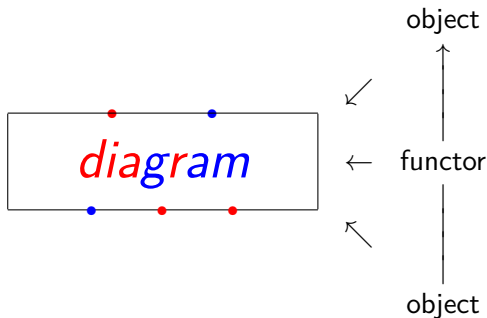
element	object
$+$	$\oplus$ or formal sum
$\cdot$	$\otimes$ , composition etc
$=$	$\cong$ , exact sequence, etc
$q$	degree shift
pairing	<b>dim</b> $\text{Hom}^\bullet(-, -)$
involution	duality




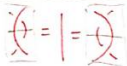
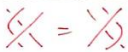
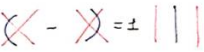


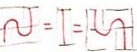
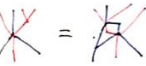






- ▶ **Subalgebra/subspace?** We may consider subcategory, for example, the category of modules over a quotient algebra.
- ▶ **Subalgebra fixed by a group  $G$  action?** We can consider the functor category from  $G$  to  $\mathcal{C}$ ,  $\text{Fun}(G, \mathcal{C})$ . For example, Lusztig uses similar idea to construct non-simply-laced quantum groups.
- ▶ **Quotient algebra/space?** Well, it adds more relation so we should consider the localization of the category. For example, cyclotomic KLR algebra.
- ▶ **Module  $\approx$  representation?** We can consider the 2-category generated by the functor  $X \otimes -$  and natural transform  $\phi \otimes -$  for all object  $X$  and functor  $\phi$ . For example, the 2-category  $\mathcal{U}\mathfrak{g}$ .

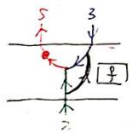
## When Computation of Categories, usually ...







- ▶ We only compute a category  $\mathcal{C}$  whose *Karoubi envelope* (=idempotent completion) is what we want.




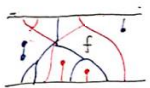
Usually, the diagram is a “cobordism” of the boundary.

 crossing	 (braid relation)	 (skein relation)	 (quiver-Hecke relation)
 cap/cup	 (adjointness)	 (Frobenius)	 (Elias relation)
 branching	$Y = Y$ (associativity)	$X = X = X$ (Frobenius)	
 trace	$Y =   = Y$	$X - X = (q - q^{-1})$ (skein relation)	
 dot	$X - X = \pm 1$	$\smile = X + q$ (skein relation)	
 identity	$\cup = \cap$	$\bigcirc = [2]$ (bubble)	



- objects morphisms
- colored 
  - oriented 
  - valued 
- functions
- 
  - 
  - 

- dot 
- label  $|s$
- chord  $\leftarrow$



# Reference for Diagrammatic Categorification

- [REF](#) The category of  $\mathfrak{sl}_2$ -representations (=Temperley–Lieb category). [bibibili]
- [REF](#) The category of  $\mathfrak{sl}_3$ -representations.
- [REF](#) The category of  $\mathfrak{sl}_N$ -representations (=  $\mathfrak{sl}_N$ -web).
- [REF](#) Heisenberg algebra. [bibibili]
- [REF](#) Realization of  $\dot{\mathbf{U}}_q(\mathfrak{gl}_n)$  (=ladder diagram)
- [REF](#) Realization of  $\dot{\mathbf{U}}_q(\mathfrak{g})$  (=KLR algebra). [bibibili]
- [REF](#) Realization of Hecke algebras (=Soergel bimodules). [bibibili, Youtube]